

# Programming Languages and Techniques (CIS120)

## Lecture 15

October 4<sup>th</sup>, 2015

### Mutable Queues

Lecture notes: Chapter 16

# Announcements

- No Lab Sections this week
  - Fall break!
- HW 4: Mutable Queues is available
  - Due: Tuesday, October 13<sup>th</sup> at 11:59 pm
- Midterm 1 is graded.
  - Scores available on the HW submission web page
  - More details on Weds. (after make-up exams have been taken)

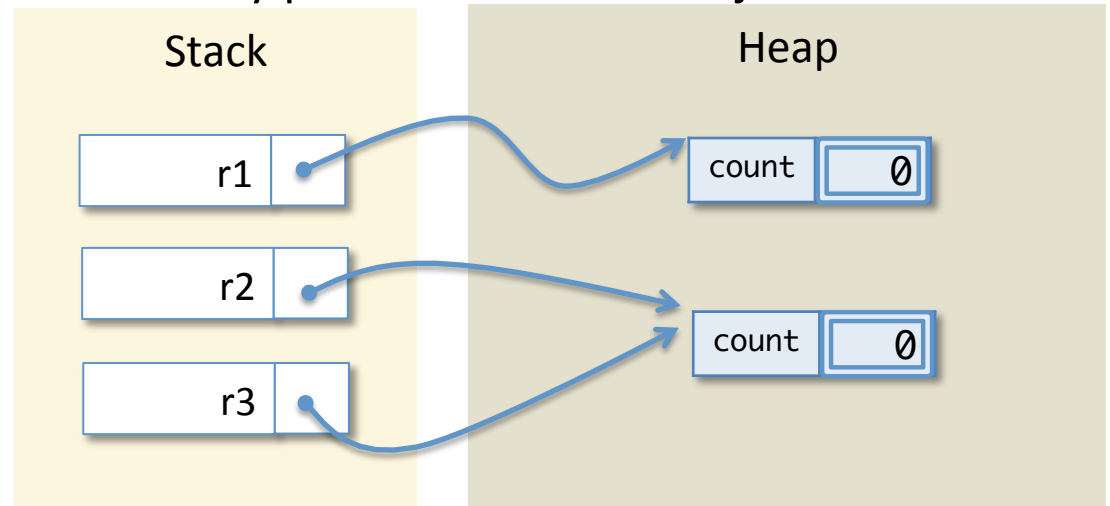
# Reference and Equality

= vs. ==

# Reference Equality

- Suppose we have two counters. How do we know whether they share the same internal state?
  - `type counter = { mutable count : int }`
  - We could increment one and see whether the other's value changes.
  - But we could also just test whether the references alias directly.
- Ocaml uses `'=='` to mean *reference equality*:
  - two reference values are `'=='` if they point to the same object in the heap; so:

```
r2 == r3
not (r1 == r2)
r1 = r2
```



# Structural vs. Reference Equality

- *Structural (in)equality*:  $v1 = v2$        $v1 \neq v2$ 
  - recursively traverses over the *structure* of the data, comparing the two values' components for structural equality
  - function values are never structurally equivalent to anything
  - structural equality can go into an infinite loop (on cyclic structures)
  - appropriate for comparing *immutable* datatypes
- *Reference (in)equality*:  $v1 == v2$        $v1 \neq v2$ 
  - Only looks at where the two references point in the heap
  - function values are only equal to themselves
  - equates strictly fewer things than structural equality
  - appropriate for comparing *mutable* datatypes

What is the result of evaluating the following expression?

```
let p1 : point = { x = 0; y = 0; } in  
let p2 : point = p1 in  
  
p1 = p2
```

1. true
2. false
3. runtime error
4. compile-time error

Answer: true

What is the result of evaluating the following expression?

```
let p1 : point = { x = 0; y = 0; } in  
let p2 : point = p1 in  
  
p1 == p2
```

1. true
2. false
3. runtime error
4. compile-time error

Answer: true

What is the result of evaluating the following expression?

```
let p1 : point = { x = 0; y = 0; } in  
let p2 : point = { x = 0; y = 0; } in  
  
p1 == p2
```

1. true
2. false
3. runtime error
4. compile-time error

Answer: false

What is the result of evaluating the following expression?

```
let p1 : point = { x = 0; y = 0; } in
let p2 : point = { x = 0; y = 0; } in
let l1 : point list = [p1] in
let l2 : point list = [p2] in

l1 = l2
```

1. true
2. false
3. runtime error
4. compile-time error

Answer: true

What is the result of evaluating the following expression?

```
let p1 : point = { x = 0; y = 0; } in
let p2 : point = p1 in
let l1 : point list = [p1] in
let l2 : point list = [p2] in

l1 == l2
```

1. true
2. false
3. runtime error
4. compile-time error

Answer: false

# Putting State to Work

Mutable Queues

Have you ever implemented the mutable data structure called a **linked list**, in any language?

