

# Programming Languages and Techniques (CIS120)

Lecture 8

Feb 5, 2014

Abstract Types: Sets  
Modules and Interfaces

# Announcements

- Homework 3 is available
  - Due **TUESDAY**, February 11<sup>th</sup> at 11:59:59pm
  - Practice with BSTs, generic functions, and *abstract types*
- If you added CIS 120 recently, make sure that you can see your scores online.
  - If you get feedback about your scores, you are in our database.
  - If not, please send mail to [tas120@lists.seas.upenn.edu](mailto:tas120@lists.seas.upenn.edu)
  - If you see unsubmitted “quizzes”, you may need to register your clicker
- Read chapter 9 of the lecture notes

Do these two declarations produce the same BST?

```
let t1 = insert Empty 2  
let t2 = insert (insert Empty 2) 2
```

1. yes
2. no

Answer: yes

Do these two declarations produce the same BST?

```
let t1 = insert (insert (insert Empty 2) 1) 3  
let t2 = insert (insert (insert Empty 2) 3) 1
```

1. yes
2. no

Answer: yes

Are you familiar with the idea of a *set* from mathematics?

1. yes
2. no

Answer: about 70% reported yes

# Abstract Collections

# A *set* is an abstraction

- A set is a collection of data
  - In math, we typically write sets like this:  $\emptyset$   $\{1,2,3\}$   $\{\text{true},\text{false}\}$  with operations like:  $S \cup T$  or  $S \cap T$  for union and intersection; we write  $x \in S$  to mean that “x is a member of the set S”
- A set is a lot like a list, except:
  - Order doesn't matter
  - Duplicates don't matter
  - *It isn't built into OCaml*
- Sets show up frequently in applications
  - Examples: set of students in a class, set of coordinates in a graph, set of answers to a survey, set of data samples from an experiment, ...

# Abstract type: set

- A binary search tree is an *implementation* of a *set*
  - *there is an empty set*
  - *there is a way to list all elements contained in the set (inorder)*
  - *there is a way to test membership (lookup)*
  - *could define union/intersection with insert and delete*
- Order doesn't matter
  - We create BSTs by adding elements to an empty BST
  - The BST data structure doesn't remember what order we added the elements
- Duplicates don't matter
  - Our implementation doesn't keep track of how many times an element is added
- BSTs are not the only way to implement sets, let's generalize



# Abstract type: set

- An abstract type is defined by its *interface* and its *properties*
- The interface defines how sets can be created and used
  - There is an empty set
  - There is a way to add elements to a set to make a bigger set
  - There is a way to list all elements in a set
  - There is a way to remove elements from the set to make a smaller set
  - There is a way to test membership
- The properties define how these operations interact with each other
  - Elements that were added can be found in the set
  - Adding a twice doesn't change the elements of a set
  - Adding in a different order doesn't change the elements of a set
  - ....
- Any type that can implement this interface while satisfying the properties can be a set

# Sets in action

# A design problem

*As a high-school student, Stephanie had the job of reading books and finding which words, out of a list of the 1000-most common SAT vocabulary words, appeared in a particular book. She enjoyed being paid to read, but she would have enjoyed being paid to program more. How could she have automated this task?*

1. What are the important concepts or *abstractions* for this problem?
  - The list of words that appear in a book
  - The set of 1000-most common SAT words
  - The set of words from the list that are contained in the set

## 2. Formalize the Interface

- Suppose we had a generic type of sets:

'a set

(We'll get to the details of that in a moment.)

- We can formalize the interface for our problem:

```
let countVocab (text : string list)
               (vocab : string set)
               : int =
  failwith "write me"
```

## 3. Write Test Cases

```
let vocab : string set =  
  list_to_set ["induce"; "crouching"; "reprieve";  
              "indigent"; "arrogate"; "coalesce";  
              "temerity"]  
  
let text1 = ["i"; "looked"; "up"; "again"; "at";  
            "the"; "crouching"; "white"; "shape"; "and";  
            "the"; "full"; "temerity"; "of"; "my"; "voyage"]  
  
let test () : bool =  
  countVocab text1 vocab = 2  
;; run_test "countVocab" test
```

Test cases specify the *interface* and the *properties* of the necessary abstractions.

## 4. Implement the Required Behavior

```
let countVocab (text : string list)
               (vocab : string set)
               : int =
  failwith "write me"
```

- Easy recursive programming task
  - (we'll leave the details to you)
- Requires set membership test

```
let member (x:'a) (s:'a set) : bool =
  failwith "unimplemented"
```

# The set interface in OCaml (*a signature*)

```
module type Set = sig
  type 'a set

  val empty : 'a set
  val add    : 'a -> 'a set -> 'a set
  val remove : 'a -> 'a set -> 'a set
  val list_to_set : 'a list -> 'a set
  val member : 'a -> 'a set -> bool
  val elements : 'a set -> 'a list
end
```

Keyword 'val' names values that must be defined and their types.

