

### Markov Chains

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September 15, 2014

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Markov chains. Definition and examples

- Chapman Kolmogorov equations
- Gambler's ruin problem
- Queues in communication networks: Transition probabilities
- Classes of States
- Limiting distributions
- Ergodicity

Queues in communication networks: Limit probabilities

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- Consider time index n = 0, 1, 2, ... & time dependent random state  $X_n$
- State X<sub>n</sub> takes values on a countable number of states
  - ▶ In general denotes states as *i* = 0, 1, 2, ...
  - Might change with problem
- Denote the history of the process  $\mathbf{X}_n = [X_n, X_{n-1}, \dots, X_0]^T$
- Denote stochastic process as  $X_{\mathbb{N}}$
- The stochastic process  $X_{\mathbb{N}}$  is a Markov chain (MC) if

$$P[X_{n+1} = j | X_n = i, \mathbf{X}_{n-1}] = P[X_{n+1} = j | X_n = i] = P_{ij}$$

▶ Future depends only on current state X<sub>n</sub>

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- Process's history  $X_{n-1}$  irrelevant for future evolution of the process
- Probabilities P<sub>ij</sub> are constant for all times (time invariant)
- $\blacktriangleright$  From the definition we have that for arbitrary m

$$\mathsf{P}\left[X_{n+m} \mid X_n, \mathbf{X}_{n-1}\right] = \mathsf{P}\left[X_{n+m} \mid X_n\right]$$

- ►  $X_{n+m}$  depends only on  $X_{n+m-1}$ , which depends only  $onX_{n+m-2}$ , ... which depends only on  $X_n$
- Since  $P_{ij}$ 's are probabilities they're positive and sum up to 1

$$P_{ij} \geq 0$$
  $\sum_{j=1}^{\infty} P_{ij} = 1$ 

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▶ Group transition probabilities *P*<sub>ij</sub> in a "matrix" **P** 

$$\mathbf{P} := \begin{pmatrix} P_{00} & P_{01} & P_{02} & \dots \\ P_{10} & P_{11} & P_{12} & \dots \\ \vdots & \vdots & \vdots & \vdots \\ P_{i0} & P_{i1} & P_{i2} & \dots \\ \vdots & \vdots & \vdots & \ddots \end{pmatrix}$$

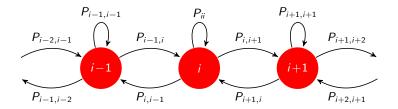
Not really a matrix if number of states is infinite

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A graph representation is also used

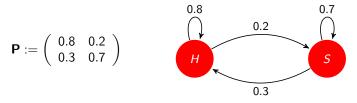


Useful when number of states is infinite

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- I can be happy  $(X_n = 0)$  or sad  $(X_n = 1)$ .
- Happiness tomorrow affected by happiness today only
- Model as Markov chain with transition probabilities



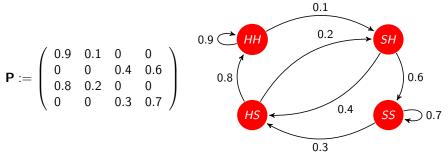
▶ Inertia ⇒ happy or sad today, likely to stay happy or sad tomorrow ( $P_{00} = 0.8, P_{11} = 0.7$ )

• But when sad, a little less likely so  $(P_{00} > P_{11})$ 

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# Example: Happy - Sad, version 2

- Happiness tomorrow affected by today and yesterday
- ► Define double states HH (happy-happy), HS (happy-sad), SH, SS
- Only some transitions are possible
  - ► HH and SH can only become HH or HS
  - HS and SS can only become SH or SS

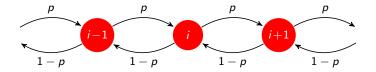


- More time happy or sad increases likelihood of staying happy or sad
- ► State augmentation ⇒ Capture longer time memory



# Random (drunkard's) walk

• Step to the right with probability p, to the left with prob. (1-p)



- States are  $0, \pm 1, \pm 2, \ldots$ , number of states is infinite
- Transition probabilities are

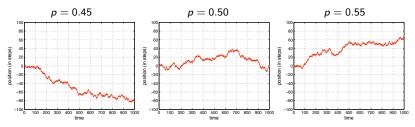
$$P_{i,i+1} = p,$$
  $P_{i,i-1} = 1 - p,$ 



# Random (drunkard's) walk - continued



▶ Random walks behave differently if p < 1/2, p = 1/2 or p > 1/2



- With p > 1/2 diverges to the right (grows unbounded almost surely)
- With p < 1/2 diverges to the left
- With p = 1/2 always come back to visit origin (almost surely)
- Because number of states is infinite we can have all states transient
  - They are not revisited after some time (more later)

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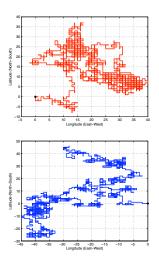
- Take a step in random direction East, West, South or North
  - $\Rightarrow$  E, W, S, N chosen with equal probability
- States are pairs of coordinates (x, y)
  - $x = 0, \pm 1, \pm 2, \dots$  and  $y = 0, \pm 1, \pm 2, \dots$
- Transiton probabilities are not zero only for points adjacent in the grid

$$P[x(t+1) = i+1, y(t+1) = j | x(t) = i, y(t) = j] = \frac{1}{4}$$

$$P[x(t+1) = i-1, y(t+1) = j | x(t) = i, y(t) = j] = \frac{1}{4}$$

$$P[x(t+1) = i, y(t+1) = j+1 | x(t) = i, y(t) = j] = \frac{1}{4}$$

$$P[x(t+1) = i, y(t+1) = j-1 | x(t) = i, y(t) = j] = \frac{1}{4}$$





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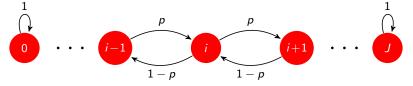


- Some random facts of life for equiprobable random walks
- $\blacktriangleright$  In one and two dimensions probability of returning to origin is 1
- Will almost surely return home
- In more than two dimensions, probability of returning to origin is less than 1
- In three dimensions probability of returning to origin is 0.34
- ▶ Then 0.19, 0.14, 0.10, 0.08, ...

# Random walk with borders (gambling)



- As a random walk, but stop moving when i = 0 or i = J
  - Models a gambler that stops playing when ruined,  $X_n = 0$
  - Or when reaches target gains  $X_n = J$



▶ States are 0, 1, ..., J. Finite number of states (J). Transition probs.

$$P_{i,i+1} = p, \quad P_{i,i-1} = 1 - p, \qquad P_{00} = 1, \quad P_{JJ} = 1$$

•  $P_{ij} = 0$  for all other transitions

- ▶ States 0 and J are called absorbing. Once there stay there forever
- The rest are transient states. Visits stop almost surely

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Queues in communication networks: Limit probabilities



- What can be said about multiple transitions ?
- Transition probabilities between two time slots

$$P_{ij}^2 := \mathsf{P}\left[X_{m+2} = j \mid X_m = i\right]$$

▶ Probabilities of  $X_{m+n}$  given  $X_n \Rightarrow n$ -step transition probabilities

$$P_{ij}^{n} := \mathsf{P}\left[X_{m+n} = j \mid X_m = i\right]$$

- ▶ Relation between *n*-step, *m*-step and (m + n)-step transition probs.
  - Write P<sup>m+n</sup><sub>ij</sub> in terms of P<sup>m</sup><sub>ij</sub> and P<sup>n</sup><sub>ij</sub>
- All questions answered by Chapman-Kolmogorov's equations

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## 2-step transition probabilities



Start considering transition probs. between two time slots

$$P_{ij}^2 = \mathsf{P}\left[X_{n+2} = j \mid X_n = i\right]$$

Using the theorem of total probability

$$P_{ij}^{2} = \sum_{k=1}^{\infty} P\left[X_{n+2} = j \mid X_{n+1} = k, X_{n} = i\right] P\left[X_{n+1} = k \mid X_{n} = i\right]$$

▶ In the first probability, conditioning on  $X_n = i$  is unnecessary. Thus

$$P_{ij}^{2} = \sum_{k=1}^{\infty} P\left[X_{n+2} = j \mid X_{n+1} = k\right] P\left[X_{n+1} = k \mid X_{n} = i\right]$$

