CIS 1200 Final Exam May 9, 2023

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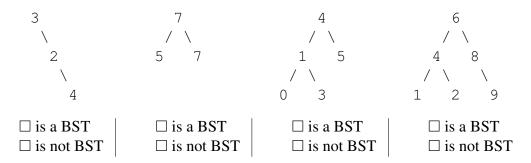
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I certify that I have complied with th completing this examination.	te University of Pennsylvania's Code of Academic Integrity in
Signature:	Date:

- There are 120 total points. The exam period is 120 minutes long.
- There are 13 pages in the exam and an Appendix for your reference.
- Please begin by writing your PennKey (e.g., stevez) at the bottom of all the odd-numbered pages in the rest of the exam.
- Please skim the entire exam first—some of the questions will take significantly longer than others.
- Do not spend too much time on any one question. Be sure to recheck all of your answers.
- We will ignore anything you write on the Appendix.
- For coding problems: aim for accurate syntax, but we will not grade your code style for indentation, spacing, etc.
- If you need extra space for an answer, you may use the scratch page at the end of the exam; make sure to clearly indicate that you have done this in the normal answer space for the problem.
- Good luck!

OCaml Concepts (27 points total)

Appendix A contains the definitions of several OCaml types and functions that should be familiar from class: 'a tree, 'a ref, transform, and fold.

Binary Search Trees (8 points) For each tree below, check the box indicating whether it satisfies the *binary search tree* invariants. **NOTE:** If it does not, cross out and replace *one* node value such that the resulting **int** tree *does* satisfy the invariants (the result does not need to have the same nodes).



Higher-order Functions (8 points)

Consider the following four functions defined using transform or fold.

```
let hof1 = transform (fun x -> insert x 120)
let hof2 = fold (fun x acc -> insert acc x) Empty
let hof3 = fold (fun x acc -> (insert x 120) :: acc) []
let hof4 = fold (fun x acc -> insert x 120) Empty
```

Match each of the functions to one of the following English descriptions of its behavior. (choose one option for each function; a choice may be used by more than one function)

- (A) Returns a list of trees obtained by adding the value 120 to every tree in a given list.
- (B) Returns either Empty or the tree obtained by adding 120 to the first tree in a given list.
- (C) Returns a (binary search) tree containing the elements of a provided list.
- (D) Returns a list of trees obtained by adding the value 120 to Empty for every tree in a given list.

• hof1 is described by: □ A	\square B	\Box C	\Box D
• hof2 is described by: □ A	\square B	\Box C	\Box D
• hof3 is described by: □ A	\square B	\Box C	\Box D
• hof4 is described by: □ A	\square B	\Box C	\Box D

Aliasing and mutable state Recall that OCaml supports mutable state via type declarations such as for the 'a ref type shown below (and in Appendix A). Consider the following well-typed program:

type 'a re	f = { mutable	contents : 'a	}		
let y = { ;; x.conte let z = y	<pre>contents = "F nts <- y nts <- "RULES</pre>		O " } }		
<pre>;; print_e ;; print_e</pre>	ndline ("y.co	ontents.contents ontents = " ^ y.o ontents = " ^ z.o " ^ w)	contents)	is.contents)	
(a) (4 poin	ts) What is the	type of each variabl	le in the above code	e? (choose one each)	
• x :	\square string	\square string ref	\square unit ${f ref}$	\square (string ref)	ref
• y :	\square string	\square string ref	\square unit ${f ref}$	\square (string ref)	ref
• z :	\square string	\square string ref	\square unit ${ t ref}$	\square (string ref)	ref
• w :	\square string	\square string ref	\square unit ${f ref}$	\square (string ref)	ref
each slot wit x.contents y.contents z.contents	*	ngs 1200, PENN, or	0 1	when this program	is run. Fill
□ x □ x.co		z z.content		ll <i>alias</i> y? (mark all t	:hat apply)

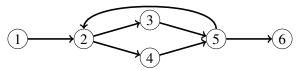
Java Programming

The next several questions refer to the Java code found in Appendix C. Some of the questions test your understanding of Java concepts; other questions will ask you to follow our design process to implement (parts of) a collection datatype for *graphs*. You may find the (excerpt of) the Java Documentation found in Appendix B to be useful.

