CIS 1200 Midterm II November 15, 2024

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Name:	
PennKey (penn login, e.g., bcpierce):	
PennID (the "numbers", e.g., 12001200):	

I certify that I have complied with the University of Pennsylvania's Code of Academic Integrity in completing this examination.

Signature: Date: _____

- Please wait to begin the exam until you are told it is time for everyone to start.
- When you begin, start by writing your PennKey at the bottom of all the odd-numbered pages in the rest of the exam.
- There are 120 total points. The time for the exam is 60 minutes.
- For coding problems, aim for accurate syntax, but we will not grade your code for indentation, spacing, etc.
- There are 13 pages in the exam and an appendix for your reference. Do not write any answers in the appendix as they will not be graded.
- Do not spend too much time on any one question. Be sure to recheck all of your answers.
- If you need extra space for an answer, you may use the scratch page at the end of the exam; make sure to clearly indicate that you have done this in the normal answer space for the problem.
- Good luck!

1. Tail Recursion (18 points total)

Which of the following functions are tail recursive?

1.1 (2 points) let rec mystery (l: int list) (count: int) : int = begin match 1 with | [] -> count | hd::tl -> mystery tl (count + 1) end □ Not Tail Recursive □ Tail Recursive 1.2 (2 points) let rec mystery (l: bool list) : bool = begin match 1 with | [] -> **false** | hd::tl -> not (mystery tl) end □ Tail Recursive □ Not Tail Recursive 1.3 (2 points) let rec mystery (1: bool list) : bool = begin match 1 with | [] -> true | hd::tl -> hd && mystery tl end □ Tail Recursive □ Not Tail Recursive 1.4 (2 points) let rec mystery (l: 'a list) (f: 'a -> 'a) : 'a = begin match 1 with | [] -> failwith "error" | hd::[] -> f hd | hd::tl -> f (mystery tl f) end \Box Tail Recursive \Box Not Tail Recursive

1.5 (10 points)

This function is *not* tail recursive:

```
let rec mystery (f : int -> int) (x : int) : int =
    if x = 0 then 0
    else (f x) + (mystery f (x - 1))
```

In the space below, define a tail-recursive function that behaves the same as mystery on all inputs (except that it might loop where the version above would overflow the stack).

```
let rec mystery (f : int -> int) (x : int) : int =
```

2. Deques (29 points total)

Recall the type definitions for the deque data structure in Homework 4.

```
type 'a dqnode = {
   v: 'a;
   mutable next: 'a dqnode option;
   mutable prev: 'a dqnode option;
}
type 'a deque = {
   mutable head: 'a dqnode option;
   mutable tail: 'a dqnode option;
```

The *deque invariant* is as follows:

```
(1) head and tail are both tagged None or both tagged Some, and
(2) if head = Some h and tail = Some t, then
(2a) t is reachable from h by following next pointers
(2b) t.next = None
(2c) h is reachable from t by following prev pointers
(2d) h.prev = None
and, for every node n in the queue,
(2e) if n.next = Some m, then m.prev = Some n
(2f) if n.prev = Some m, then m.next = Some n.
```

Your job in this problem will be to help define a function $split_deque$ that, given a deque d and a ' a element x, splits d into two deques at the first occurrence of x, putting x at the head of the second deque.

For instance, if d is a deque with three elements such that to_list d = [1; 2; 3], then split_deque d 3 will return deques (d1, d2) such that to_list d1 = [1; 2] and to_list d2 = [3].

If x is not equal to any element of d, then d1 should be the same as d and d2 should be empty. If d is empty, then two empty deques should be returned.

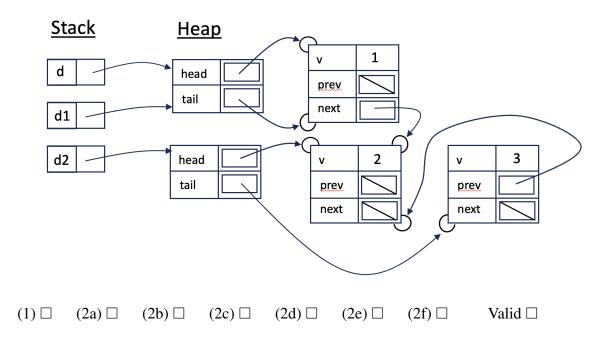
Note that the split_deque operation does not allocate any new dqnodes—it reuses the ones from the original deque d.