1. yes
2. no
3. not sure

# A design problem

*Suppose you are implementing a website to sell tickets to a very popular music event. To be fair, you would like to allow people to select seats first come, first served. How would you do it?*

- Understand the problem
  - Some people may visit the website to buy tickets while others are still selecting their seats
  - Need to remember the order in which people purchase tickets
- Define the interface
  - Need a data structure to store ticket purchasers
  - Need to add purchasers to the *end* of the line
  - Need to allow purchasers at the *beginning* of the line to select seats
  - Both kinds of access must be efficient to handle the volume

# (Mutable) Queue Interface

```
module type QUEUE =  
sig  
  (* abstract type *)  
  type 'a queue  
  
  (* Make a new, empty queue *)  
  val create : unit -> 'a queue  
  
  (* Determine if the queue is empty *)  
  val is_empty : 'a queue -> bool  
  
  (* Add a value to the end of the queue *)  
  val enq : 'a -> 'a queue -> unit  
  
  (* Remove the first value (if any) and return it *)  
  val deq : 'a queue -> 'a  
  
end
```

We can tell, just looking at this interface, that it is for a MUTABLE data structure. How?

Because queues are mutable, we must allocate a new one every time we need one.

Adding an element to the queue returns unit because it modifies the given queue.

# Specify the behavior via test cases

```
let test () : bool =  
  let q : int queue = create () in  
  enq 1 q;  
  enq 2 q;  
  1 = deq q  
;; run_test "queue test 1" test
```

```
let test () : bool =  
  let q : int queue = create () in  
  enq 1 q;  
  enq 2 q;  
  let _ = deq q in  
  2 = deq q  
;; run_test "queue test 2" test
```

What value should replace ??? so that the following test passes?

```
let test () : bool =  
  let q : int queue = create () in  
  enq 1 q;  
  let _ = deq q in  
  enq 2 q;  
  ??? = deq q  
  
;; run_test "enq after deq" test
```

1. 1
2. 2
3. None
4. failwith "empty queue"

Answer: 2

# Implementing Linked Queues

Representing links

# Data Structure for Mutable Queues

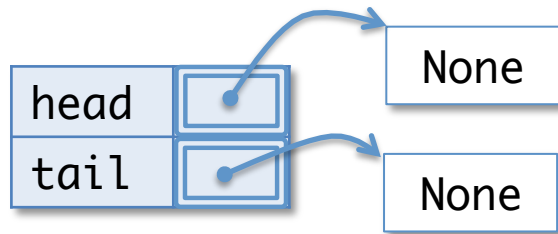
```
type 'a qnode = {  
    v: 'a;  
    mutable next : 'a qnode option  
}  
  
type 'a queue = { mutable head : 'a qnode option;  
                  mutable tail : 'a qnode option }
```

There are two parts to a mutable queue:

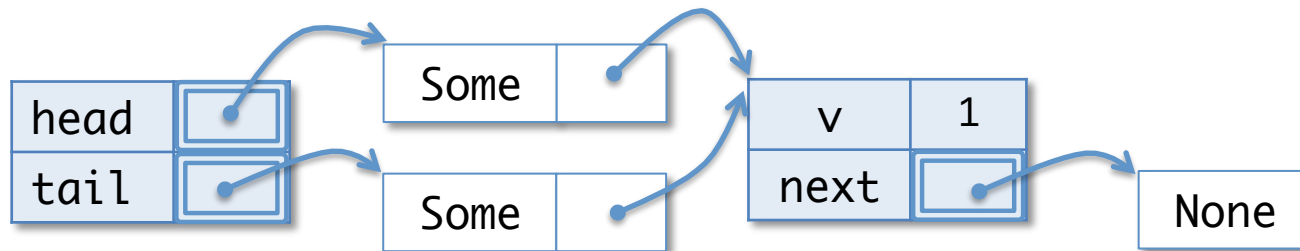
1. the “internal nodes” of the queue, with links from one to the next
2. a record with links to the head and tail nodes

All of the links are *optional* so that the queue can be empty.