Step 1: Understand the problem (4 points total)

Besides the *set* and *finite map* collections that we have studied in class, another frequently used collection type is the *graph*. A *graph* contains a set of *nodes* along with a collection of (directed) *edges*, each of which connects a *source* node to a *target* node. We consider graphs where there is at most one edge between any pair of nodes.

We draw nodes as labeled circles (n) and edges as arrows pointing from the source to the target (s)—(t). For example, the diagram below depicts a graph whose nodes are the integers 1 through 6, and whose edges are indicated by the seven arrows.



As with other collections, there are many operations that we might want our graphs to support. In addition to methods to add nodes and edges to the graph, we focus on just two other operations. First, determining the *neighbors* of a node, which is just the set of nodes reachable via an edge. For instance, in the graph above, the neighbors of 2 are 3 and 4, but node 6 has no neighbors. Second, we can determine whether there is a *path* from some node 6 to a node 6 that follows the edges in the graph. For instance, in the example graph above, there are paths $1 \rightarrow 2 \rightarrow 4 \rightarrow 5$ and $5 \rightarrow 2 \rightarrow 3$ but there is no path from node $6 \rightarrow 6$ to node $6 \rightarrow 6$. Note that there is always a (trivial) path from any node to itself and that *a path might visit the same node multiple times*.

(a)	(2 points) that apply)		mple graph	above, whi	ich nodes a	re <i>neighbor</i>	s of node (5)?	(mark all
		\square (2)	\square (3)	\Box \bigcirc	\Box \bigcirc \bigcirc	\Box (6)		
(b)	(2 points)	How many	paths are t	here from n	ode 4 to 1	node 3 ? (c	choose one)	
	$\square 0$							
	\Box 1							
	\square 2							
	☐ there are	e infinitely	many paths	due to the	cycle in the	graph		

Step 2: Design the interface (14 points total)

Just as with Java's <code>Set<E></code> and <code>Map<K</code>, <code>V></code> interfaces, which are generic in the data they store, we will make the interface <code>Graph<Node></code> polymorphic in the type <code>Node</code>. The <code>Graph<Node></code> interface supports add and <code>contains</code> operations with the same specification of those methods in <code>Set</code>; it also supports three new graph-specific methods: <code>addEdge</code>, <code>neighbors</code>, and <code>hasPath</code>. Code defining the interface is shown in Appendix C.1.

Java concepts: (12 points) With respect to Graph<Node>, indicate whether the following statements are true or false; if false give a brief justification.

1. According to the method signature on line 25, addEdge might throw an IOException.
□ True
☐ False because:
2. According to the method signature on line 25, addEdge might throw a NullPointerException
□ True
☐ False because:
3. The following snippet of code will compile successfully (but may produce warnings):
<pre>Graph<integer> g = null;</integer></pre>
g.contains("CIS1200");
☐ False because:
4. The following snippet of code will compile successfully (but may produce warnings):
<pre>Graph<integer> g = new Graph<>(); g.contains(3);</integer></pre>
☐ True
☐ False because:
Design question: (2 points) Suppose that, rather than neighbors, the Graph <node> interface provides a method boolean hasEdge(Node src, Node tgt), which returns true if there is an edge from src to tgt (and throws NoSuchElementException if src is not a node in the graph).</node>
1. Consider implementing hasEdge using neighbors: (choose one)
\square It is not possible to implement has Edge using neighbors.
\square g.contains(src)&& g.neighbors(tgt) implements g.hasEdge(src,tgt)
\square g.neighbors(src).contains(tgt) implements g.hasEdge(src,tgt)
\square g.neighbors(tgt).contains(src) implements g.hasEdge(src,tgt)
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Step 3: Write test code for Graph (12 points total) Even before we have code that implements the Graph<Node> interface, we can write test cases that check our examples and properties. These test cases can use *subtype polymorphism* to implement tests that work for any implementation: we will assume that each such generic test is provided a instance of a graph object freshly created by new. For instance, in Step 4 (later in the exam) you will work with TreeGraph. We will assume that a test for TreeGraph is called via test (new TreeGraph()).

```
(a) (4 points) Consider the following test case:
      private void testNoSuchSource(Graph<Integer> g) {
2
           assertFalse(g.contains(0));
3
           assertThrows(NoSuchElementException.class, () -> q.neighbors(0));
4
       Which of the following statements are true? (mark all that apply)
       ☐ This test case will succeed only if g.contains(0) returns false.
       ☐ The syntax NoSuchElementException.class (line 3) creates an instance of an anonymous
          inner class.
       ☐ The syntax () -> g.neighbors (0) (line 3) creates an instance of an anonymous inner class.
       \Box The following test case is equivalent to the one above.
            private void testNoSuchSourceAlternate(Graph<Integer> g) {
                 assertFalse(g.contains(0));
                 try {
                     g.neighbors(0);
                 } catch (NoSuchElementException e) { fail(); }
             }
```