Which by definition yields

$$P_{ij}^2 = \sum_{k=1}^{\infty} P_{kj} P_{ik}$$



 Identical argument can be made (condition on X<sub>0</sub> to simplify notation, possible because of time invariance)

$$P_{ij}^{m+n} = \mathsf{P}\left[X_{n+m} = j \mid X_0 = i\right]$$

Use theorem of total probability, remove unnecessary conditioning and use definitions of *n*-step and *m*-step transition probabilities

$$P_{ij}^{m+n} = \sum_{k=1}^{\infty} P[X_{m+n} = j | X_m = k, X_0 = i] P[X_m = k | X_0 = i]$$

$$P_{ij}^{m+n} = \sum_{k=1}^{\infty} P[X_{m+n} = j | X_m = k] P[X_m = k | X_0 = i]$$

$$P_{ij}^{m+n} = \sum_{k=1}^{\infty} P_{kj}^n P_{ik}^m$$



Chapman Kolmogorov is intuitive. Recall

$$P_{ij}^{m+n} = \sum_{k=1}^{\infty} P_{kj}^n P_{ik}^m$$

- Between times 0 and m + n time m occurred
- At time m, the chain is in some state X<sub>m</sub> = k
  ⇒ P<sup>m</sup><sub>ik</sub> is the probability of going from X<sub>0</sub> = i to X<sub>m</sub> = k
  ⇒ P<sup>n</sup><sub>kj</sub> is the probability of going from X<sub>m</sub> = k to X<sub>m+n</sub> = j
  ⇒ Product P<sup>m</sup><sub>ik</sub>P<sup>n</sup><sub>kj</sub> is then the probability of going from X<sub>0</sub> = i to X<sub>m+n</sub> = j passing through X<sub>m</sub> = k at time m
  Since any k might have occurred sum over all k

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- ▶ Define matrices  $\mathbf{P}^{(m)}$  with elements  $P_{ij}^m$ ,  $\mathbf{P}^{(n)}$  with elements  $P_{ij}^n$  and  $\mathbf{P}^{(m+n)}$  with elements  $P_{ij}^{m+n}$
- $\sum_{k=1}^{\infty} P_{kj}^n P_{ik}^m$  is the (i,j)-th element of matrix product  $\mathbf{P}^{(m)} \mathbf{P}^{(n)}$

Chapman Kolmogorov in matrix form

$$\mathbf{P}^{(m+n)} = \mathbf{P}^{(m)}\mathbf{P}^{(n)}$$

• Matrix of (n + m)-step transitions is product of *n*-step and *m*-step

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• For m = n = 1 (2-step transition probabilities) matrix form is

 $\mathbf{P}^{(2)} = \mathbf{P}\mathbf{P} = \mathbf{P}^2$ 

Proceed recursively backwards from n

$$\mathbf{P}^{(n)} = \mathbf{P}^{(n-1)}\mathbf{P} = \mathbf{P}^{(n-2)}\mathbf{P}\mathbf{P} = \ldots = \mathbf{P}^n$$

Have proved the following

#### Theorem

The matrix of n-step transition probabilities  $P^{(n)}$  is given by the n-th power of the transition probability matrix **P**. *i.e.*,

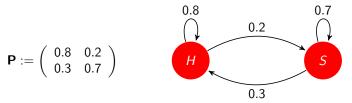
 $\mathbf{P}^{(n)}=\mathbf{P}^n$ 

Henceforth we write  $\mathbf{P}^n$ 

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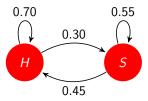


Happiness transitions in one day (not the same as earlier example)



> Transition probabilities between today and the day after tomorrow?

$$\mathbf{P}^2 := \left( \begin{array}{cc} 0.70 & 0.30 \\ 0.45 & 0.55 \end{array} \right)$$



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... After a week and after a month

$$\mathbf{P}^7 := \left( \begin{array}{ccc} 0.6031 & 0.3969 \\ 0.5953 & 0.4047 \end{array} \right) \qquad \qquad \mathbf{P}^{30} := \left( \begin{array}{ccc} 0.6000 & 0.4000 \\ 0.6000 & 0.4000 \end{array} \right)$$

▶ Matrices  $\mathbf{P}^7$  and  $\mathbf{P}^{30}$  almost identical  $\Rightarrow \lim_{n\to\infty} \mathbf{P}^n$  exists

Note that this is a regular limit

- After a month transition from H to H with prob. 0.6 and from S to H also 0.6
- State becomes independent of initial condition
- ▶ Rationale: 1-step memory  $\Rightarrow$  initial condition eventually forgotten

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## Unconditional probabilities



- ► All probabilities so far are conditional, i.e.,  $P[X_n = j | X_0 = i]$
- Want unconditional probabilities  $p_j(n) := P[X_n = j]$
- Requires specification of initial conditions p<sub>i</sub>(0) := P [X<sub>0</sub> = i]
- Using theorem of total probability and definitions of  $P_{ij}^n$  and  $p_j(n)$

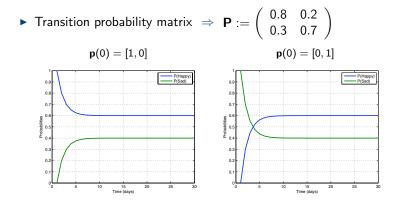
$$p_j(n) := P[X_n = j] = \sum_{i=1}^{\infty} P[X_n = j | X_0 = i] P[X_0 = i]$$
$$= \sum_{i=1}^{\infty} P_{ij}^n p_i(0)$$

• Or in matrix form (define vector  $\mathbf{p}(n) := [p_1(n), p_2(n), \ldots]^T)$ 

$$\mathbf{p}(n) = \mathbf{P}^{n\,T}\mathbf{p}(0)$$

# Example: Happy-Sad





For large *n* probabilities  $\mathbf{p}(t)$  are independent of initial state  $\mathbf{p}(0)$ 

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Queues in communication networks: Limit probabilities

## Gambler's ruin problem



- You place \$1 bets,
  - (a) With probability p you gain \$1 and
  - (b) With probability q = (1 p) you loose your \$1 bet
- Start with an initial wealth of \$i<sub>0</sub>
- Define bias factor  $\alpha := q/p$ 
  - If  $\alpha > 1$  more likely to loose than win (biased against gambler)
  - lpha < 1 favors gambler (more likely to win than loose )
  - $\alpha = 1/2$  game is fair
- You keep playing until
  - (a) You go broke (loose all your money)
  - (b) You reach a wealth of N
- Prob. S<sub>i</sub> of reaching \$N before going broke for initial wealth \$i?
  - S stands for success

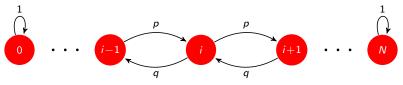
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• Model as Markov chain  $X_{\mathbb{N}}$ . Transition probabilities

$$P_{i,i+1} = p, \quad P_{i,i-1} = q, \quad P_{00} = P_{NN} = 1$$

• Realizations  $x_{\mathbb{N}}$ . Initial state = initial wealth =  $i_0$ 



▶ Sates 0 and N said absorbing. Eventually end up in one of them

- Remaining states said transient (visits eventually stop)
- Being absorbing states says something about the limit wealth

$$\lim_{n\to\infty} x_n = 0, \text{ or } \lim_{n\to\infty} x_n = N, \quad \Rightarrow \quad S_i := \mathsf{P}\left(\lim_{n\to\infty} X_n = N \mid X_0 = i\right)$$



- Prob.  $S_i$  of successful betting run depends on current state *i* only
- We can relate probabilities of SBR from adjacent states

$$S_i = S_{i+1}P_{i,i+1} + S_{i-1}P_{i,i-1} = S_{i+1}p + S_{i-1}q$$

• Recall p + q = 1. Reorder terms

$$p(S_{i+1}-S_i) = q(S_i-S_{i-1})$$

• Recall definition of bias  $\alpha = q/p$ 

$$S_{i+1}-S_i=\alpha(S_i-S_{i-1})$$

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### Recursive relations (continued)



• If current state is 0 then  $S_i = S_0 = 0$ . Can write

$$S_2 - S_1 = \alpha(S_1 - S_0) = \alpha S_1$$

• Substitute this in the expression for  $S_3 - S_2$ 

$$S_3 - S_2 = \alpha(S_2 - S_1) = \alpha^2 S_1$$

• Apply recursively backwards from  $S_i - S_{i-1}$ 

$$S_i - S_{i-1} = \alpha(S_{i-1} - S_{i-2}) = \ldots = \alpha^{i-1}S_1$$

Sum up all of the former to obtain

$$S_i - S_1 = S_1 \left( \alpha + \alpha^2 + \ldots + \alpha^{i-1} \right)$$

The latter can be written as a geometric series

$$S_i = S_1 \left( 1 + \alpha + \alpha^2 + \ldots + \alpha^{i-1} \right)$$



• Geometric series can be summed. Assuming  $\alpha \neq 1$ 

$$S_i = \frac{1 - \alpha^i}{1 - \alpha} S_1$$

• Write for i = 1. When in state N,  $S_N = 1$ 

$$1 = S_N = \frac{1 - \alpha^N}{1 - \alpha} S_1$$

Compute S<sub>1</sub> from latter and substitute into expression for S<sub>i</sub>

$$S_i = \frac{1 - \alpha^i}{1 - \alpha^N}$$

► For 
$$\alpha = 1 \implies S_i = iS_1$$
,  $1 = S_N = NS_1$ ,  $\implies S_i = (i/N)$ 

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► Consider exit bound *N* arbitrarily large.