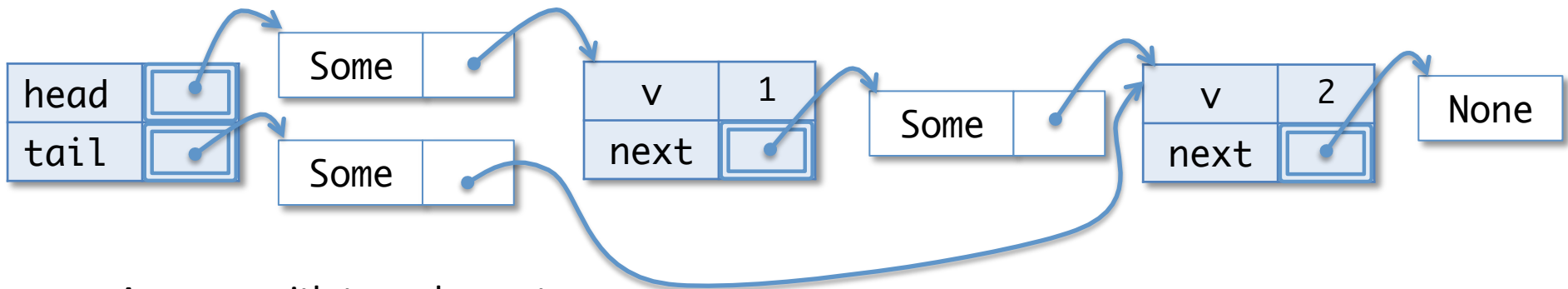
# Queues in the Heap



An empty queue



A queue with one element

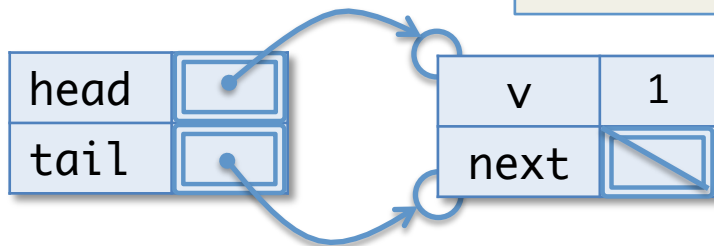


A queue with two elements

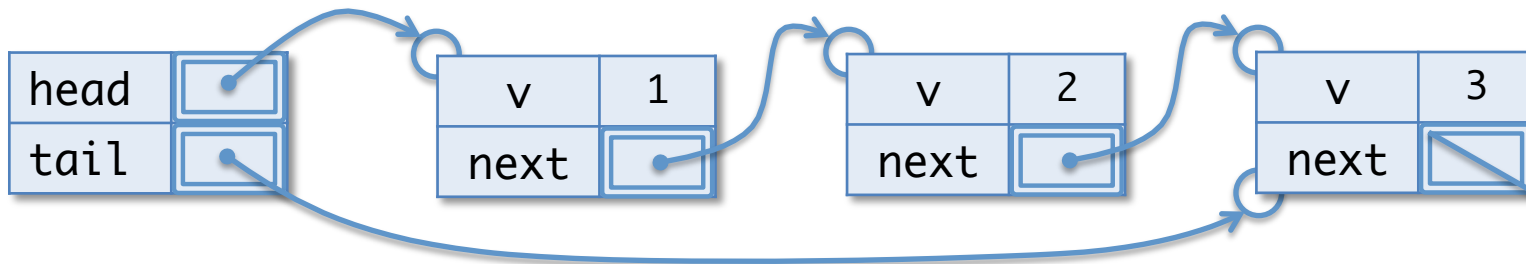
# Visual Shorthand: Abbreviating Options



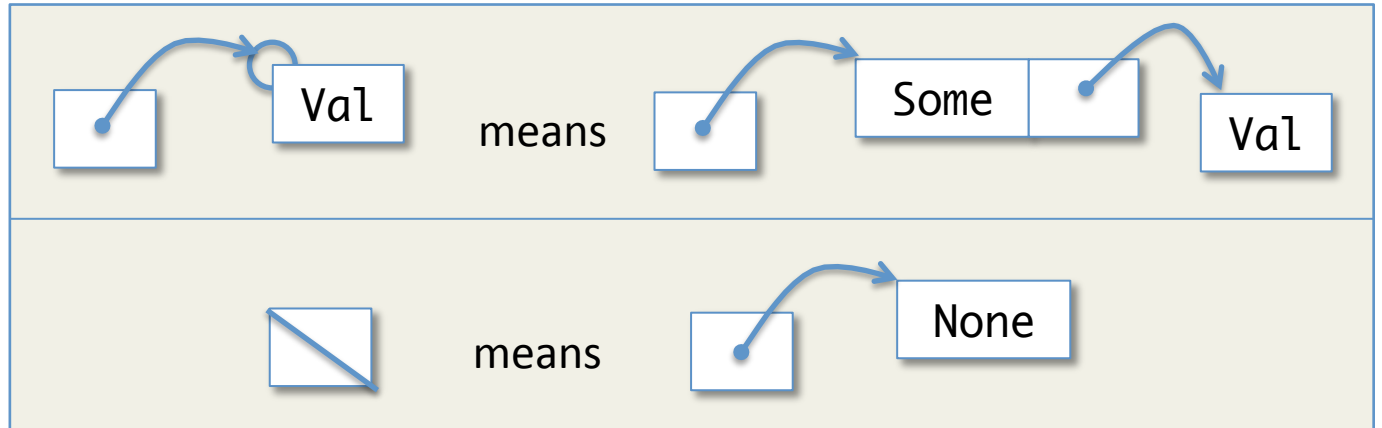
An empty queue



A queue with one element



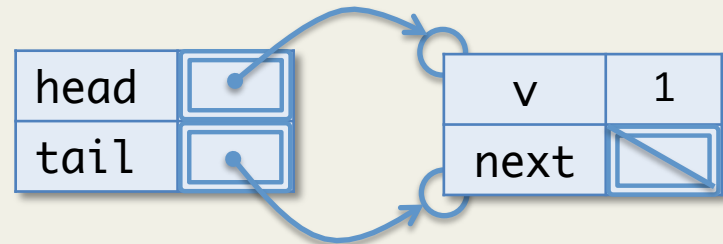
A queue with three elements



Given the queue datatype shown below, which expression creates a 1-element queue in the heap:

```
type 'a qnode = {  
  v: 'a;  
  mutable next : 'a qnode option  
}
```

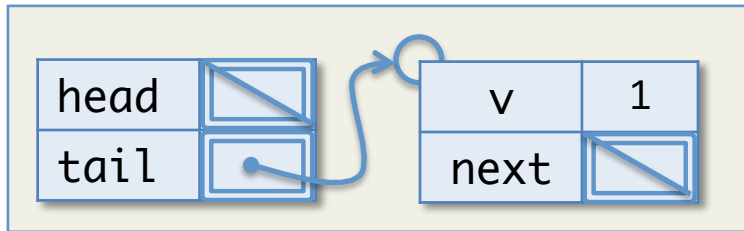
```
type 'a queue = { mutable head : 'a qnode option;  
                  mutable tail : 'a qnode option }
```



