

Automata, Computability and Complexity

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Homework 1

January 16, 2020; Due January 23, 2020, beginning of class

Mostly a review of induction.

Problem B1 (50 pts). Let $S: \mathbb{N} \rightarrow \mathbb{N}$ be the function given by

$$S(n) = n + 1, \quad \text{for all } n \in \mathbb{N}.$$

Define the function $add: \mathbb{N} \times \mathbb{N} \rightarrow \mathbb{N}$ recursively as follows for all $m, n \in \mathbb{N}$:

$$add(m, 0) = m \tag{A}$$

$$add(m, S(n)) = S(add(m, n)). \tag{B}$$

(1) Prove that

$$add(add(m, n), p) = add(m, add(n, p))$$

for all $m, n, p \in \mathbb{N}$. In other words, add is associative.

Hint. Use induction on p . It turns out that $add(m, n) = m + n$, where $+$ is the usual addition of natural numbers but **you can't use this fact!**

(2) We would like to prove that

$$add(m, n) = add(n, m), \quad \text{for all } m, n \in \mathbb{N}$$

but this is a little tricky. First prove

(2a)

$$add(0, n) = n, \quad \text{for all } n \in \mathbb{N}.$$

Also prove

(2b)

$$add(S(m), n) = S(add(m, n)), \quad \text{for all } m, n \in \mathbb{N}.$$

Hint. Use induction on n .

Finally, prove that

$$add(m, n) = add(n, m), \quad \text{for all } m, n \in \mathbb{N}.$$

In other words, *add* is commutative.

Hint. Use induction on m .

Problem B2 (10 pts). Let Σ be any alphabet. For any string $w \in \Sigma^*$ recall that w^n is defined inductively as follows:

$$\begin{aligned}w^0 &= \epsilon \\w^{n+1} &= w^n w, \quad n \in \mathbb{N}.\end{aligned}$$

For any string $w \in \Sigma^*$ and any natural numbers $m, n \in \mathbb{N}$, prove that

$$w^m w^n = w^{m+n}.$$

Hint: Use induction on n .

Problem B3 (40 pts). Let Σ be any alphabet. Given a string $w \in \Sigma^*$, its reversal w^R is defined inductively as follows: $\epsilon^R = \epsilon$, and $(ua)^R = au^R$, where $a \in \Sigma$ and $u \in \Sigma^*$.

(1) Prove that $(uv)^R = v^R u^R$, for all $u, v \in \Sigma^*$.

(2) Prove that $(w^R)^R = w$, for all $w \in \Sigma^*$.

TOTAL: 100 points.