

# Hidden flaw jeopardizes millions of online transactions

Mathematicians discover weakness in commonly used encryption schemes

Jump to discuss

comments below

X

 Recommend 38

 Tweet 9

+1 3

 Share 20

Below:  Discuss  Related

By Paul Wagenseil



updated 2/15/2012 6:01:37 PM ET

Mathematicians based in Switzerland and the United States have discovered a small but the most commonly used digital encryption schemes, a flaw that could undermine the security of communication.

Topics ▾

News

In Depth

Reviews

Blogs ▾

Opini

## Blogs



### IT Blogwatch

*A Daily Digest of IT Blogs from Richi Jennings*

[More posts](#) | [Read bio](#)



February 15, 2012 - 6:00 A.M.

## RSA crypto: 'flawed', 'risky', 'quagmire of vulnerabilities'

2 Comments

Like < 5

+1 < 1

**TAGS:** [certificate](#), [crypto](#), [cryptography](#), [Diffie-Hellman](#), [encryption](#), [key management](#), [RSA](#), [SSL](#)

**IT TOPICS:** [Applications](#), [Cybercrime & Hacking](#), [Financial IT](#), [Government & Regulation](#), [Internet](#), [Privacy](#), [Security](#), [Security Hardware & Software](#)

**A group of six academic researchers have concluded that real-world RSA encryption keys are riskier than Diffie-Hellman-based ones. It seems that some of the random numbers used to generate them weren't, errm, random. In [IT Blogwatch](#), bloggers wonder if the sky is falling.**



Print



Comment



Tweet



Like < 1



Alert

## 'Predictably random' public keys can be cracked - crypto boffins Battling researchers argue over whether you should panic

By [John Leyden](#) • [Get more from this author](#)

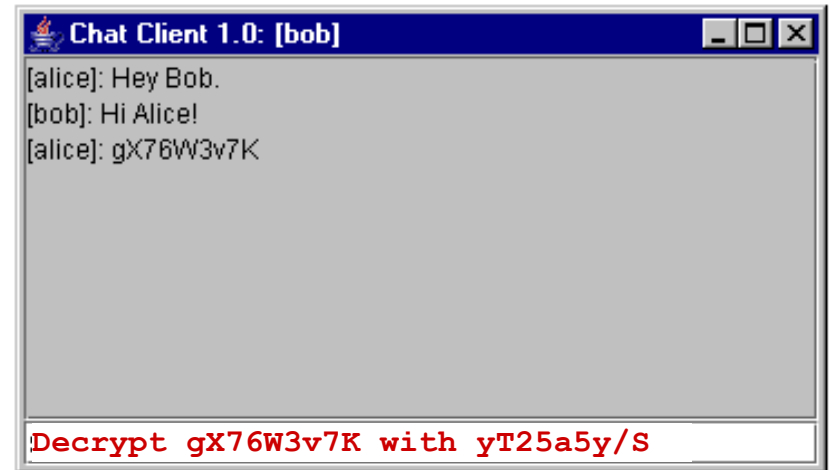
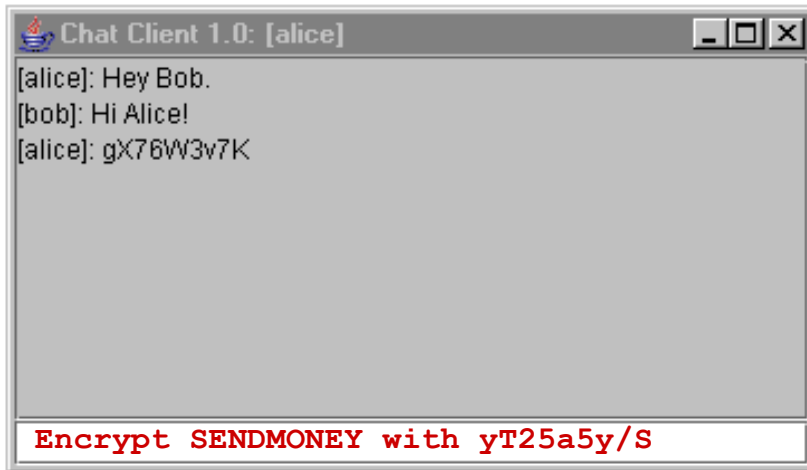
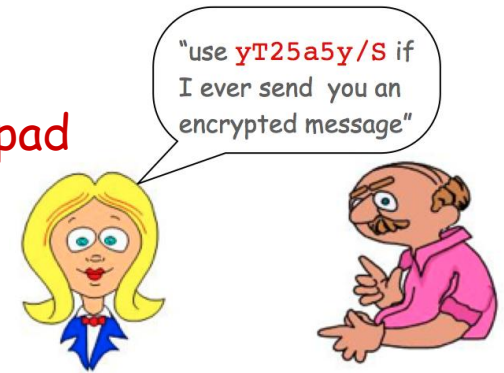
Posted in [Enterprise Security](#), 16th February 2012 13:38 GMT

**Analysis** Cryptography researchers have discovered flaws in the key generation that underpins the security of important cryptography protocols, including SSL.