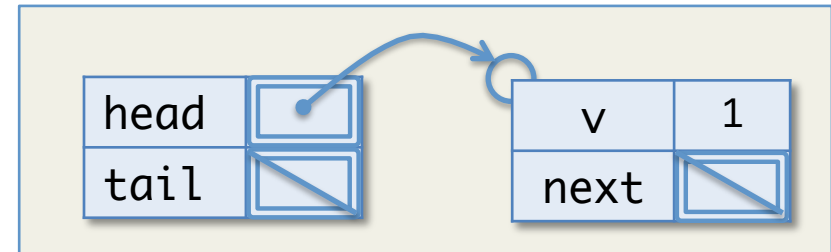
1. `let q = { head = None; tail = None }`
2. `let q = { head = 1; tail = None }`
3. `let q = let qn = { v= 1; next = None } in  
          { head = qn; tail = None }`
4. `let q = let qn = { v= 1; next = None } in  
          { head = Some qn; tail = Some qn }`

Answer: 4

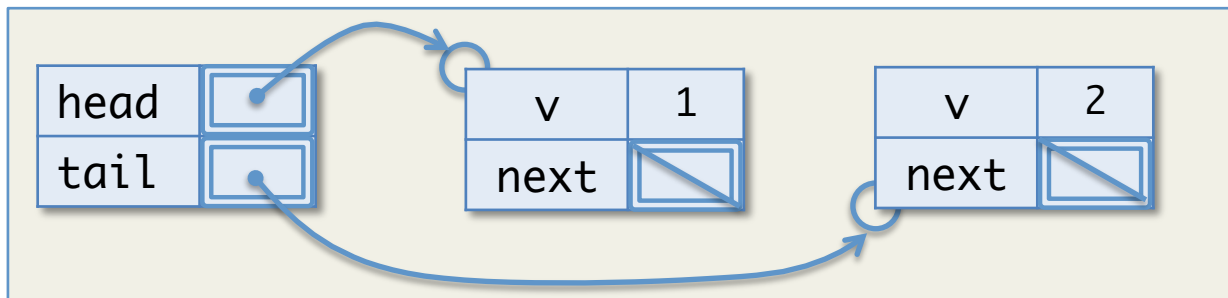
# “Bogus” values of type `int` queue



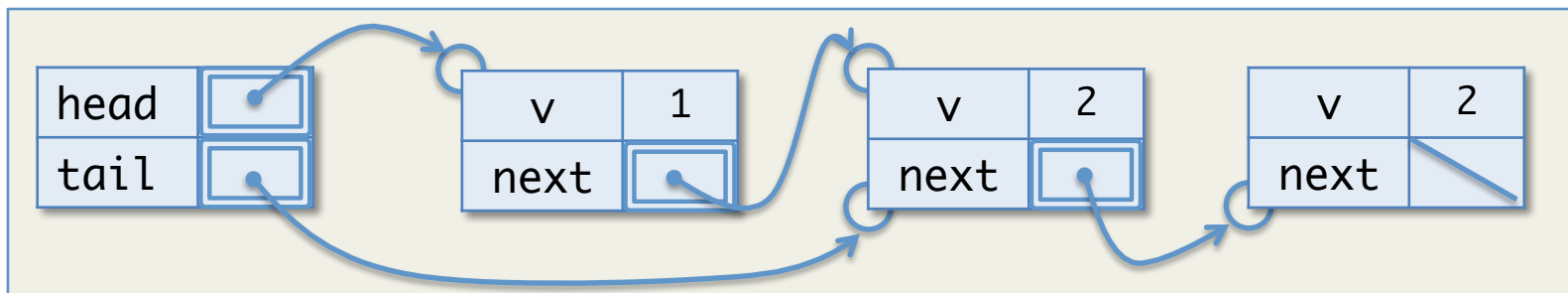
head is `None`, tail is `Some`



head is `Some`, tail is `None`



tail is not reachable from the head



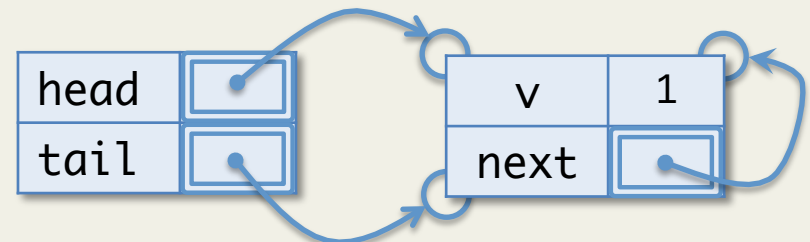
tail doesn't point to the last element of the queue

Given the queue datatype shown below, is it possible to create a *cycle* of references in the heap. (i.e. a way to get back to the same place by following references.)