Nothing for you to answer on this page

2.1 (6 points) To begin, let's check our understanding of the deque invariant. In this part of the question (and the next), we'll show you a picture representing the state of the ASM after running a candidate implementation of split_deque. The implementation splits the deque in the right place, but the resulting two deques may or may not be valid—i.e., they might not maintain the deque invariants.

The input d is a valid deque with to_list d = [1; 2; 3] before a call to (split_deque d 2).

Indicate which, if any, of the deque invariants are *broken* in the picture. Check all boxes that apply, or "Valid" if no invariants are broken in d1 or d2.



The deque invariant again, for quick reference:

- (1) head and tail are both tagged None or both tagged Some, and
- (2) if head = Some h and tail = Some t, then
 - (2a) t is reachable from h by following next pointers
 - (2b) t.next = None
 - (2c) h is reachable from t by following prev pointers
 - (2d) h.prev = None

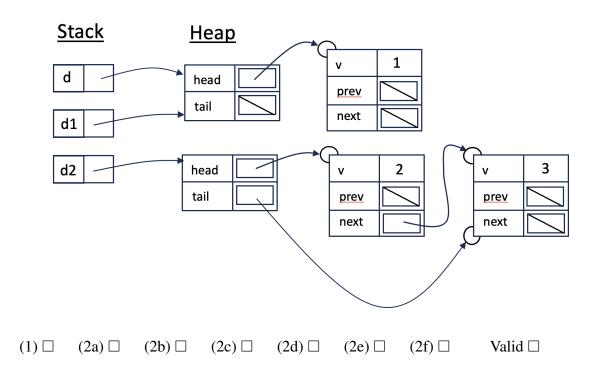
and, for every node n in the queue,

(2e) if n.next = Some m, then m.prev = Some n

(2f) if n.prev = Some m, then m.next = Some n.

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2.2 (6 points) Again, indicate which, if any, of the deque invariants are *broken* in the picture.



The deque invariant again, for quick reference:

- (1) head and tail are both tagged None or both tagged Some, and
- (2) if head = Some h and tail = Some t, then
 - (2a) t is reachable from h by following next pointers
 - (2b) t.next = None
 - (2c) h is reachable from t by following prev pointers
 - (2d) h.prev = None

and, for every node n in the queue,

- (2e) if n.next = Some m, then m.prev = Some n
- (2f) if n.prev = Some m, then m.next = Some n.

```
2.3 (17 points) Complete the code below for split_deque.
    let split_deque (d : 'a deque) (x: 'a) : 'a deque * 'a deque =
      let rec loop (dno: 'a dqnode option) : 'a deque * 'a deque =
        begin match dno with
          | None -> (* consider which two cases can lead here... *)
          | Some n ->
             if n.v = x then
                (* n becomes head of second deque *)
               begin match n.prev with
                 | None ->
                    (* consider which edge case this is... *)
                 | Some p ->
                   (* create new deque... *)
                   let d2 = {head = ____; tail = ____} in
                   (* make any other appropriate updates... *)
                   (* fill in appropriate return value... *)
               end
             else
        end
      in loop d.head
```

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3. OCaml ASM (23 points total)

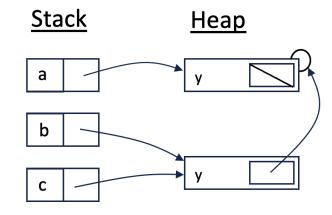
Suppose we've defined the following OCaml type:

type t = {mutable y : t option}

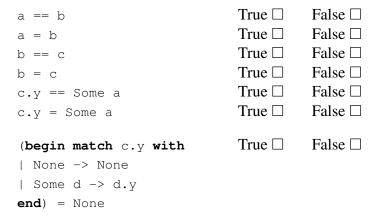
Then executing the following code...

let a : t = {y = None}
let b : t = {y = Some a}
let c : t = b

... will yield this state of the OCaml ASM:



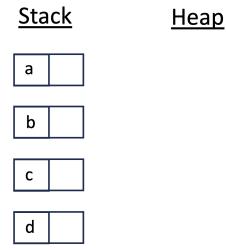
3.1 (7 points) Does evaluating the following expressions now yield true or false? (Recall that == in OCaml is reference equality and = is structural equality.)