# Secure Chat

Alice wants to send a secret message to Bob?

- Sometime in the past, they exchange a **one-time pad**
- Alice uses the pad to encrypt the message
- Bob uses the same pad to decrypt the message



**Key point** Without the pad, Eve cannot understand the message



# Encryption Machine

Goal Design a machine to encrypt and decrypt data

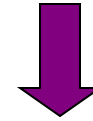
S	E	N	D	M	O	N	E	Y
---	---	---	---	---	---	---	---	---

g	x	7	6	w	3	v	7	k
---	---	---	---	---	---	---	---	---

S	E	N	D	M	O	N	E	Y
---	---	---	---	---	---	---	---	---



encrypt



decrypt

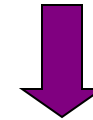
# Encryption Machine

**Goal** Design a machine to encrypt and decrypt data

S	E	N	D	M	O	N	E	Y
---	---	---	---	---	---	---	---	---

g	x	7	6	w	3	v	7	k
---	---	---	---	---	---	---	---	---

S	E	N	D	M	O	N	E	Y
---	---	---	---	---	---	---	---	---



encrypt



decrypt

## Enigma encryption machine

- "Unbreakable" German code during WWII
- Broken by Turing bombe
- One of first uses of computers
- Helped win Battle of Atlantic by locating U-boats





# A Digital World

Data is a sequence of bits [bit = 0 or 1]

- **Text**
- Programs, executables
- Documents, pictures, sounds, movies, ...

File formats txt, pdf, java, exe, docx, pptx, jpeg, mp3, divx, ...



computer with  
a lens



computer with  
earbuds



computer with  
a radio

# A Digital World

Data is a sequence of bits [bit = 0 or 1]

- **Text**
- Programs, executables
- Documents, pictures, sounds, movies, ...

File formats txt, pdf, java, exe, docx, pptx, jpeg, mp3, divx, ...



computer with  
a cash dispenser

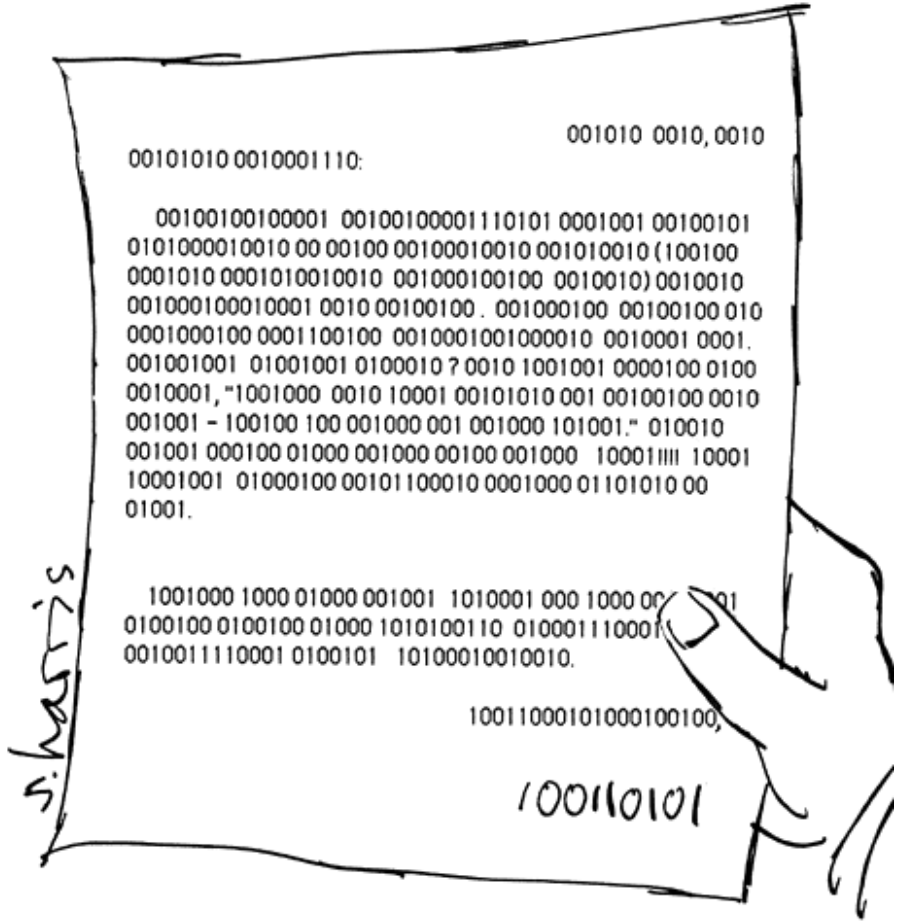


computer with  
a ballot box



computer with  
a heating element





00101010 0010001110: 001010 0010, 0010

00100100100001 00100100001110101 0001001 00100101  
0101000010010 00 00100 00100010010 001010010 (100100  
0001010 0001010010010 001000100100 0010010)0010010  
001000100010001 0010 00100100 . 001000100 00100100 010  
0001000100 0001100100 0010001001000010 0010001 0001.  
001001001 01001001 0100010 ? 0010 1001001 0000100 0100  
0010001, "1001000 0010 10001 00101010 001 00100100 0010  
001001 - 100100 100 001000 001 001000 101001." 010010  
001001 000100 01000 001000 00100 001000 100011111 10001  
10001001 01000100 00101100010 0001000 01101010 00  
01001.

s. harris

1001000 1000 01000 001001 1010001 000 1000 001 001  
0100100 0100100 01000 1010100110 010001110001  
0010011110001 0100101 10100010010010.