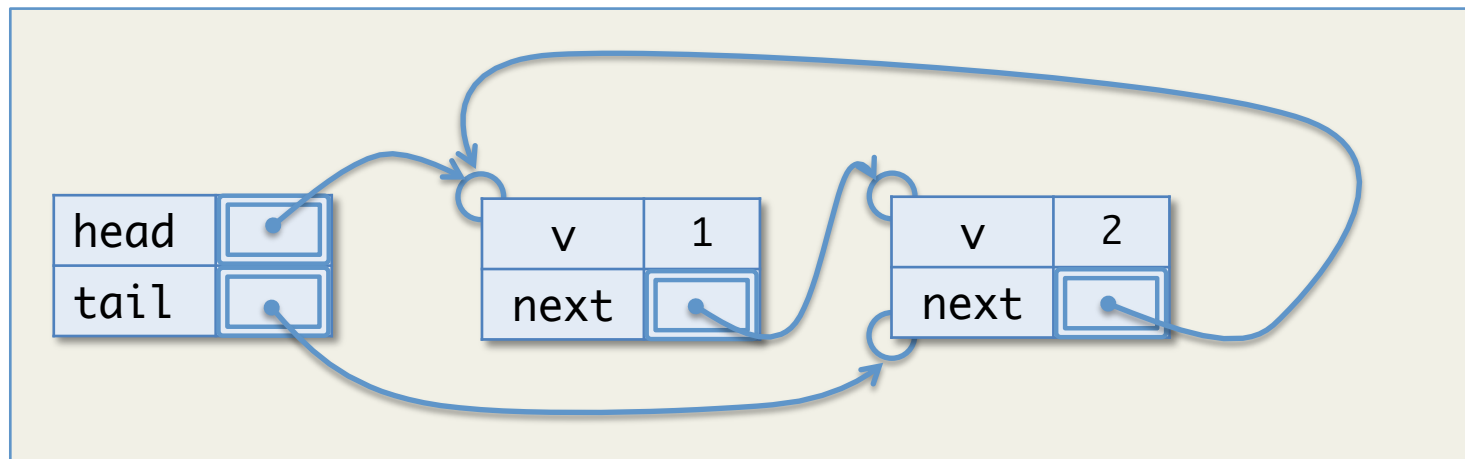
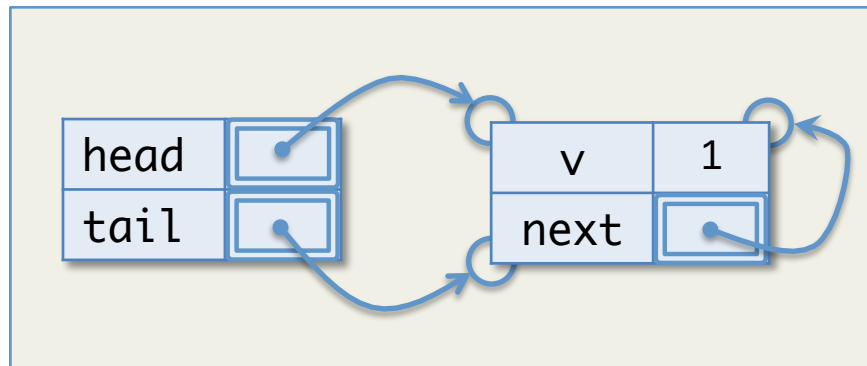
```
type 'a qnode = {  
  v: 'a;  
  mutable next : 'a qnode option  
}  
  
type 'a queue = { mutable head : 'a qnode option;  
                  mutable tail : 'a qnode option }
```

1. yes
2. no
3. not sure

Answer: 1



# Cyclic queues



(And infinitely many more...)

# Linked Queue Invariants

- Just as we imposed some restrictions on which trees count as legitimate Binary Search Trees, Linked Queues must also satisfy representation *invariants*:

Either:

(1) `head` and `tail` are both `None` (i.e. the queue is empty)

or

(2) `head` is `Some n1`, `tail` is `Some n2` and

- `n2` is reachable from `n1` by following 'next' pointers
- `n2.next` is `None`

- We can check that these properties rule out all of the “bogus” examples.
- Each queue operation may assume that these invariants hold of its inputs, and must ensure that the invariants hold when it's done.

Either:

(1) `head` and `tail` are both `None` (i.e. the queue is empty)

or

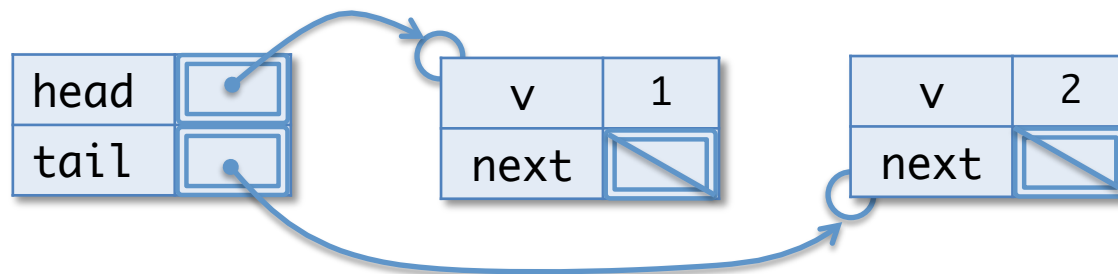
(2) `head` is `Some n1`, `tail` is `Some n2` and

- `n2` is reachable from `n1` by following 'next' pointers
- `n2.next` is `None`

Is this a valid queue?

1. Yes

2. No



Either:

(1) `head` and `tail` are both `None` (i.e. the queue is empty)

or

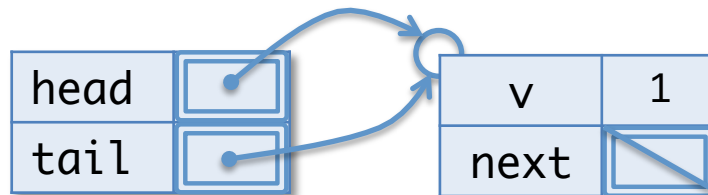
(2) `head` is `Some n1`, `tail` is `Some n2` and

- `n2` is reachable from `n1` by following 'next' pointers
- `n2.next` is `None`

Is this a valid queue?