1

(b) (2 points) Complete the following test case so that all assertions succeed. It is based on the example graph pictured earlier:

```
private void testNeighbors(Graph<Integer> g) {
        g.addEdge(1,2); g.addEdge(2,3); g.addEdge(2,4); g.addEdge(3,5);
        g.addEdge(4,5); g.addEdge(5,6); g.addEdge(5,2);
        TreeSet<Integer> expected1 = new TreeSet<>();
       expected1.add(_____);
       assertEquals(expected1, g.neighbors(1));
       TreeSet<Integer> expected2 = new TreeSet<>();
       expected2.add(3);
       expected2.add(____
       assertEquals(expected2, g.neighbors(_____));
       TreeSet<Integer> expected3 = new TreeSet<>();
       assertEquals(expected3, g.neighbors(_____));
    }
```

	h and/or pseudocode) how to write a test case that can perly encapsulate some state associated with the graph
•	
tations of the Graph interface that use diffe	Collections library, we might have different implementerent internal representations. We will follow the design uplement the hasPath algorithm once so that different can share that code.
representations of the graph accuracy type of	
	nestions are about the implementation of the AbstractGraph be able to answer these questions without understanding n; they are just about Java concepts.)
1. (4 points) Which of the following a that apply)	re supertypes of AbstractGraph <integer>? (mark all</integer>
☐ AbstractGraph <integer></integer>	AbstractGraph <object></object>
\square AbstractSet <integer> \square AbstractGraph<node></node></integer>	☐ Set <node> ☐ Object</node>
2. (4 points) Suppose a well-typed pr	rogram declares a variable Graph <integer> g = (* ng statements are true? (mark all that apply)</integer>
☐ The <i>static type</i> of g is Graph <i< th=""><td>nteger>.</td></i<>	nteger>.
**	lass associated with g to be AbstractGraph <integer></integer>
☐ It is possible for the <i>dynamic class</i>	ass associated with g to be a subtype of AbstractGraph
☐ In code after this declaration it exception.	is possible for the expression g.equals(g) to throw an
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, -	

that the implementation assumes that Node implements the Comparable <node> interface. Which of the following best explains why? (choose one)</node>
☐ The implementation of this.contains, as used on line 11, requires compareTo in its implementation.
☐ The method toSearch.removeFirst(), as used on line 24, needs compareTo to find the smallest node to remove from the list.
☐ The implementation of hasPath uses a TreeSet <node> to store the alreadyVisited nodes, as seen on lines 16, 25, and 30, and TreeSet requires its elements to support compareTo.</node>
☐ The method current.equals(tgt), as used on line 27, requires compareTo in its implementation.
4. (4 points) Which of the following best describes what would happen if we removed the use of this from the conditional guard on line 11, i.e., so it reads:
<pre>if (!contains(src) !contains(tgt)){</pre>
(choose one)
\square The hasPath method would always throw a NoSuchElementException.
\square The hasPath method would never throw a NoSuchElementException.
☐ The hasPath method would always go into an infinite loop.
☐ The behavior of the hasPath method would be unaffected.

Design question Now we will implement *two* different instances of the <code>Graph<Integer></code> interface. Both inherit from <code>AbstractGraph<Integer></code> and provide the missing <code>Graph<Integer></code> methods. (Note that, for these types, the nodes are specialized to be <code>Integer</code> objects.)

(no more questions on this page)

Implementation 1: ArrayGraph (14 points total) The first implementation represents the nodes and edges of the graph using arrays of booleans (sometimes called the "adjacency matrix"). Appendix C.3 contains code for this class.

The following are the *representation invariants* embodied by this code.