10011000101000100100,

100110101

BINARY LETTER FROM GRANDMA

# A Digital World

Data is a sequence of bits [bit = 0 or 1]

- **Text**
- Programs, executables
- Documents, pictures, sounds, movies, ...

**Base64 encoding** Use 6 bits to represent each alphanumeric symbol.

 very weak type of encryption

Binary	Char	Binary	Char	Binary	Char	Binary	Char	Binary	Char	Binary	Char
000000	A	001011	L	010110	W	100001	h	101100	s	110111	3
000001	B	001100	M	010111	X	100010	i	101101	t	111000	4
000010	C	001101	N	011000	Y	100011	j	101110	u	111001	5
000011	D	001110	O	011001	Z	100100	k	101111	v	111010	6
000100	E	001111	P	011010	a	100101	l	110000	w	111011	7
000101	F	010000	Q	011011	b	100110	m	110001	x	111100	8
000110	G	010001	R	011100	c	100111	n	110010	y	111101	9
000111	H	010010	S	011101	d	101000	o	110011	z	111110	+
001000	I	010011	T	011110	e	101001	p	110100	0	111111	/
001001	J	010100	U	011111	f	101010	q	110101	1		
001010	K	010101	V	100000	g	101011	r	110110	2		

# One-Time Pad Encryption

## Encryption

- Convert text message to  $N$  bits

### *Base64 Encoding*

char	dec	binary
A	0	000000
B	1	000001
...	...	...
M	12	001100
...	...	...

S	E	N	D	M	O	N	E	Y	message
010010	000100	001101	000011	001100	001110	001101	000100	011000	base64

# One-Time Pad Encryption

## Encryption

- Convert text message to N bits
- Generate N random bits (one-time pad)

S	E	N	D	M	O	N	E	Y	message
010010	000100	001101	000011	001100	001110	001101	000100	011000	base64
110010	010011	110110	111001	011010	111001	100010	111111	010010	random bits

# One-Time Pad Encryption

## Encryption

- Convert text message to N bits
- Generate N random bits (one-time pad)
- Take bitwise XOR of two bitstrings

↑  
sum corresponding pair of bits: 1 if sum is odd, 0 if even

*XOR Truth Table*

x	y	$x \oplus y$
0	0	0
0	1	1
1	0	1
1	1	0

<b>S</b>	<b>E</b>	<b>N</b>	<b>D</b>	<b>M</b>	<b>O</b>	<b>N</b>	<b>E</b>	<b>Y</b>	message
010010	000100	001101	000011	001100	001110	001101	000100	011000	base64
110010	010011	110110	111001	011010	111001	100010	111111	010010	random bits
100000	010111	111011	111010	010110	110111	101111	111011	001010	XOR

↑  
 $0 \oplus 1 = 1$

# One-Time Pad Encryption

## Base64 Encoding

char	dec	binary
A	0	000000
B	1	000001
...	...	...
w	22	010110
...	...	...