1. Yes

2. No



# Implementing Linked Queues

LinkedQ.ml

# create and is\_empty

```
(* create an empty queue *)  
let create () : 'a queue =  
  { head = None;  
    tail = None }
```

```
(* determine whether a queue is empty *)  
let is_empty (q: 'a queue) : bool =  
  q.head = None
```

- `create` *establishes* the queue invariants
  - both head and tail are None
- `is_empty` *assumes* the queue invariants
  - it doesn't have to check that q.tail is None

# enq

```
(* add an element to the tail of a queue *)
let enq (x: 'a) (q: 'a queue) : unit =
  let newnode = {v=x; next=None} in
  begin match q.tail with
  | None ->
      q.head <- Some newnode;
      q.tail <- Some newnode
  | Some n ->
      n.next <- Some newnode;
      q.tail <- Some newnode
  end
```

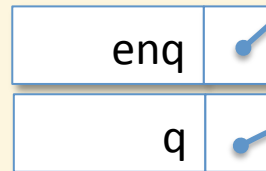
- The code for `enq` is informed by the queue invariant:
  - either the queue is empty, and we just update head and tail, or
  - the queue is non-empty, in which case we have to “patch up” the “next” link of the old tail node to maintain the queue invariant.

# Calling Enq on a non-empty queue

Workspace

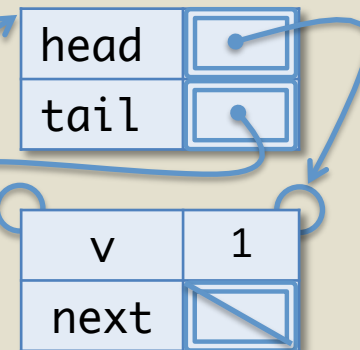
enq 2 q

Stack



Heap

```
fun (x: 'a) (q: 'a queue) ->  
  let newnode = {v=x; next=None}  
  in begin match q.tail with  
    | None -> ...  
    | Some n -> ...  
  end
```