• For 
$$\alpha \geq 1$$
,  $S_i \approx (\alpha^i - 1)/\alpha^N \to 0$ 

- If win prob. does not exceed loose probability will almost surely loose all money
- For  $\alpha < 1$ ,  $P_i \rightarrow 1 \alpha^i$
- If win prob. exceeds loose probability might win
- If initial wealth i sufficiently high, will most likely win
  - $\Rightarrow$  Which explains what we saw on first lecture and homework



Markov chains. Definition and examples

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Queues in communication networks: Limit probabilities



- Communication systems goal
  - $\Rightarrow$  Move packets from generating sources to intended destinations
- Between arrival and departure we hold packets in a memory buffer
- Want to design buffers appropriately

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- Time slotted in intervals of duration  $\Delta t$
- *n*-th slot between times  $n\Delta t$  and  $(n+1)\Delta t$
- Average arrival rate is  $\bar{\lambda}$  packets per unit time
- During slot of duration  $\Delta t$  probability of packet arrival is  $\lambda = \overline{\lambda} \Delta t$
- Packets are transmitted (depart) at a rate of  $\bar{\mu}$  packets per unit time
- During interval  $\Delta t$  probability of packet departure is  $\mu = \bar{\mu} \Delta t$
- Assume no simultaneous arrival and departure (no concurrence)
  - Reasonable for small  $\Delta t$  ( $\mu$  and  $\lambda$  are likely small)

## Queue evolution equations



- $q_n$  denotes number of packets in queue in *n*-th time slot
- ▶  $A_n = nr$ . of packet arrivals,  $\mathbb{D}_n = nr$ . of departures (during *n*-th slot)
- If there are no packets in queue  $q_n = 0$  then there are no departures
- Queue length at time n + 1 can be written as

$$q_{n+1} = q_n + \mathbb{A}_n, \quad \text{if } q_n = 0$$

• If  $q_n > 0$ , departures and arrivals may happen

$$q_{n+1} = \left[q_n + \mathbb{A}_n - \mathbb{D}_n\right]^+, \quad \text{if } q_n > 0$$

•  $\mathbb{A}_n \in \{0,1\}$ ,  $\mathbb{D}_n \in \{0,1\}$  and either  $\mathbb{A}_n = 1$  or  $\mathbb{D}_n = 1$  but not both • Arrival and departure probabilities are

$$\mathsf{P}\left[\mathbb{A}_n=1\right]=\lambda,\qquad\mathsf{P}\left[\mathbb{D}_n=1\right]=\mu$$

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- Future queue lengths depend on current length only
- Probability of queue length increasing

$$\mathsf{P}\left[q_{n+1}=i+1 \mid q_n=i\right] = \mathsf{P}\left[\mathbb{A}_n=1\right] = \lambda, \quad \text{for all } i$$

• Queue length might decrease only if  $q_n > 0$ . Probability is

$$\mathsf{P}\left[q_{n+1}=i-1 \mid q_n=i\right] = \mathsf{P}\left[\mathbb{D}_n=1\right] = \mu, \qquad \text{for all } i > 0$$

Queue length stays the same if it neither increases nor decreases

$$\mathsf{P}\left[q_{n+1}=i \mid q_n=i\right] = 1 - \lambda - \mu, \quad \text{for all } i > 0 \\ \mathsf{P}\left[q_{n+1}=0 \mid q_n=0\right] = 1 - \lambda$$

• No departures when  $q_n = 0$  explain second equation

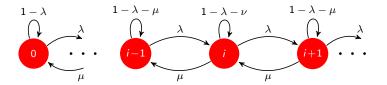
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- MC with states  $0, 1, 2, \ldots$  Identify states with queue lengths
- Transition probabilities for  $i \neq 0$  are

$$P_{i,i-1} = \lambda, \qquad P_{i,i} = 1 - \lambda - \mu, \qquad P_{i,i+1} = \mu$$

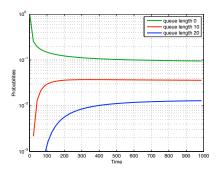
• For 
$$i = 0$$
  $P_{0,0} = 1 - \lambda$  and  $P_{01} = \lambda$ 



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- Build matrix **P** truncating at maximum queue length L = 100
- Arrival rate  $\lambda = 0.3$ . Departure rate  $\mu = 0.33$
- ▶ Initial probability distribution  $\mathbf{p}(0) = [1, 0, 0, ...]^T$  (queue empty)



- Propagate probabilities with product P<sup>n</sup>p(0)
- Probabilities obtained are

$$\mathsf{P}\left[q_n=i\,\big|\,q_0=0\right]=p_i(n)$$

- A few i's (0, 10, 20) shown
- Probability of empty queue  $\approx 0.1$ .
- Occupancy decrease with index



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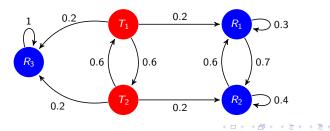
Queues in communication networks: Limit probabilities

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#### Transient and recurrent states



- States of a MC can be recurrent or transient
- Transient states might be visited at the beginning but eventually visits stop
- Almost surely,  $X_n \neq i$  for *n* sufficiently large (qualifications needed)
- Visits to recurrent states keep happening forever
- ► Fix arbitrary *m*
- ▶ Almost surely,  $X_n = i$  for some  $n \ge m$  (qualifications needed)





• Let  $f_i$  be the probability that starting at i, MC ever reenters state i

$$f_i := \mathsf{P}\left[\bigcup_{n=1}^{\infty} X_n = i \, \big| \, X_0 = i\right] = \mathsf{P}\left[\bigcup_{n=m+1}^{\infty} X_n = i \, \big| \, X_m = i\right]$$

• State *i* is recurrent if  $f_i = 1$ 

- ▶ Process reenters *i* again and again (almost surely). Infinitely often
- State *i* is transient if  $f_i < 1$
- Positive probability  $(1 f_i)$  of never coming back to *i*

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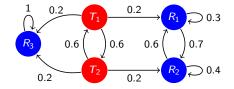
- State  $R_3$  is recurrent because  $P[X_1 = R_3 | X_0 = R_3] = 1$
- ► State  $R_1$  is recurrent because  $P \begin{bmatrix} X_1 = R_1 \mid X_0 = R_1 \end{bmatrix} = 0.3$   $P \begin{bmatrix} X_2 = R_1, X_1 \neq R_1 \mid X_0 = R_1 \end{bmatrix} = (0.7)(0.6)$   $P \begin{bmatrix} X_3 = R_1, X_2 \neq R_1, X_1 \neq R_1 \mid X_0 = R_1 \end{bmatrix} = (0.7)(0.4)(0.6)$   $\vdots$   $P \begin{bmatrix} X_n = R_1, X_{n-1} \neq R_1, \dots, X_1 \neq R_1 \mid X_0 = R_1 \end{bmatrix} = (0.7)(0.4)^{n-1}(0.6)$

• Sum up: 
$$f_i = \sum_{n=1}^{\infty} P\left[X_n = R_1, X_{n-1} \neq R_1, \dots, X_1 \neq R_1 \mid X_0 = R_1\right]$$
  
= 0.3 + 0.7  $\left(\sum_{n=1}^{\infty} 0.4^{n-1}\right)$  0.6 = 0.3 + 0.7  $\left(\frac{1}{1-0.4}\right)$  0.6 = 1

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- States  $T_1$  and  $T_2$  are transient
- Probability of returning to  $T_1$  is  $f_{T_1} = (0.6)^2 = 0.36$
- Might come back to  $T_1$  only if it goes to  $T_2$  (with prob. 0.6)
- Will come back only if it moves back from  $T_2$  to  $T_1$  (with prob. 0.6)