## Encryption

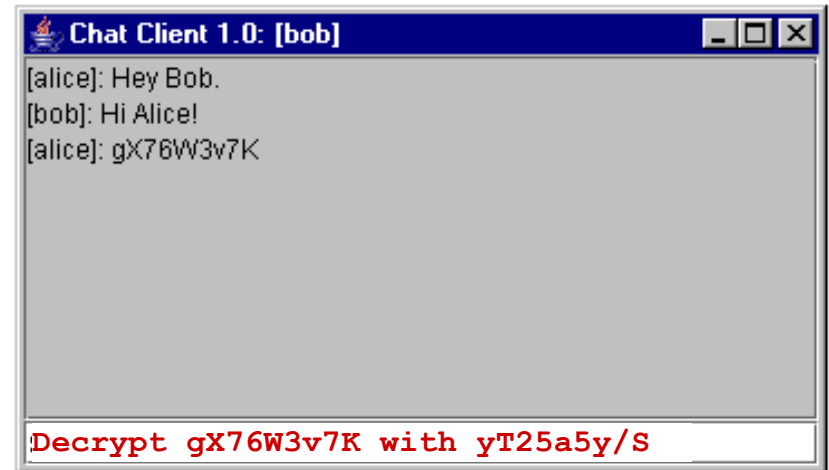
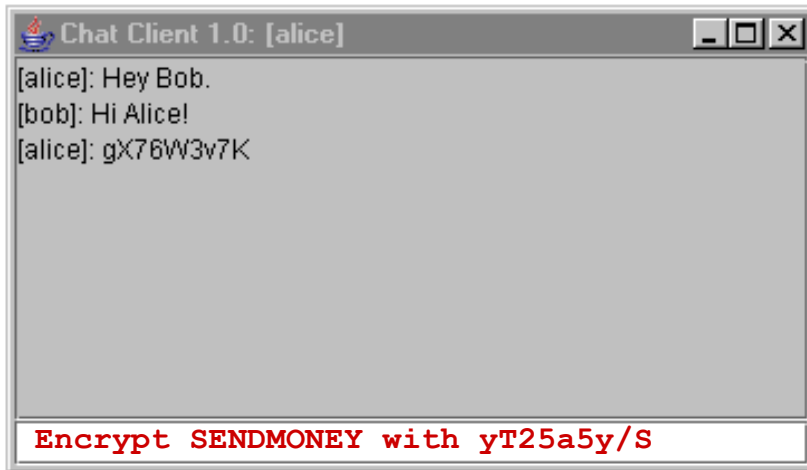
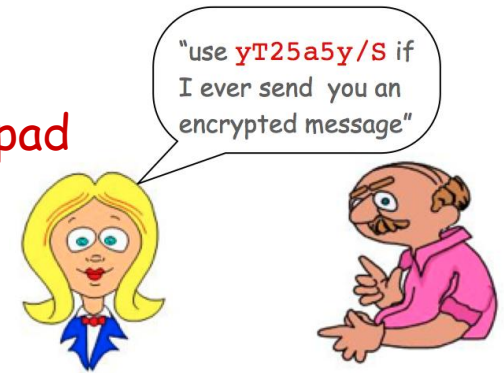
- Convert text message to N bits
- Generate N random bits (one-time pad)
- Take bitwise XOR of two bitstrings
- Convert binary back into text

S	E	N	D	M	O	N	E	Y	message
010010	000100	001101	000011	001100	001110	001101	000100	011000	base64
110010	010011	110110	111001	011010	111001	100010	111111	010010	random bits
100000	010111	111011	111010	010110	110111	101111	111011	001010	XOR
g	x	7	6	w	3	v	7	k	encrypted

# Secure Chat (review)

Alice wants to send a secret message to Bob?

- Sometime in the past, they exchange a **one-time pad**
- Alice uses the pad to encrypt the message
- Bob uses the same pad to decrypt the message



**Key point** Without the pad, Eve cannot understand the message





# One-Time Pad Decryption

## Decryption

- Convert encrypted message to binary

g	x	7	6	w	3	v	7	k
---	---	---	---	---	---	---	---	---

encrypted

# One-Time Pad Decryption

## Decryption

- Convert encrypted message to binary

### *Base64 Encoding*

char	dec	binary
A	0	000000
B	1	000001
...	...	...
W	22	010110
...	...	...

g	x	7	6	W	3	v	7	K	encrypted
100000	010111	111011	111010	010110	110111	101111	111011	001010	base64

# One-Time Pad Decryption

## Decryption

- Convert encrypted message to binary
- Use **same** N random bits (one-time pad)

g	x	7	6	w	3	v	7	k	encrypted
100000	010111	111011	111010	010110	110111	101111	111011	001010	base64
110010	010011	110110	111001	011010	111001	100010	111111	010010	random bits

# One-Time Pad Decryption

## Decryption

- Convert encrypted message to binary
- Use same N random bits (one-time pad)
- Take bitwise XOR of two bitstrings

*XOR Truth Table*

x	y	$x \wedge y$
0	0	0
0	1	1
1	0	1
1	1	0

g	x	7	6	w	3	v	7	K	encrypted
100000	010111	111011	111010	010110	110111	101111	111011	001010	base64
110010	010011	110110	111001	011010	111001	100010	111111	010010	random bits
010010	000100	001101	000011	001100	001110	001101	000100	011000	XOR

↙  
 $1 \wedge 1 = 0$

# One-Time Pad Decryption

## Base64 Encoding

char	dec	binary
A	0	000000
B	1	000001
...	...	...
M	12	001100
...	...	...

## Decryption

- Convert encrypted message to binary
- Use same N random bits (one-time pad)
- Take bitwise XOR of two bitstrings
- Convert back into text

<b>g</b>	<b>x</b>	<b>7</b>	<b>6</b>	<b>w</b>	<b>3</b>	<b>v</b>	<b>7</b>	<b>K</b>	encrypted
100000	010111	111011	111010	010110	110111	101111	111011	001010	base64
110010	010011	110110	111001	011010	111001	100010	111111	010010	random bits
010010	000100	001101	000011	001100	001110	001101	000100	011000	XOR
<b>S</b>	<b>E</b>	<b>N</b>	<b>D</b>	<b>M</b>	<b>O</b>	<b>N</b>	<b>E</b>	<b>Y</b>	message

# Why Does It Work?

Crucial property Decrypted message = original message

Notation	Meaning
$a$	original message bit
$b$	one-time pad bit
$\wedge$	XOR operator
$a \wedge b$	encrypted message bit
$(a \wedge b) \wedge b$	decrypted message bit

*XOR Truth Table*

$x$	$y$	$x \wedge y$
0	0	0
0	1	1
1	0	1
1	1	0

Why is crucial property true?