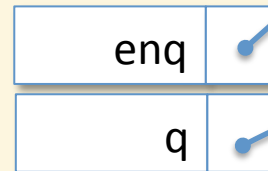


# Calling Enq on a non-empty queue

Workspace

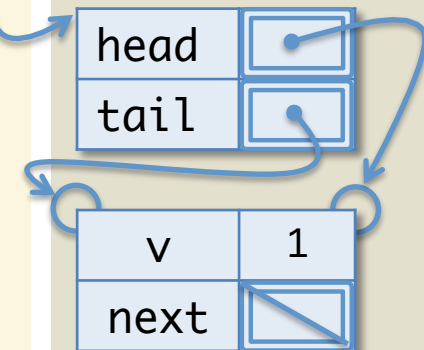
enq 2 q

Stack

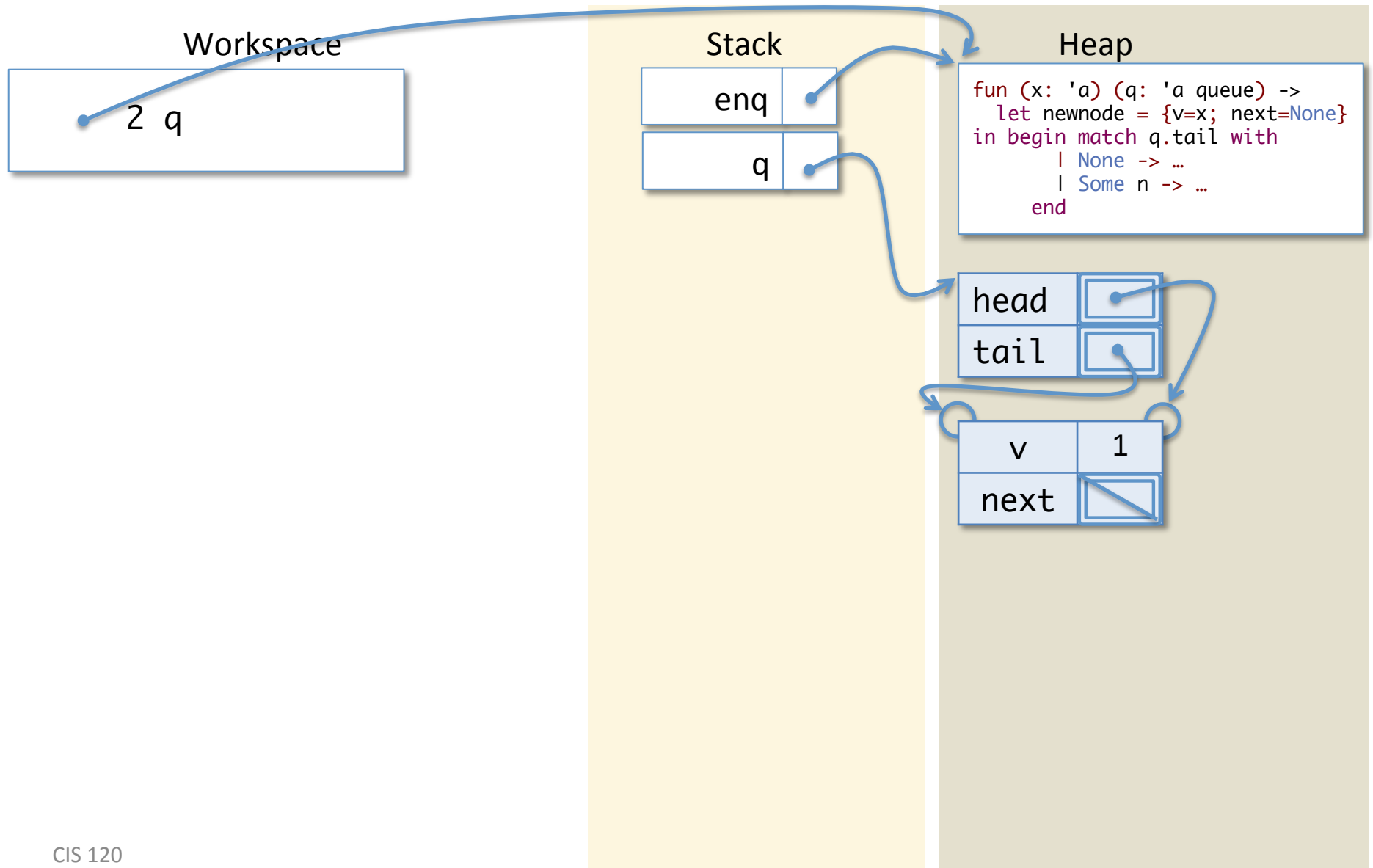


Heap

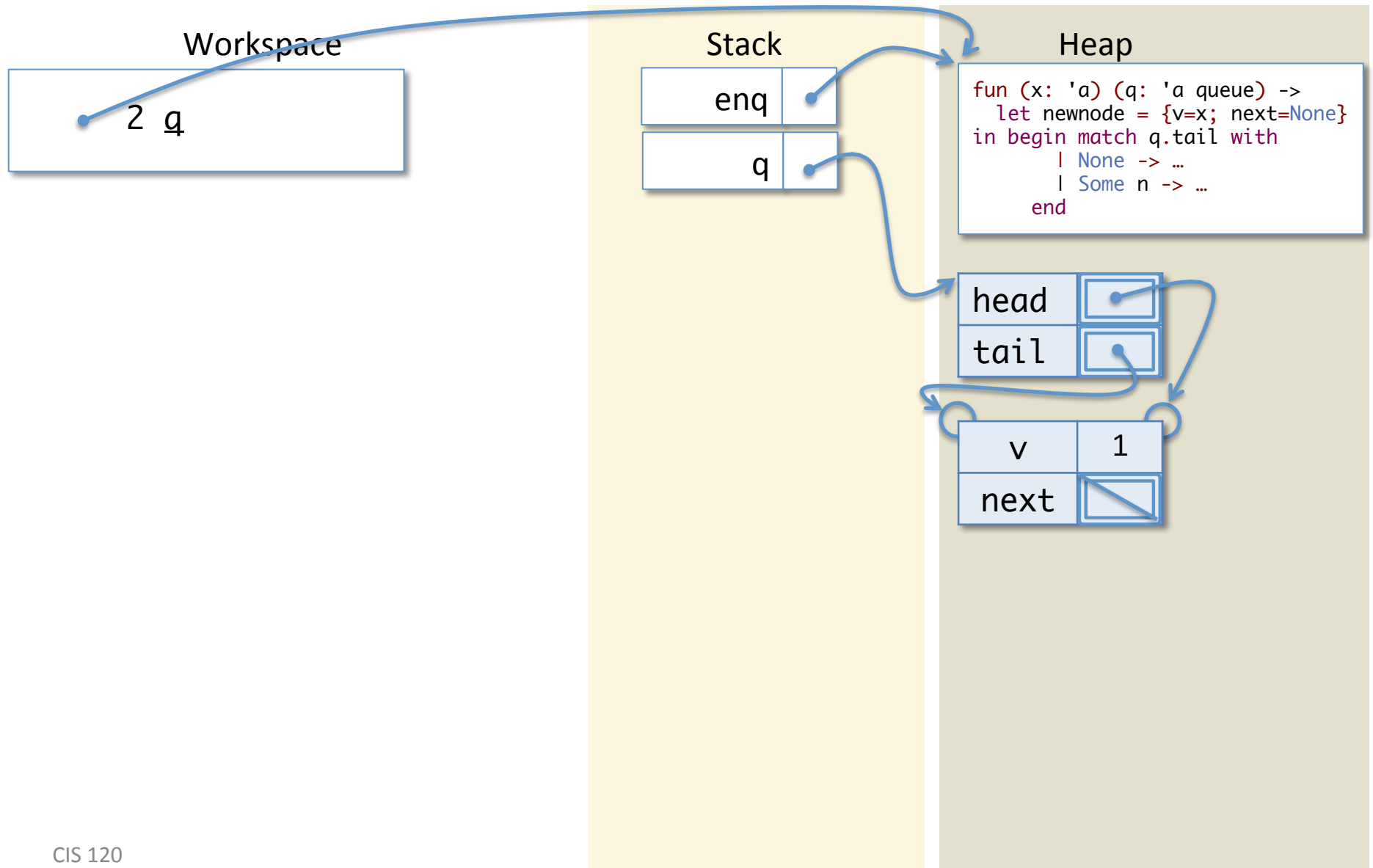
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fun (x: 'a) (q: 'a queue) ->  
  let newnode = {v=x; next=None}  
  in begin match q.tail with  
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    | Some n -> ...  
  end
```



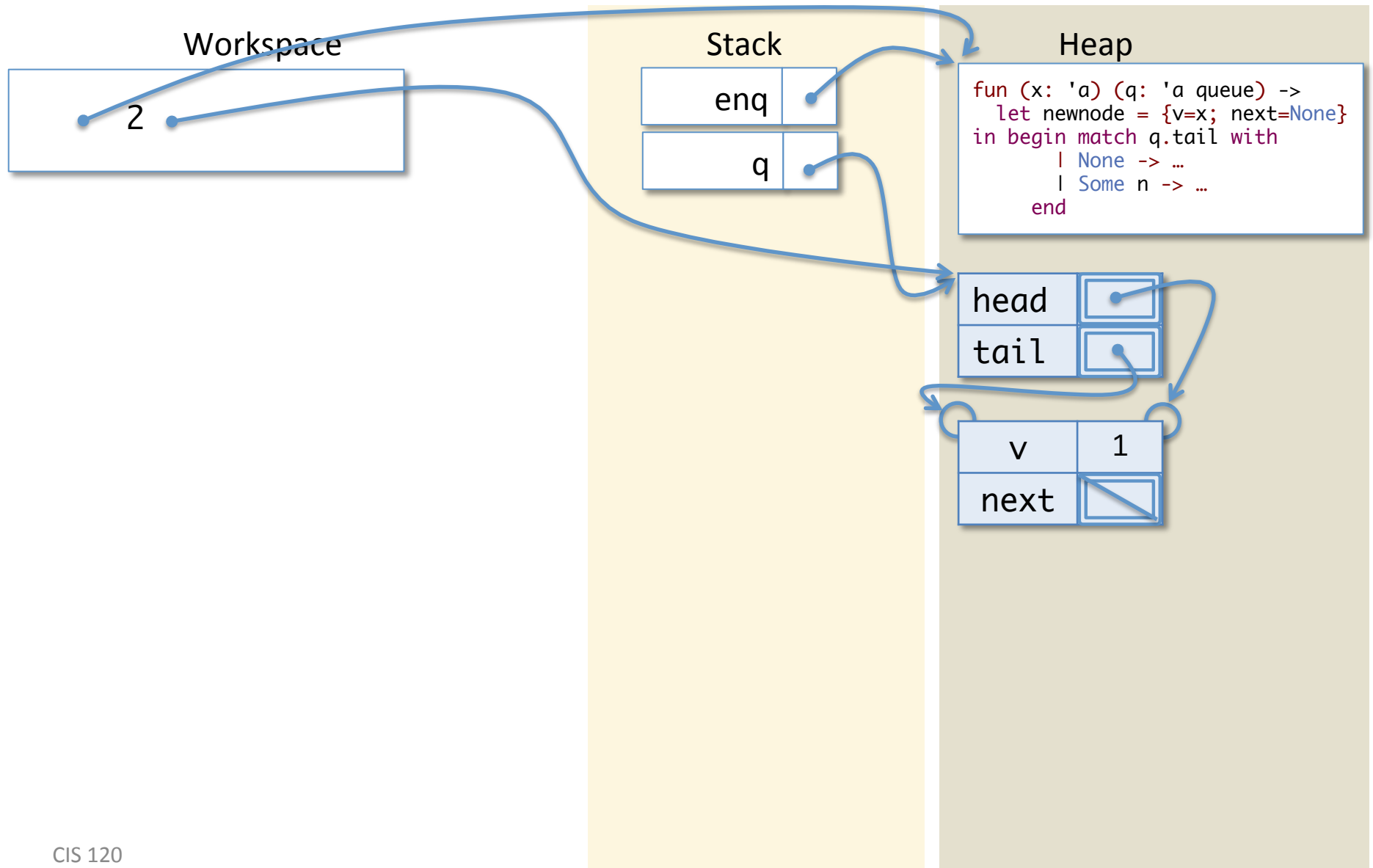