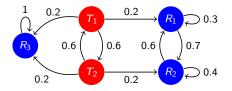


• Likewise, 
$$f_{T_2} = (0.6)^2 = 0.36$$

# Accessibility



- State *j* is accessible from state *i* if  $P_{ii}^n > 0$  for some  $n \ge 0$
- It is possible to enter j if MC initialized at  $X_0 = i$
- ► Since  $P_{ii}^0 = P[X_0 = 1 | X_0 = i] = 1$ , state *i* is accessible from itself



- All states accessible from  $T_1$  and  $T_2$
- Only  $R_1$  and  $R_2$  accessible from  $R_1$  or  $R_2$
- None other than itself accessible from  $R_3$

#### Communication



- States *i* and *j* are said to communicate (*i* ↔ *j*) if
   ⇒ *i* is accessible from *j*, P<sup>n</sup><sub>ij</sub> > 0 for some *n*; and
   ⇒ *j* is accessible from *i*, P<sup>m</sup><sub>ii</sub> > 0 for some *m*
- Communication is an equivalence relation
- Reflexivity:  $i \leftrightarrow i$ 
  - true because  $P_{ii}^0 = 1$
- **Symmetry**: If  $i \leftrightarrow j$  then  $j \leftrightarrow i$ 
  - If  $i \leftrightarrow j$  then  $P_{ij}^n > 0$  and  $P_{ji}^m > 0$  from where  $j \leftrightarrow i$
- Transitivity: If  $i \leftrightarrow j$  and  $j \leftrightarrow k$ , then  $i \leftrightarrow k$ 
  - Just notice that  $P_{ik}^{n+m} \ge P_{ij}^n P_{jk}^m > 0$
- Partitions set of states into disjoint classes (as all equivalences do)
- What are these classes? (start with a brief detour)

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• Define  $N_i$  as the number of visits to state *i* given that  $X_0 = i$ 

$$N_i := \sum_{n=1}^{\infty} \mathbb{I}\left\{X_n = i\right\}$$

- If  $X_n = i$ , this is the last visit to *i* with probability  $1 f_i$
- Prob. revisiting state *i* exactly *n* times is (*n* visits  $\times$  no more visits)

$$\mathsf{P}\left[N_i=n\right]=f_i^n(1-f_i)$$

- Number of visits  $N_i$  has a geometric distribution with parameter  $f_i$
- Expected number of visits is

$$\mathbb{E}[N_i] = \sum_{n=1}^{\infty} nf_i^n(1-f_i) = \frac{1}{1-f_i}$$

▶ For recurrent states  $N_i = \infty$  almost surely and  $\mathbb{E}[N_i] = \infty$   $(f_i = 1)$ 

• Another way of writing  $\mathbb{E}[N_i]$ 

$$\mathbb{E}[N_i] = \sum_{n=1}^{\infty} \mathbb{E}\Big[\mathbb{I}\{X_n = i\}\Big] = \sum_{n=1}^{\infty} P_{ii}^n$$

- ► Recall that: for transient states  $\mathbb{E}[N_i] = 1/(1 f_1)$ for recurrent states  $\mathbb{E}[N_i] = \infty$
- Therefore proving

Theorem

- State i is transient if and only if  $\sum_{n=1}^{\infty} P_{ii}^n < \infty$
- State *i* is recurrent if and only if  $\sum_{n=1}^{\infty} P_{ii}^n = \infty$
- Number of future visits to transient states is finite
- If number of states is finite some states have to be recurrent



#### Theorem

If state i is recurrent and  $i \leftrightarrow j$ , then j is recurrent

Proof.

- If  $i \leftrightarrow j$  then there are I, m such that  $P'_{ji} > 0$  and  $P^m_{ij} > 0$
- ▶ Then, for any *n* we have

$$P_{jj}^{l+n+m} \ge P_{ji}^l P_{ii}^n P_{ij}^m$$

▶ Sum for all *n*. Note that since *i* is recurrent  $\sum_{n=1}^{\infty} P_{ii}^n = \infty$ 

$$\sum_{n=1}^{\infty} P_{jj}^{l+n+m} \geq \sum_{n=1}^{\infty} P_{ji}^{l} P_{ii}^{n} P_{ij}^{m} = P_{ji}^{l} \left( \sum_{n=1}^{\infty} P_{ii}^{n} \right) P_{ij}^{m} = \infty$$

Which implies j is recurrent

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#### Corollary

If state i is transient and  $i \leftrightarrow j$  then j is transient

Proof.

- If j were recurrent, then i would be recurrent from previous theorem
- Since communication defines classes and recurrence is shared by elements of this class, we say that recurrence is a class property
- Likewise, transience is also a class property
- States of a MC are separated in classes of transient and recurrent states



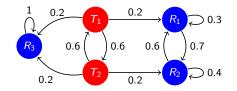
A MC is called irreducible if it has only one class

- All states communicate with each other
- If MC also has finite number of states the single class is recurrent
- If MC infinite, class might be transient
- When it has multiple classes (not irreducible)
  - Classes of transient states  $T_1, T_2, \ldots$
  - Classes of recurrent states  $\mathcal{R}_1, \mathcal{R}_2, \dots$
  - If MC initialized in a recurrent class  $\mathcal{R}_k$ , stays within the class
  - ► If starts in transient class T<sub>k</sub>, might stay on T<sub>k</sub> or end up in a recurrent class R<sub>l</sub>
- ► For large time index *n*, MC restricted to one class
- Can be separated into irreducible components

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Example





- Three classes
  - $\Rightarrow \mathcal{T} := \{T_1, T_2\}$ , class with transient states
  - $\Rightarrow \mathcal{R}_1 := \{R_1, R_2\}$ , class with recurrent states
  - $\Rightarrow \mathcal{R}_2 := \{R_3\}$ , class with recurrent states
- Asymptotically in n suffices to study behavior for the irreducible components R<sub>1</sub> and R<sub>2</sub>

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- States of a MC can be transient of recurrent
- A MC can be partitioned in classes of states reachable from each other
- Elements of the class are either all recurrent or all transient
- A MC with only one class is irreducible
- If not irreducible can be separated in irreducible components

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Queues in communication networks: Limit probabilities

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#### Limiting distributions

- MCs have one-step memory. Eventually they forget initial state
- ▶ What can we say about probabilities for large *n*?

$$\pi_j := \lim_{n \to \infty} \mathsf{P}\left[X_n = j \,\middle|\, X_0 = i\right] = \lim_{n \to \infty} \mathcal{P}_{ij}^n$$

- Implicitly assumed that limit is independent of initial state  $X_0 = i$
- ▶ We've seen that this problem is related to the matrix power **P**<sup>n</sup>

$$\mathbf{P} := \begin{pmatrix} 0.8 & 0.2 \\ 0.3 & 0.7 \end{pmatrix} \qquad \mathbf{P}^7 := \begin{pmatrix} 0.6031 & 0.3969 \\ 0.5953 & 0.4047 \end{pmatrix}$$
$$\mathbf{P}^2 := \begin{pmatrix} 0.7 & 0.3 \\ 0.45 & 0.55 \end{pmatrix} \qquad \mathbf{P}^{30} := \begin{pmatrix} 0.6000 & 0.4000 \\ 0.6000 & 0.4000 \end{pmatrix}$$

- Matrix product converges  $\Rightarrow$  probs. independent of time (large *n*)
- All columns are equal  $\Rightarrow$  probs. independent of initial condition

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#### Periodicity



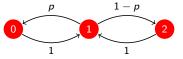
• The period of a state *i* is defined as (d is set of multiples of d)

$$d = \max\left\{d: P_{ii}^n = 0 \text{ for all } n \notin \dot{d}\right\}$$

State i is periodic with period d if and only if

$$\Rightarrow P_{ii}^n 
eq 0$$
 only if *n* is a multiple of *d*  $(n \in \dot{d})$ 

- $\Rightarrow$  *d* is the largest number with this property
- Positive probability of returning to *i* only every *d* time steps
- If period d = 1 state is aperiodic (most often the case)
- Periodicity is a class property



- State 1 has period 2. So do 0 and 2 (class property)
- One dimensional random walk also has period 2

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- ▶ Recall: state *i* is recurrent if chain returns to *i* with probability 1
- Proved it was equivalent to  $\sum_{n=1}^{\infty} P_{ii}^n = \infty$
- Positive recurrent when expected value of return time is finite

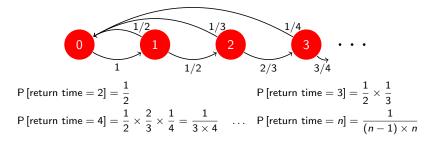
$$\mathbb{E}\left[\text{return time}\right] = \sum_{n=1}^{\infty} n P_{ii}^n \prod_{m=0}^{n-1} (1 - P_{ii}^m) < \infty$$