- Use properties of XOR.
- $(a \wedge b) \wedge b = a \wedge (b \wedge b) = a \wedge 0 = a$ 
  - ↑ associativity of  $\wedge$
  - ↑ always 0
  - ↑ identity

# One-Time Pad Decryption (with the wrong pad)

## Decryption

- Convert encrypted message to binary

g	x	7	6	w	3	v	7	k
---	---	---	---	---	---	---	---	---

 encrypted



# One-Time Pad Decryption (with the wrong pad)

## Decryption

- Convert encrypted message to binary

g	x	7	6	w	3	v	7	K	encrypted
100000	010111	111011	111010	010110	110111	101111	111011	001010	base64

# One-Time Pad Decryption (with the wrong pad)

## Decryption

- Convert encrypted message to binary
- Use **wrong** N bits (bogus one-time pad)

g	x	7	6	w	3	v	7	K	encrypted
100000	010111	111011	111010	010110	110111	101111	111011	001010	base64
101000	011100	110101	101111	010010	111001	100101	101010	001010	wrong bits

# One-Time Pad Decryption (with the wrong pad)

## Decryption

- Convert encrypted message to binary
- Use **wrong** N bits (bogus one-time pad)
- Take bitwise XOR of two bitstrings

g	x	7	6	w	3	v	7	K	encrypted
100000	010111	111011	111010	010110	110111	101111	111011	001010	base64
101000	011100	110101	101111	010010	111001	100101	101010	001010	wrong bits
001000	001011	001110	010101	000100	001110	001010	010001	000000	XOR

# One-Time Pad Decryption (with the wrong pad)

## Decryption

- Convert encrypted message to binary
- Use **wrong** N bits (bogus one-time pad)
- Take bitwise XOR of two bitstrings
- Convert back into text: **Oops**

g	x	7	6	w	3	v	7	k	encrypted
100000	010111	111011	111010	010110	110111	101111	111011	001010	base64
101000	011100	110101	101111	010010	111001	100101	101010	001010	wrong bits
001000	001011	001110	010101	000100	001110	001010	010001	000000	XOR
I	L	O	V	E	O	K	R	A	wrong message



# Goods and Bads of One-Time Pads

## Good

- Easily computed by hand
- Very simple encryption/decryption processes
- Provably unbreakable if bits are truly random [Shannon, 1940s]

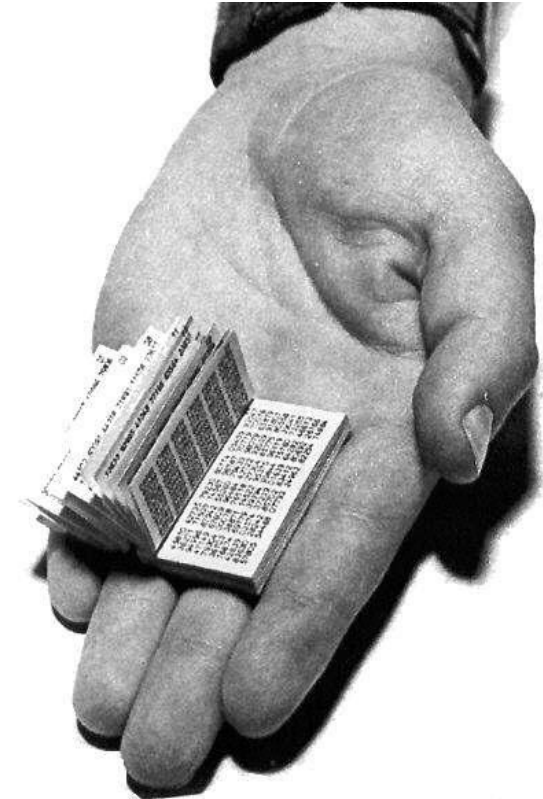
← eavesdropper Eve sees only random bits

"one time" means one time only

## Bad

- Easily breakable if pad is re-used
- Pad must be as long as the message
- Truly random bits are very hard to come by
- **Pad must be distributed securely**

← impractical for Web commerce



a Russian one-time pad

# Pseudo-Random Bit Generator

## Practical middle-ground

- Let's make a "random"-bit generator gadget
- Alice and Bob each get identical small gadgets

← instead of identical large one-time pads

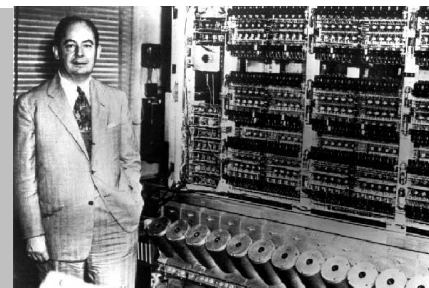
## How to make small gadget that produces "random" bits

- Enigma machine
- **Linear feedback shift register**
- Linear congruential generator
- Blum-Blum-Shub generator
- ...