3.2 (8 points) Suppose we start from an empty stack and heap and execute the following code.

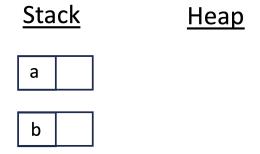
let a : t = {y = None}
let b : t = {y = Some a}
let c : t = {y = Some b}
let d : t = {y = Some c}
;; a.y <- Some d</pre>

Complete the following ASM diagram to show the state after this code is executed.



3.3 (8 points) Suppose, instead, that we start from an empty stack and heap and execute this code.

Complete the following ASM diagram.



4. GUI Programming in OCaml (30 points total)

Key definitions for our OCaml GUI library can be found on page 1 of the Appendix.

We can use the GUI library to create a rudimentary gaming console where we can control the position of a ball (the circle in the leftmost box below) on a 1D board, using two buttons labeled Left and Right.



4.1 (21 points) Below is incomplete code to build this widget. Fill in the blanks to create a widget where pressing "Left" will shift the ball 1 pixel to the left. (Pressing "Right" would shift the ball 1 pixel to the right, but we have omitted that part of the code for brevity.) Select functions from gctx.mli and widget.mli are provided in the Appendix. See the description for each blank in the comments above the blank.

```
1 let make_game (init_pos : int) : widget =
    (* Allocate an internal state to track the ball's position: *)
 2
    let pos =
 3
    let circle drawing (g: Gctx.gctx) : unit =
       (* draw a 5 by 5 ellipse for the ball at its current position: *)
 4
       Gctx.draw_ellipse _____ (____, 10) 5 5 in
    let board, canvas_nc = canvas (100,20) circle_drawing in
 5
    (* Generate a button with text "Left"; make sure to name the
       first component of the result left_w: *)
    let (_____, _____) = _____
 6
                                                      in
 7
    let left () =
        (* Move ball's position 1 pixel to the left: *)
                  <- max ( - 1) 5 in
 8
    (* Install an event listener to handle clicks on the left button: *)
 9
        _____.add_event_listener (mouseclick_listener _____);
    <...code omitted for right button function and event_listener...>
10
```

11 let full_game = hlist [board; border left_w; border right_w] in
12 border full_game

4.2 For each of the following edits, decide whether they would (1) continue to compile, (2) maintain the appearance of the make_game widget on the screen, and (3) maintain the functionality / playability of the game.

(3 points) Replace line 12 with bor	der board
Still compiles	
Maintains appearance	🗆 No
Maintains function	□ No
(3 points) Move lines 3–5 below lin Still compiles □ Yes	e 10
☐ Tes Maintains appearance ☐ Yes	\Box No
Maintains function	□ No
(3 points) Replace line 11 with let full_game = hpai	r (border left_w) (hpair (border right_w) board) in
Still compiles	□ No
Maintains appearance	□ No
Maintains function	□ No

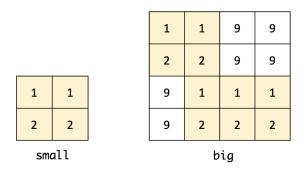
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5. Java Arrays (20 points)

The goal in this problem will be to write a function numberOfAppearances that takes two square 2D arrays as parameters:

- array small of size $m \times m$, where m > 0
- array big of size $n \times n$, where $m \le n$

The numberOfAppearances function should return an integer indicating the number of times small appears as a *subarray* of big. For example, if small and big look like this...



... then numberOfAppearances(small, big) should return 3. Note that different occurrences of small inside big may overlap.

You may assume the arrays are non-null.

On the next page, we've given you the start of a helper method called *isSubArray*, which checks whether one array is a subarray of another *at a given position*. First, implement this method; then use it to complete the implementation of numberOfAppearances.

Nothing for you to answer on this page

```
public static int numberOfAppearances(int[][] small, int[][] big) {
    int m = small.length;
    int n = big.length;
    int count = 0;
```

}

}

Scratch Space

Use this page for work that you do not want us to grade. If you run out of space elsewhere in the exam and you **do** want to put something here that we should grade, make sure to put a clear note in the normal answer space for the problem in question.