- Null recurrent if recurrent but  $\mathbb{E}[return time] = \infty$
- Positive and null recurrence are class properties
- Recurrent states in a finite-state MC are positive recurrent
- ► Ergodic states are those that are positive recurrent and aperiodic
- ► An irreducible MC with ergodic states is said to be an ergodic MC

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## Example of a null recurrent MC





It is recurrent because probability of returning is 1 (use induction)

$$\sum_{m=2}^{n} \mathsf{P}\left[\text{return time} = m\right] = \sum_{m=2}^{n} \frac{1}{(m-1) \times m} = \frac{n-1}{n} \to 1$$

Null recurrent because expected return time is infinite

$$\sum_{n=2}^{\infty} n \mathbb{P}\left[\text{return time} = n\right] = \sum_{n=2}^{\infty} \frac{n}{(n-1) \times n} = \sum_{n=2}^{\infty} \frac{1}{(n-1)} = \infty$$

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#### Theorem

For an irreducible ergodic MC,  $\lim_{n\to\infty}P_{ij}$  exists and is independent of the initial state i. That is

$$\pi_j = \lim_{n \to \infty} P_{ij}^n$$
 exists

Furthermore, steady state probabilities  $\pi_j \ge 0$  are the unique nonnegative solution of the system of linear equations

$$\pi_j = \sum_{i=0}^\infty \pi_i P_{ij}, \qquad \sum_{j=0}^\infty \pi_j = 1$$

- As observed, limit probs. independent of initial condition exist
- Simple algebraic equations can be solved to find  $\pi_j$
- ▶ No periodic states, transient states, multiple classes or null recurrent



- Difficult part of theorem is to prove that  $\pi_j = \lim_{n \to \infty} P_{ii}^n$  exists
- ► To see that algebraic relation is true use theorem of total probability (omit conditioning on X<sub>0</sub> to simplify notation)

$$P[X_{n+1} = j] = \sum_{i=1}^{\infty} P[X_{n+1} = j | X_n = i] P[X_n = i]$$
$$= \sum_{i=1}^{\infty} P_{ij} P[X_n = i]$$

▶ If limits exists,  $P[X_{n+1} = j] \approx P[X_n = j] \approx \pi_j$  (sufficiently large *n*)

$$\pi_j = \sum_{i=0}^{\infty} \pi_i P_{ij}$$

• The other equation is true because the  $\pi_j$  are probabilities

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# Vector/matrix notation: Matrix limit



- More compact and illuminating on vector/matrix notation
- Finite MC with J states
- First part of theorem says that  $\lim_{n\to\infty} \mathbf{P}^n$  exists and

$$\lim_{n \to \infty} \mathbf{P}^n = \begin{pmatrix} \pi_1 & \pi_2 & \dots & \pi_J \\ \pi_1 & \pi_2 & \dots & \pi_J \\ \vdots & \vdots & \vdots & \vdots \\ \pi_1 & \pi_2 & \dots & \pi_J \end{pmatrix}$$

▶ Same probs. for all rows  $\Rightarrow$  independent of initial state

Probability distribution for large n.

$$\lim_{n\to\infty}\mathbf{p}(n) = \lim_{n\to\infty}\mathbf{P}^{T^n}\mathbf{p}(0) = [\pi_0, \pi_1, \dots, \pi_J]^T$$

Independent of initial condition

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- Define vector stationary distribution  $\boldsymbol{\pi} := [\pi_0, \pi_1, \dots, \pi_J]^T$
- Limit distribution is unique solution of  $(\mathbf{1} = [1, 1, \ldots]^T)$

$$\boldsymbol{\pi} = \mathbf{P}^T \boldsymbol{\pi}, \qquad \boldsymbol{\pi}^T \mathbf{1} = 1$$

- $\pi$  eigenvector associated with eigenvalue 1 of  $\mathbf{P}^{T}$ 
  - Eigenvectors are defined up to a constant
  - Normalize to sum 1
- ▶ All other eigenvectors of  $\mathbf{P}^{T}$  have modulus smaller than 1
  - ▶ If not,  $\mathbf{P}^n$  diverges, but we know  $\mathbf{P}^n$  contains *n*-step transition probs.
  - $\pi$  eigenvector associated with largest eigenvalue of  $\mathbf{P}^{T}$
- ► Computing *π* as eigenvector is computationally efficient and robust in some problems



• Can also write as (**I** is identity matrix,  $\mathbf{0} = [0, 0, ...]^T$ )

$$\left(\mathbf{I} - \mathbf{P}^{T}\right) \boldsymbol{\pi} = \mathbf{0} \qquad \boldsymbol{\pi}^{T} \mathbf{1} = 1$$

▶  $\pi$  has J elements, but there are J+1 equations  $\Rightarrow$  overdetermined

- ▶ If 1 is eigenvalue of  $\mathbf{P}^{T}$ , then 0 is eigenvalue of  $\mathbf{I} \mathbf{P}^{T}$ 
  - ▶  $\mathbf{I} \mathbf{P}^{\mathsf{T}}$  is rank deficient, in fact rank $(\mathbf{I} \mathbf{P}^{\mathsf{T}}) = J 1$
  - Then, there are in fact only J equations
- $\pi$  is eigenvector associated with eigenvalue 0 of  $I P^T$ 
  - $\pi$  spans null space of  $\mathbf{I} \mathbf{P}^T$  (not much significance)



MC with transition probability matrix

$$\mathbf{P} := \left( \begin{array}{ccc} 0 & 0.3 & 0.7 \\ 0.1 & 0.5 & 0.4 \\ 0.1 & 0.2 & 0.7 \end{array} \right)$$

Does P correspond to an ergodic MC?

- All states communicate with state 2 (full row and column P<sub>2j</sub> ≠ 0 and P<sub>j2</sub> ≠ 0 for all j)
- No transient states (irreducible with one recurrent state and finite)
- Aperiodic (period of state 2 is 1)
- ▶ Then, there exist  $\pi_1$ ,  $\pi_2$  and  $\pi_3$  such that  $\pi_j = \lim_{n\to\infty} P_{ij}^n$
- Limit is independent of *i*

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## Example: Aperiodic, irreducible MC (continued)



- How do we determine limit probabilities  $\pi_j$ ?
- Solve system of linear equations  $\pi_j = \sum_{i=0}^{\infty} \pi_i P_{ij}$  and  $\sum_{j=0}^{\infty} \pi_j = 1$

$$\begin{pmatrix} \pi_1 \\ \pi_2 \\ \pi_3 \\ 1 \end{pmatrix} = \begin{pmatrix} 0 & 0.1 & 0.1 \\ 0.3 & 0.5 & 0.2 \\ 0.7 & 0.4 & 0.7 \\ 1 & 1 & 1 \end{pmatrix} \begin{pmatrix} \pi_1 \\ \pi_2 \\ \pi_3 \end{pmatrix}$$

- The upper part of matrix above is P<sup>T</sup>
- There are three variables and four equations
  - Some equations might be linearly dependent
  - Indeed, summing first three equations:  $\pi_1 + \pi_2 + \pi_3 = \pi_1 + \pi_2 + \pi_3$
  - Always true, because probabilities in rows of P sum up to 1
  - This is because of rank deficiency of  $I P^T$
- Solution yields  $\pi_1 = 0.0909$ ,  $\pi_2 = 0.2987$  and  $\pi_3 = 0.6104$

## Stationary distribution



- Limit distributions are sometimes called stationary distributions
- ► Select initial distribution such that  $P[X_0 = i] = \pi_i$  for all *i*
- Probabilities at time n = 1 follow from theorem of total probability

$$P[X_1 = i] = \sum_{i=1}^{\infty} P[X_1 = j | X_0 = i] P[X_0 = i]$$

▶ Definitions of  $P_{ij}$ , and  $P[X_0 = i] = \pi_i$ . Algebraic property of  $\pi_j$ 

$$\mathsf{P}\left[X_1=i\right]=\sum_{i=1}^{\infty}P_{ij}\pi_i=\pi_j$$

Probability distribution is unchanged

- ▶ Proceeding recursively, system initialized with  $P[X_0 = i] = \pi_i$ ,
  - $\Rightarrow$  Probability distribution constant,  $P[X_n = i] = \pi_i$  for all n
- MC stationary in a probabilistic sense (states change, probs. do not)

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- Gambler's ruin problem
- Queues in communication networks: Transition probabilities
- Classes of States
- Limiting distributions
- Ergodicity