<pre>INV1 nodes[n] == true if and only if n is a node in the graph INV2 edges[src][tgt] == true if and only if src and tgt are nodes in the graph and there an edge from src to tgt</pre>	is
 (a) (2 points) Which of the following best explains why the constructor (line 10) establish these invariants? (choose one) The graph starts out with maxNodes nodes, and the default initializer for boolean is true The graph starts out with maxNodes nodes, and the default initializer for boolean is false. The graph starts out empty, and the default initializer for boolean is true. The graph starts out empty, and the default initializer for boolean is false. 	
 (b) (4 points) The add method is supposed to add a new node to the graph. • Which line of code <i>relies on</i> INV1? (choose one): □ line 22 □ line 23 □ line 24 □ line 25 • Which line of code <i>establishes</i> INV1? (choose one): □ line 22 □ line 23 □ line 24 □ line 25 	
 (c) (2 points) Which invariant would break if we delete line 37 from addEdge? (choose one) INV1 breaks	
Set <integer> nbrs = new TreeSet<>();</integer>	
return nbrs; } PennKey:	9

Implementation 2: TreeGraph (17 points) The second implementation represents the nodes and edges of the graph using the Java TreeMap and/or TreeSet collection(s). From the options below, choose appropriate *representation type(s)* or mark "not needed" if you don't need that field. Then write down the invariant and complete the missing parts of the TreeGraph following that plan.

```
class TreeGraph extends AbstractGraph<Integer> {
   private □ Set<Integer>
                          ☐ Set<Set<Integer>>
           ☐ Set<Map<Integer, Integer>>
                                                nodes; // or □ NOT NEEDED
   ☐ Set<Map<Integer, Integer>
                                               edges; // or □ NOT NEEDED
    /* INVARIANT: FILL IN BELOW
    * n is a node exactly when:
    * s -> t is an edge exactly when:
    */
   public TreeGraph() { /* Initializes the field(s) using new TreeSet or new TreeMap */ }
   public boolean add(Integer node) {
       if (!this.contains(node)) { /* FILL IN HERE: */
           return true;
       return false;
   public boolean contains(Object o) { /* FILL IN HERE: */
   public void addEdge(Integer src, Integer tgt) {
       this.add(src);
       this.add(tgt); /* FILL IN HERE: */
   public Set<Integer> neighbors(Integer src) {
       if (!this.contains(src)) { throw new NoSuchElementException(); }
       Set<Integer> nbrs = new TreeSet<>(); /* FILL IN HERE: */
       return nbrs;
   }
```

Using a Graph (16 points total) Appendix C.4 contains a program that reads the edges of a graph from a file example.txt and then prints out, for each node src in the graph, a list of nodes that can be reached via a path from src. The file format is simple: nodes are numbers and each line of the file is of the form s -> t, representing an edge from s to t in the graph. Below you can see the sample example.txt associated with the example graph, along with the output printed to the console:

example.txt:	console c	output:	
1 -> 2	1 ==> 1 2	2 3 4 5	6
2 -> 3	2 ==> 2 3	3 4 5 6	
2 -> 4	3 ==> 2 3	3 4 5 6	
3 -> 5	4 ==> 2 3	3 4 5 6	
4 -> 5	5 ==> 2 3	3 4 5 6	
5 -> 6	6 ==> 6		
5 -> 2			
Java Concepts (a) (2 points) this keyword:	Suppose that v	we chang	e line 25 of GraphApp to instead use the
Graph <integer> g = this.read</integer>	Graph(new FileR	eader(fil	.ename));
What would be the result? (cho	ose one)		
☐ The program would comp	pile successfully	and its b	ehavior would be unchanged.
☐ The program would comp	pile successfully	but it wo	ould throw an exception when run.
\Box The program would not c	ompile because	main $is\ d$	eclared as static.
\Box The program would not c	ompile because	readGra	ph is declared as static.
2	wing are true pr	roperties	BufferedReader as part of your TwitterBoot the BufferedReader class. (Note: we leader.) (mark all that apply)
☐ BufferedReader is more from the input.	efficient than ju	ust using	FileReader to read individual characters
\square The BufferedReader me	thods do not thr	row IOEx	ceptionS.
☐ The BufferedReader re once as a String.	adLine method	provides	the ability to read a whole line of input at
\square A BufferedReader cons	tructor can acce	pt any Re	ader as an input .

Scratch Space

Use this page for work that you do not want us to grade. If you run out of space elsewhere in the exam and you do want to put something here that we should grade, make sure to put a clear note in the normal answer space for the problem in question.
PennKey: 13