*“ Anyone who considers arithmetical methods of producing random digits is, of course, in a state of sin. ”*

*– Jon von Neumann (left)*

*– ENIAC (right)*

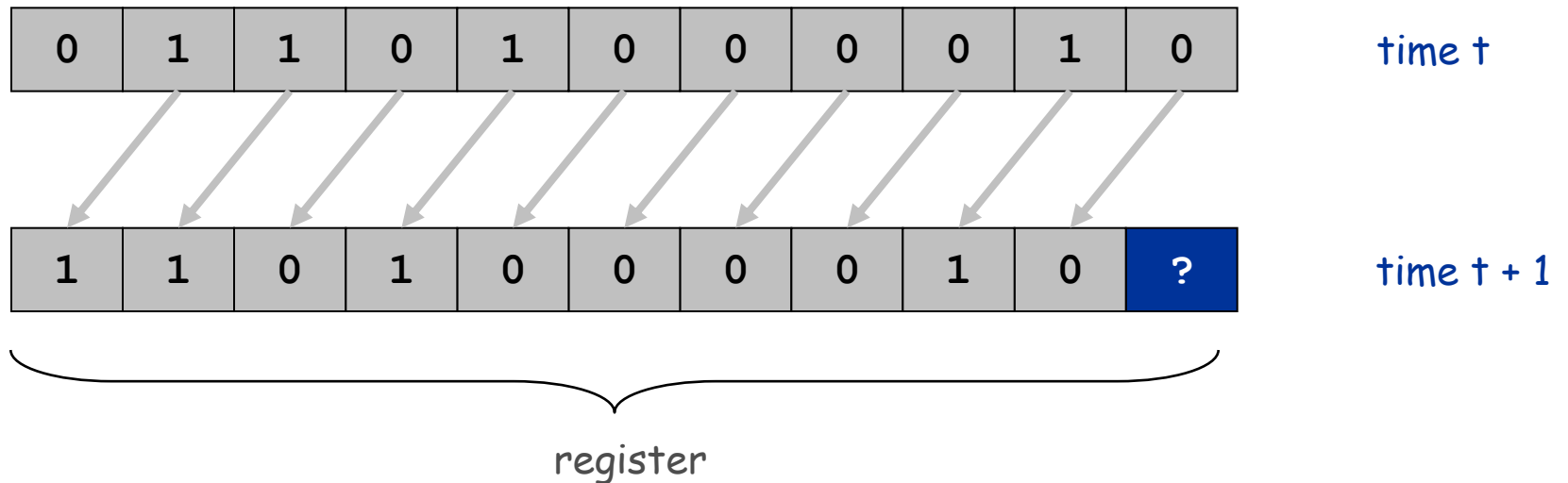




# Shift Register

## Shift register terminology.

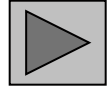
- Bit: 0 or 1
- Cell: storage element that holds one bit
- Register: sequence of cells
- Seed: initial sequence of bits
- Shift register: when clock ticks, bits propagate one position to left



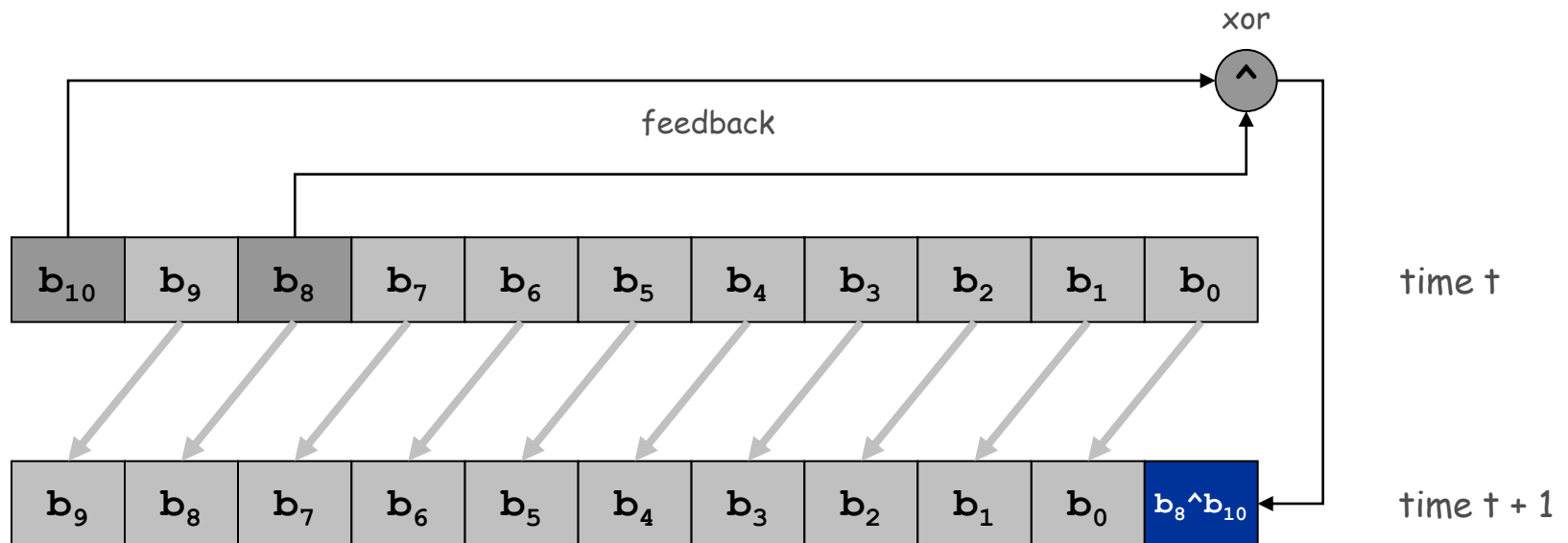
# Linear Feedback Shift Register (LFSR)