Queues in communication networks: Limit probabilities

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## Ergodicity



• Define  $T_i^{(n)}$  as fraction of time spent in *i*-th state up to time *n* 

$$T_i^{(n)} := \frac{1}{n} \sum_{m=1}^n \mathbb{I} \{ X_m = i \}$$

• Compute expected value of  $T_i^{(n)}$ 

$$\mathbb{E}\left[T_i^{(n)}\right] = \frac{1}{n} \sum_{m=1}^n \mathbb{E}\left[\mathbb{I}\left\{X_m = i\right\}\right] = \frac{1}{n} \sum_{m=1}^n \mathbb{P}\left[X_m = i\right] \to \pi_i$$

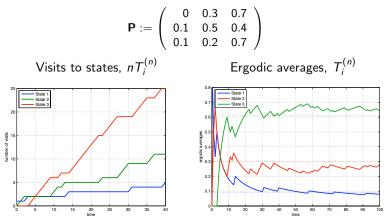
► As time  $n \to \infty$ , probabilities  $P[X_m = i]$  approach  $\pi_i$ . Then  $\lim_{t \to \infty} \mathbb{E}\left[T_i^{(n)}\right] = \lim_{t \to \infty} \frac{1}{n} \sum_{m=1}^n P[X_m = i] = \pi_i$ 

• For ergodic MCs same is true without expected value  $\Rightarrow$  ergodicity

$$\lim_{n \to \infty} T_i^{(n)} = \lim_{n \to \infty} \frac{1}{n} \sum_{m=1}^n \mathbb{I} \{ X_m = i \} = \pi_i, \quad \text{a.s.}$$



Recall transition probability matrix



Ergodic averages slowly converge to limit probabilities

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• Use of ergodic averages is more general than  $T_i^{(n)}$ 

#### Theorem

Consider an irreducible Markov chain with states  $X_n = 0, 1, 2, ...$  and stationary probabilities  $\pi_j$ . Let  $f(X_n)$  be a bounded function of the state X(n). Then, with probability 1

$$\lim_{n\to\infty}\frac{1}{n}\sum_{m=1}^n f(X_m) = \sum_{i=1}^\infty f(i)\pi_i$$

- $T_i^{(n)}$  is a particular case with  $f(X_m) = \mathbb{I}\{X_m = i\}$
- Think of  $f(X_m)$  as a reward associated with state X(m)
- $(1/n) \sum_{m=1}^{n} f(X_m)$  is the time average of rewards

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Proof.

• Because  $\mathbb{I} \{X_m = i\} = 1$  if and only if  $X_m = i$  we can write

$$\frac{1}{n}\sum_{m=1}^{n}f(X_m)=\frac{1}{n}\sum_{m=1}^{n}\left(\sum_{i=1}^{\infty}f(i)\mathbb{I}\left\{X_m=i\right\}\right)$$

• Change order of summations. Use definition of  $T_i^{(n)}$ 

$$\frac{1}{n}\sum_{m=1}^{n}f(X_{m}) = \sum_{i=1}^{\infty}f(i)\left(\frac{1}{n}\sum_{m=1}^{n}\mathbb{I}\{X_{m}=i\}\right) = \sum_{i=1}^{\infty}f(i)T_{i}^{(n)}$$

• Let  $n \to \infty$  in both sides

• Use ergodic average result for  $\lim_{n\to\infty} T_i^{(n)} = \pi_i$  [cf. page 67]

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- > There's more depth to ergodic results than meets the eye
- Ensemble average: across different realizations of the MC

$$\mathbb{E}\left[f(X_n)\right] = \sum_{i=1}^{\infty} f(i) \mathsf{P}\left(X_n = i\right) \to \sum_{i=1}^{\infty} f(i) \pi_i$$

Ergodic average: across time for a single realization of the MC

$$\bar{f}(n) = \frac{1}{n} \sum_{m=1}^{n} f(X_n)$$

- These quantities are fundamentally different but their values coincide asymptotically in n
- Observing one realization of the MC provides as much information as observing all realizations
- Practical consequence: Observe/simulate only one path of the MC

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- ► In some sense, still true if MC is periodic
- ► For irreducible positive recurrent MC (periodic or aperiodic) define

$$\pi_j = \sum_{i=0}^{\infty} \pi_i P_{ij}, \qquad \sum_{j=0}^{\infty} \pi_j = 1$$

- A unique solution exists (we say  $\pi_j$  are well defined)
- The fraction of time spent in state *i* converges to  $\pi_i$

$$\lim_{n \to \infty} T_i^{(n)} = \lim_{n \to \infty} \frac{1}{n} \sum_{m=1}^n \mathbb{I} \{ X_m = i \} \to \pi_i$$

 If MC is periodic, probabilities oscillate, but fraction of time spent in state *i* converges to π<sub>i</sub>

# Example: Periodic irreducible Markov chain



► Matrix **P** and state diagram of a periodic MC  $\mathbf{P} := \begin{pmatrix} 0 & 1 & 0 \\ 0.3 & 0 & 0.7 \\ 0 & 1 & 0 \end{pmatrix}$ 

• MC has period 2. If initialized with  $X_0 = 1$ , then

$$\begin{aligned} P_{11}^{2n+1} &= \mathsf{P}\left[X_{2n+1} = 1 \mid X_0 = 1\right] = 0, \\ P_{11}^{2n} &= \mathsf{P}\left[X_{2n} = 1 \mid X_0 = 1\right] = 1 \neq 0 \end{aligned}$$

• Define  $\boldsymbol{\pi} := [\pi_1, \pi_2, \pi_3]^T$  as solution of

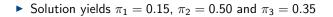
$$\begin{pmatrix} \pi_1 \\ \pi_2 \\ \pi_3 \\ 1 \end{pmatrix} = \begin{pmatrix} 0 & 0.3 & 0 \\ 1 & 0 & 1 \\ 0 & 0.7 & 0 \\ 1 & 1 & 1 \end{pmatrix} \begin{pmatrix} \pi_1 \\ \pi_2 \\ \pi_3 \end{pmatrix}$$

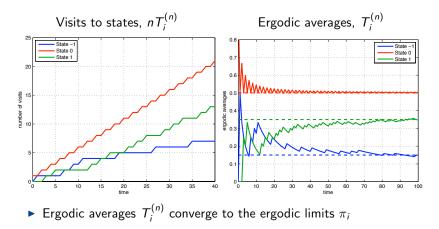
• Normalized eigenvector for eigenvalue 1 ( $\pi = \mathbf{P}^T \pi$ ,  $\pi^T \mathbf{1} = 1$ )

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# Example: Periodic irreducible MC (continued)









Powers of the transition probability matrix do not converge

$$\mathbf{P}^2 = \begin{pmatrix} 0.3 & 0 & 0.7 \\ 0 & 1 & 0 \\ 0.3 & 0 & 0.7 \end{pmatrix} \qquad \mathbf{P}^3 = \begin{pmatrix} 0 & 1 & 0 \\ 0.3 & 0 & 0.7 \\ 0 & 1 & 0 \end{pmatrix} = \mathbf{P}$$

▶ In general we have  $\mathbf{P}^{2n} = \mathbf{P}^2$  and  $\mathbf{P}^{2n+1} = \mathbf{P}$ 

At least one other eigenvalue of the transition probability matrix has modulus 1

$$\left|\operatorname{eig}_{2}\left(\mathbf{P}^{T}\right)\right|=1$$

• In this example, eigenvalues of  $\mathbf{P}^{T}$  are 1, -1 and 0

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- ► If MC is not irreducible it can be decomposed in transient (T<sub>k</sub>), ergodic (E<sub>k</sub>), periodic (P<sub>k</sub>) and null recurrent (N<sub>k</sub>) components
  - All of these are class properties
- ▶ Limit probabilities for transient states are null  $P[X_n = i] \rightarrow 0$ , for all  $X_n \in T_k$
- For arbitrary ergodic component  $\mathcal{E}_k$ , define conditional limits

$$\pi_{i} = \lim_{n \to \infty} \mathsf{P}\left[X_{n} = i \, \big| \, X_{0} \in \mathcal{E}_{k}\right], \quad \text{for all } i \in \mathcal{E}_{k}$$

• Result in page 58 is true with this (re)defined  $\pi_i$ 



• Likewise, for arbitrary periodic component  $\mathcal{P}_k$  (re)define  $\pi_j$  as

$$\pi_j = \sum_{i \in \mathcal{P}_k} \pi_i P_{ij}, \quad \sum_{j \in \mathcal{P}_k} \pi_j = 1, \quad \text{for all } j \in \mathcal{P}_k$$