# Aside: Function Types

- In OCaml, the type of functions from input `t` to output `u` is written:

```
t -> u
```

- Functions with multiple arguments are written with multiple arrows
- Examples:

```
size : tree -> int
hamming_distance : helix -> helix -> int
acids_of_helix   : helix -> acids list
length : 'a list -> int
zip      : 'a list -> 'b list -> ('a*'b) list
lookup  : tree -> int -> bool
insert  : 'a tree -> 'a -> 'a tree
```



# A *module* of sets

- An implementation of the set interface will look like this:

Name of the module

Signature that it implements

```
module Myset : Set = struct
  ...
  (* implementations of all the operations *)
  ...
end
```

# Testing (and using) sets

- To use the values defined in the set module use the “dot” syntax:

`Myset.<member>`

- Note: Module names are always capitalized in OCaml

```
let s1 = Myset.add 3 Myset.empty
let s2 = Myset.add 4 Myset.empty
let s3 = Myset.add 4 s1
```

```
let test () : bool = (Myset.member 3 s1) = true
;; run_test "Myset.member 3 s1" test
```

```
let test () : bool = (Myset.member 4 s3) = true
;; run_test "Myset.member 4 s3" test
```

# Testing (and using) sets

- Alternatively, use “`open`” to bring all of the names defined in the interface into scope.

```
;; open Myset
```

```
let s1 = add 3 empty
```

```
let s2 = add 4 empty
```

```
let s3 = add 4 s1
```

```
let test () : bool = (member 3 s1) = true
```

```
;; run_test "Myset.member 3 s1" test
```

```
let test () : bool = (member 4 s3) = true
```

```
;; run_test "Myset.member 4 s3" test
```

# Implementing sets

- There are many ways to implement sets.
  - lists, trees, arrays, etc.
- **How do we choose which implementation?**
- Many such implementations are of the flavor “a set is a ... with some invariants”
  - A set is a *list* with no repeated elements.
  - A set is a *tree* with no repeated elements
  - A set is a *binary* search tree
  - A set is an *array of bits*, where 0 = absent, 1 = present
- **How do we preserve the invariants of the implementation?**

# Abstract types

BIG IDEA: Hide the *concrete representation* of a type behind an *abstract interface* to preserve invariants.

- The interface **restricts** how other parts of the program can interact with the data.
- Benefits:
  - **Safety:** The other parts of the program can't break any invariants
  - **Modularity:** It is possible to change the implementation without changing the rest of the program

# Set signature

```
module type Set = sig
```

```
  type 'a set
```

```
  val empty    : 'a set
```

```
  val add      : 'a -> 'a set -> 'a set
```

```
  val remove   : 'a -> 'a set -> 'a set
```

```
  val list_to_set : 'a list -> 'a set
```

```
  val member   : 'a -> 'a set -> bool
```

```
  val elements : 'a set -> 'a list
```


```
end
```

Type declaration has no  
“body” – its representation  
is *abstract*!

# Implement the set Module

```
module MySet : Set =  
  struct  
    type 'a tree =  
      | Empty  
      | Node of 'a tree * 'a * 'a tree  
  
    type 'a set = 'a tree  
  
    let empty : 'a set = Empty  
  
    ...  
  end
```

Module must define the type declared in the signature




- The implementation has to include everything promised by the interface
  - It can contain *more* functions and type definitions (e.g. auxiliary functions) but those cannot be used outside the module
  - The types of the provided implementations must match the interface

# Another Implementation

```
module MySet2 : Set =  
struct
```

```
  type 'a set = 'a list
```

A different definition for  
the type set



```
  let empty : 'a set = []
```

```
  ...
```

```
end
```



Does this code type check?

```
;; open MySet  
let s1 : int set = Empty
```

1. yes
2. no

Answer: no, the Empty data constructor is not available outside the module

Does this code type check?

```
;; open MySet  
let s1 : int set = add 1 empty
```

1. yes
2. no

Answer: yes

Does this code type check?

```
;; open MySet  
let s1 : int tree = add 1 empty
```

1. yes
2. no

Answer: no, add constructs a set, not a tree

If a module works and starts with:

```
;; open MySet
```

will it continue to work if we change that line to:

```
;; open MySet2
```

1. yes
2. no

Answer: yes (caveat: times may be different)