## {8, 10} linear feedback shift register

- Shift register with 11 cells
- Bit  $b_0$  is XOR of previous bits  $b_8$  and  $b_{10}$
- Pseudo-random bit =  $b_0$



LFSR demo



# Random Numbers

Q. Are these 2000 numbers random? If not, what is the pattern?

```
110010010011110110111001011010111001100010111111010010000100110100101111001100100111111
101110000010101100010000111010100110100001111001001100111011111110101000001000010001010
010101000110000010111100010010011010110111100011010011011100111101011110010001001110101
011101000001010010001000110101010111000000010110000010011100010111011010010101100110000
111111100110000011111100011000011011110011101001111010011100100111011101110101010101000
000000010000000010100000010001000010101010010000000110100000111001000110111010111010100
010100001010001001000101011010100001100001001111001011100111001011110111001001010111011
000010101110010000101110100100101001101100011110111011001010101111000000100110000101111
100100100011101101011010110001100011101111011010100101100001100111001111110111100001010
011001000111111010110000100011100101011011100001101011001110001111101101100010110111010
0110101001111000011100110011011111111110100000001001000001011010001001100101011111100001
000011001010011111000111000110110110111011011010101101100000110111000111010110110100011
011001011101111001010100111000001110110001101011101110001010101101000000110010000111110
1001100010011111101011100010001011010101001100000011111000011000110011110111111001010000
111000100110110101111011000100101110101100101000111100010110011010011111100111000011110
110011001011111111001000000111010000110100100111001101110111110101010001000000101010000
100000100101000101100010100111010001110100101101001100110011111111111000000000110000000
111100000110011000111111110110000001011100001001011001011001111001111100111100011110011
011001111101111100010100011010001011100101001011100011001011011111001101000111110010110
001110011101101111010110100100011001101011111110001000001101010001110000101101100100110
111101111010010100100110001101111101110100010101001010000011000100011110101011001000001
111010001100100101111101100100010111101010010010000110110100111011001110101111101000100
010010101010111000000001110000001101100001110111001101010111110000010001100010101111010
```

A. No. This is output of {8, 10} LFSR with seed 01101000010!

# LFSR Encryption

## Base64 Encoding

char	dec	binary
A	0	000000
B	1	000001
...	...	...
w	22	010110
...	...	...

## LFSR encryption

- Convert text message to N bits
- Initialize LFSR with small seed
- Generate N bits **with LFSR**
- Take bitwise XOR of two bitstrings
- Convert binary back into text

S	E	N	D	M	O	N	E	Y	message
010010	000100	001101	000011	001100	001110	001101	000100	011000	base64
110010	010011	110110	111001	011010	111001	100010	111111	010010	LFSR bits
100000	010111	111011	111010	010110	110111	101111	111011	001010	XOR
g	x	7	6	w	3	v	7	K	encrypted

# LFSR Decryption

## LFSR Decryption

- Convert encrypted message to N bits
- Initialize identical LFSR with same seed
- Generate N bits **with LFSR**
- Take bitwise XOR of two bitstrings
- Convert binary back into text

### Base64 Encoding

char	dec	binary
A	0	000000
B	1	000001
...	...	...
M	12	001100
...	...	...

g	x	7	6	w	3	v	7	K	encrypted
100000	010111	111011	111010	010110	110111	101111	111011	001010	base64
110010	010011	110110	111001	011010	111001	100010	111111	010010	LFSR bits
010010	000100	001101	000011	001100	001110	001101	000100	011000	XOR
S	E	N	D	M	O	N	E	Y	message

# Goods and Bads of LFSR Encryption

## Goods

- Easily computed with simple machine
- Very simple encryption/decryption process
- Scalable: 20 cells for 1 million bits; 30 cells for 1 billion bits  
[ but need theory of finite groups to know where to put taps ]



a commercially available LFSR

## Bads

- Still need secure, independent way to distribute LFSR seed
- The bits are not truly random  
[ bits in our 11-bit LFSR cycle after  $2^{11} - 1 = 2047$  steps ]
- Experts have cracked LFSR  
[ more complicated machines needed ]