A conditional version of the result in page 72 is true

$$\lim_{n\to\infty}T_i^{(n)}:=\lim_{n\to\infty}\frac{1}{n}\sum_{m=1}^n\mathbb{I}\left\{X_m=i\,\big|\,X_0\in\mathcal{P}_k\right\}\to\pi_i$$

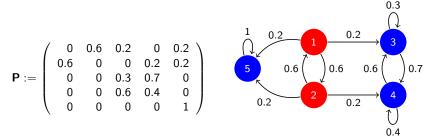
▶ For null recurrent components limit probabilities are null  $P[X_n = i] \rightarrow 0$ , for all  $X_n \in \mathcal{N}_k$ 

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Transition matrix and state diagram of a reducible MC



- States 1 and 2 are transient  $\mathcal{T} = \{1, 2\}$
- States 3 and 4 form an ergodic class  $\mathcal{E}_1 = \{3, 4\}$
- State 5 is a separate ergodic class  $\mathcal{E}_2 = \{5\}$



#### 10-step and 20 step transition probabilities

$\mathbf{P}^5 =$	/ 0	0.08	0.24	0.22	0.46 \	$\mathbf{P}^{10} =$	/0.00	0	0.23	0.27	0.50 \
	0.08	0	0.19	0.27	0.46		0	0.00	0.23	0.27	0.50
	0	0	0.46	0.54	0		0	0	0.46	0.54	0
	0	0	0.46	0.54	0		0	0	0.46	0.54	0
	\ 0	0	0	0	1/		\ 0	0	0	0	1/

- Transition into transient states is vanishing (columns 1 and 2)
- Transition from 3 and 4 into 3 and 4 only
  - If initialized in ergodic class  $\mathcal{E}_1 = \{3,4\}$  stays in  $\mathcal{E}_1$
- Transition from 5 only into 5
- ▶ From transient states T = {1,2} can go into either ergodic component E<sub>1</sub> = {3,4} or E<sub>2</sub> = {5}

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#### Matrix P can be separated in blocks

$$\mathbf{P} = \begin{pmatrix} \mathbf{0} & \mathbf{0.6} & \mathbf{0.2} & \mathbf{0} & \mathbf{0.2} \\ \mathbf{0.6} & \mathbf{0} & \mathbf{0} & \mathbf{0.2} & \mathbf{0.2} \\ \mathbf{0} & \mathbf{0} & \mathbf{0.3} & \mathbf{0.7} & \mathbf{0} \\ \mathbf{0} & \mathbf{0} & \mathbf{0.6} & \mathbf{0.4} & \mathbf{0} \\ \mathbf{0} & \mathbf{0} & \mathbf{0} & \mathbf{0} & \mathbf{1} \end{pmatrix} = \begin{pmatrix} \mathbf{P}_{\mathcal{T}} & \mathbf{P}_{\mathcal{T}\mathcal{E}_1} & \mathbf{P}_{\mathcal{T}\mathcal{E}_2} \\ \mathbf{0} & \mathbf{P}_{\mathcal{E}_1} & \mathbf{0} \\ \mathbf{0} & \mathbf{0} & \mathbf{P}_{\mathcal{E}_2} \end{pmatrix}$$

• Block  $P_T$  describes transition between transient states

- ▶ Blocks  $P_{\mathcal{E}_1}$  and  $P_{\mathcal{E}_2}$  describe transitions in ergodic components
- ▶ Blocks  $\mathbf{P}_{\mathcal{T}\mathcal{E}_1}$  and  $\mathbf{P}_{\mathcal{T}\mathcal{E}_2}$  describe transitions from  $\mathcal{T}$  to  $\mathcal{E}_1$  and  $\mathcal{E}_2$

Powers of n can be written as

$$\mathbf{P}^{n} = \begin{pmatrix} \mathbf{P}_{\mathcal{T}}^{n} & \mathbf{Q}_{\mathcal{T}\mathcal{E}_{1}} & \mathbf{Q}_{\mathcal{T}\mathcal{E}_{2}} \\ 0 & \mathbf{P}_{\mathcal{E}_{1}}^{n} & 0 \\ 0 & 0 & \mathbf{P}_{\mathcal{E}_{2}}^{n} \end{pmatrix}$$

▶ The transient transition block converges to 0,  $\lim_{n\to\infty} \mathbf{P}_{\mathcal{T}}^n = \mathbf{0}$ 

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- ▶ As *n* grows the MC hits an ergodic state with probability 1
- Henceforth, MC stays within ergodic component

$$\mathsf{P}\left[X_{n+m} \in \mathcal{E}_i \mid X_n \in \mathcal{E}_i\right] = 1, \text{ for all } m$$

For large *n* suffices to study ergodic components
 ⇒ MC behaves like a MC with transition probabilities P<sub>E1</sub>
 ⇒ Or like one with transition probabilities P<sub>E2</sub>

- We can think that all MCs as ergodic
- Ergodic behavior cannot be inferred a priori (before observing)
- Becomes known a posteriori (after observing sufficiently large time)

Culture micro: Something is known a priori if its knowledge is independent of experience (MCs exhibit ergodic behavior). A posteriori knowledge depends on experience (values of the ergodic limits). They are inherently different forms of knowledge (search for Immanuel Kant)



Markov chains. Definition and examples

- Chapman Kolmogorov equations
- Gambler's ruin problem
- Queues in communication networks: Transition probabilities
- Classes of States
- Limiting distributions
- Ergodicity

Queues in communication networks: Limit probabilities



- Communication system: Move packets from source to destination
- Between arrival and transmission hold packets in a memory buffer
- ► Example problem, buffer design: Packets arrive at a rate of 0.45 packets per unit of time and depart at a rate of 0.55. How many packets the buffer needs to hold to have a drop rate smaller than 10<sup>-6</sup> (one packet dropped for every million packets handled)
- Time slotted in intervals of duration  $\Delta t$
- During each time slot n
  - $\Rightarrow$  A packet arrives with prob.  $\lambda$ , arrival rate is  $\lambda/\Delta t$
  - $\Rightarrow$  A packet is transmitted with prob.  $\mu$ , departure rate is  $\mu/\Delta t$
- No concurrence: No simultaneous arrival and departure (small  $\Delta t$ )



- Future queue lengths depend on current length only
- Probability of queue length increasing

$$\mathsf{P}\left[q_{n+1}=i+1 \mid q_n=i\right] = \lambda, \qquad \text{for all } i$$

• Queue length might decrease only if  $q_n > 0$ . Probability is

$$\mathsf{P}\left[q_{n+1}=i-1 \mid q_n=i\right]=\mu, \qquad \text{for all } i>0$$

Queue length stays the same if it neither increases nor decreases

$$\mathsf{P}\left[q_{n+1}=i \mid q_n=i\right] = 1 - \lambda - \mu, \quad \text{for all } i > 0$$
$$\mathsf{P}\left[q_{n+1}=0 \mid q_n=0\right] = 1 - \lambda$$

• No departures when  $q_n = 0$  explain second equation

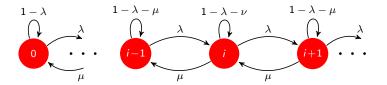
## Queue as a Markov chain (reminder)



- MC with states  $0, 1, 2, \ldots$  Identify states with queue lengths
- Transition probabilities for  $i \neq 0$  are

$$P_{i,i-1} = \lambda, \qquad P_{i,i} = 1 - \lambda - \mu, \qquad P_{i,i+1} = \mu$$

• For 
$$i = 0$$
  $P_{0,0} = 1 - \lambda$  and  $P_{01} = \lambda$ 



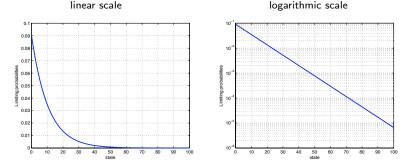
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### Numerical example: Limit probabilities

Penn

- Build matrix **P** truncating at maximum queue length L = 100
- Arrival rate  $\lambda = 0.3$ . Departure rate  $\mu = 0.33$
- Find eigenvector of  $\mathbf{P}^{T}$  associated with largest eigenvalue (i.e., 1)
- Yields limit probabilities  $\pi = \lim_{n \to \infty} \mathbf{p}(n)$

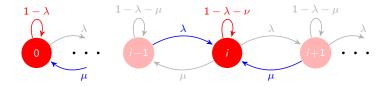


Limit probabilities appear linear in logarithmic scale

 $\Rightarrow$  Seemingly implying an exponential expression  $\pi_i \propto lpha$ 