# Other LFSR Applications

## What else can we do with a LFSR?

- DVD encryption with CSS
- DVD decryption with DeCSS!
- Subroutine in military cryptosystems



DVD Jon  
(Norwegian hacker)

```
/*  efdtt.c      Author:  Charles M. Hannum <root@ihack.net>          */
/*  Usage is:   cat title-key scrambled.vob | efdtt >clear.vob      */

#define m(i) (x[i]^s[i+84])<<

        unsigned char x[5]          ,y,s[2048];main(
n){for( read(0,x,5          );read(0,s ,n=2048
        ); write(1          ,s,n          ) if(s
[y=s          [13]%8+20] /16%4 ==1          ){int
i=m(          1)17 ^256 +m(0) 8,k          =m(2)
0,j=          m(4) 17^ m(3) 9^k*          2-k%8
^8,a          =0,c          =26;for (s[y]          -=16;
--c;j          *=2)a=          a*2^i&          1,i=i /2^j&1
<<24;for(j=          127;          ++j<n;c=c>
        y)
        c

        +=y=i^i/8^i>>4^i>>12,
i=i>>8^y<<17,a^=a>>14,y=a^a*8^a<<6,a=a
>>8^y<<9,k=s[j],k          ="7Wo~'G_\216"[k
&7]+2^"cr3sfw6v;*k+>/n."[k>>4]*2^k*257/
        8,s[j]=k^(k&k*2&34)*6^c+~y
        ;}}


```

# A Closing Profound Question

Q. What is a random number?

LFSR does not produce random numbers

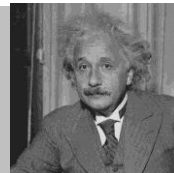
- It is a very simple deterministic machine
- But not obvious how to distinguish the bits it produces from random

Q. Are truly random processes found in nature?

- Motion of cosmic rays or subatomic particles?
- Mutations in DNA?

Q. Or, is the natural world a (not-so-simple) deterministic machine?

*“ God does not play dice. ”*  
– *Albert Einstein*

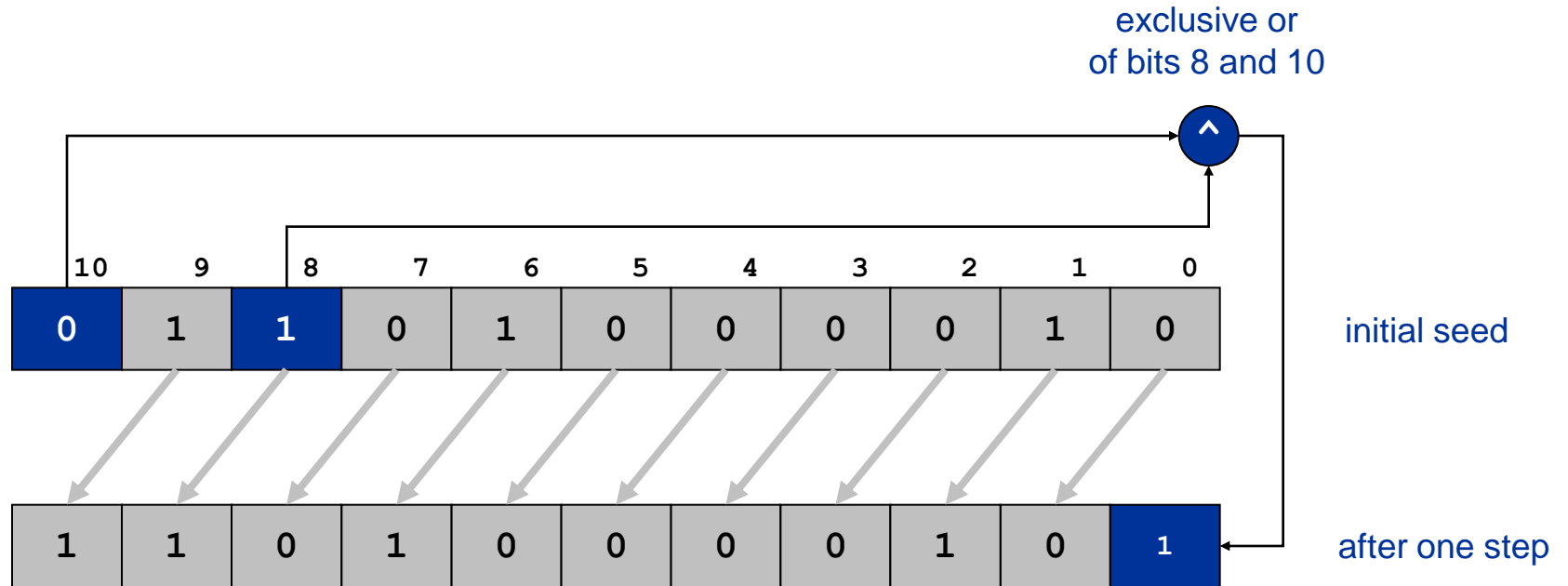




# Extra Slides

---

# Linear Feedback Shift Register



One step of an 11-bit LFSR with initial seed 01101000010 and tap at position 8