## Limit distribution equations





Limit distribution equations for state 0 (empty queue)

$$\pi_0 = (1-\lambda)\pi_0 + \mu\pi_1$$

► For the remaining states

$$\pi_i = \lambda \pi_{i-1} + (1 - \lambda - \mu)\pi_i + \mu \pi_{i+1}$$

• Propose candidate solution  $\pi_i = c\alpha^i$ 

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• Substitute candidate solution  $\pi_i = c\alpha^i$  in equation for  $\pi_0$ 

$$c\alpha^{0} = (1 - \lambda)c\alpha^{0} + \mu c\alpha^{1} \quad \Rightarrow \quad 1 = (1 - \lambda) + \mu \alpha$$

• The above equation is true if we make  $\alpha = \lambda/\mu$ 

- Does  $\alpha = \lambda/\mu$  verify the remaining equations ?
- From the equation for generic  $\pi_i$  (divide by  $c\alpha^{i-1}$ )

$$c\alpha^{i} = \lambda c\alpha^{i-1} + (1 - \lambda - \mu)c\alpha^{i} + \mu c\alpha^{i+1}$$
$$\mu\alpha^{2} - (\lambda + \mu)\alpha + \lambda = 0$$

 $\blacktriangleright$  The above quadratic equation is satisfied by  $\alpha=\lambda/\mu$ 

• And  $\alpha = 1$ , which is irrelevant



• Determine *c* so that probabilities sum 1  $(\sum_{i=0}^{\infty} \pi_i = 1)$ 

$$\sum_{i=0}^{\infty} \pi_i = \sum_{i=0}^J c (\lambda/\mu)^i = \frac{c}{1-\lambda/\mu} = 1$$

- Used geometric sum
- Solving for c and substituting in  $\pi_i = c\alpha^i$  yields

$$\pi_i = (1 - \lambda/\mu) \left(\frac{\lambda}{\mu}\right)^i$$

- $\blacktriangleright$  The ratio  $\mu/\lambda$  is the queues' stability margin
- Larger  $\mu/\lambda \Rightarrow$  larger probability of having few queued packets

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### Queue balance equations

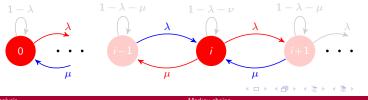


▶ Rearrange terms in equation for limit probabilities [cf. page 87]

 $\lambda \pi_0 = \mu \pi_1$  $(\lambda + \mu)\pi_i = \lambda \pi_{i-1} + \mu \pi_{i+1}$ 

- $\lambda \pi_0$  is average rate at which the queue leaves state 0
- Likewise  $(\lambda + \mu)\pi_i$  is the rate at which queue leaves state *i*
- $\mu\pi_0$  is average rate at which the queue enters state 0
- $\lambda \pi_{i-1} + \mu \pi_{i+1}$  is rate at which queue enters state *i*
- Limit equations prove validity of queue balance equations

Rate at which leaves = Rate at which enters





- Packets may arrive and depart in same time slot (concurrence)
- ▶ Queue evolution equations remain the same, [cf. 35]
- But queue probabilities change [cf. 84]
- Probability of queue length increasing (for all i)

$$P[q_{n+1} = i + 1 | q_n = i] = P[A_n = 1] P[D_n = 0] = \lambda(1 - \mu)$$

• Queue length might decrease only if  $q_n > 0$  (for all i > 0)

$$P[q_{n+1} = i - 1 | q_n = i] = P[D_n = 1] P[D_n = 0] = \mu(1 - \lambda)$$

Queue length stays the same if it neither increases nor decreases

$$\mathsf{P} [q_{n+1} = i \mid q_n = i] = \lambda \mu + (1 - \lambda)(1 - \mu) = \nu, \quad \text{for all } i > 0$$
  
$$\mathsf{P} [q_{n+1} = 0 \mid q_n = 0] = (1 - \lambda) + \lambda \mu$$

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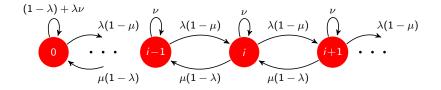
# Limit distribution / queue balance equations



- Write limit distribution equations  $\Rightarrow$  queue balance equations
- Rate at which leaves = rate at which enters

$$\lambda(1-\mu)\pi_0 = \mu(1-\lambda)\pi_1$$
$$(\lambda(1-\mu)+\mu(1-\lambda))\pi_i = \lambda(1-\mu)\pi_{i-1}+\mu(1-\lambda)\pi_{i+1}$$

• Propose exponential solution  $\pi = c\alpha^i$ 





• Substitute candidate solution in equation for  $\pi_0$ 

$$\lambda(1-\mu)c = \mu(1-\lambda)clpha \quad \Rightarrow \quad lpha = rac{\lambda(1-\mu)}{\mu(1-\lambda)}$$

• Same substitution in equation for generic  $\pi_i$ 

$$\mu(1-\lambda)c\alpha^2 + (\lambda(1-\mu) + \mu(1-\lambda))c\alpha + \lambda(1-\mu)c = 0$$

• which as before is solved for  $\alpha = \lambda(1-\mu)/\mu(1-\lambda)$ 

• Find constant c to ensure  $\sum_{i=0}^{\infty} c \alpha^i = 1$  (geometric series). Yields

$$\pi_i = (1 - lpha) lpha^i = \left(1 - rac{\lambda(1 - \mu)}{\mu(1 - \lambda)}
ight) \left(rac{\lambda(1 - \mu)}{\mu(1 - \lambda)}
ight)$$

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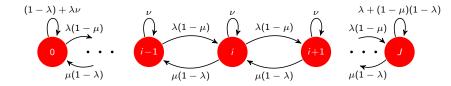
## Limited queue size



- Packets dropped if there are too many packets in queue
- Too many packets in queue, then delays too large, packets useless when they arrive. Also preserve memory
- Equation for state J requires modification (rate leaves = rate enters)

$$\mu(1-\lambda)\pi_J = \lambda(1-\mu)\pi_{J-1}$$

•  $\pi_i = c \alpha^i$  with  $\alpha = \lambda (1 - \mu) / \mu (1 - \lambda)$  also solve this equation (Yes!)





- Limit probabilities are not the same because constant c is different
- ► To compute *c*, sum a finite geometric series

$$1 = \sum_{i=0}^{J} c \alpha^{i} = c \frac{1 - \alpha^{J+1}}{1 - \alpha} \quad \Rightarrow \quad c = \frac{1 - \alpha}{1 - \alpha^{J+1}}$$

Limit distributions for the finite queue are then

$$\pi_i = \frac{1-\alpha}{1-\alpha^{J+1}} \alpha^i \approx (1-\alpha) \alpha^i$$

• with  $\alpha = \lambda(1-\mu)/\mu(1-\lambda)$ , and approximation valid for large J

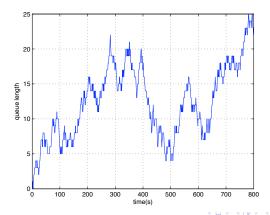
- ► Approximation for large J yields same result as infinite length queue
  - As it should

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#### Simulations



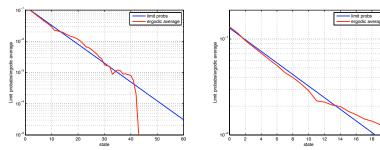
- Arrival rate  $\lambda = 0.3$ . Departure rate  $\mu = 0.33$ . Resulting  $\alpha \approx 0.87$
- Maximum queue length J = 100. Initial state  $q_0 = 0$  (queue empty)
  - Not the same as initial probability distribution



Queue lenght as function of time

# Simulations: Average occupancy and limit distributionPenn

- Average time spent at each queue state is predicted by limit distribution
- For i = 60 occupancy probability is  $\pi \approx 10^{-5}$ .
  - Explains inaccurate prediction for large i



60 states





- If  $\lambda = 0.45$  and  $\mu = 0.55$  how many packets the buffer needs to hold to have a drop rate smaller than  $10^{-6}$  (one packet dropped for every million packets handled)
- What is the probability of buffer overflow?
- ▶ Packet discarded if queue is in state J and a new packet arrives

$$\mathsf{P}\left[\text{overflow}\right] = \lambda \pi_J = \frac{1-\alpha}{1-\alpha^{J+1}} \lambda \alpha^J \approx (1-\alpha) \lambda \alpha^J$$

- With  $\lambda = 0.45$  and  $\mu = 0.55$ ,  $\alpha \approx 0.82 \Rightarrow J \approx 57$
- A final caveat
  - Still assuming only 1 packet arrives per time slot
  - Lifting this assumption requires introduction of continuous time Markov chains