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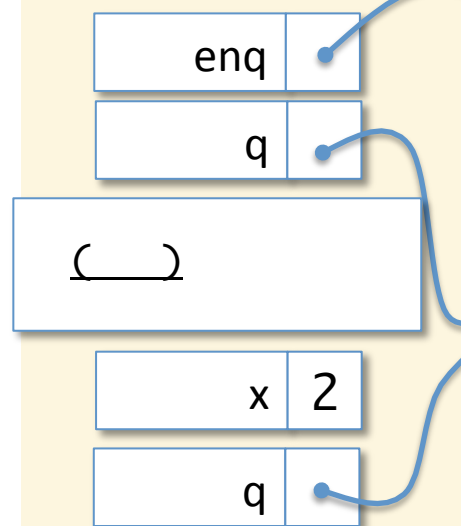


# Calling Enq on a non-empty queue

## Workspace

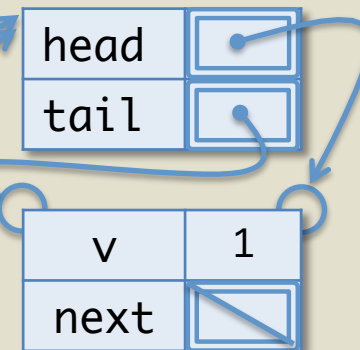
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let newnode = {v=x; next=None} in
begin match q.tail with
| None ->
  q.head <- Some newnode;
  q.tail <- Some newnode
| Some n ->
  n.next <- Some newnode;
  q.tail <- Some newnode
end
```

## Stack



## Heap

```
fun (x: 'a) (q: 'a queue) ->
  let newnode = {v=x; next=None} in
  begin match q.tail with
  | None -> ...
  | Some n -> ...
  end
```

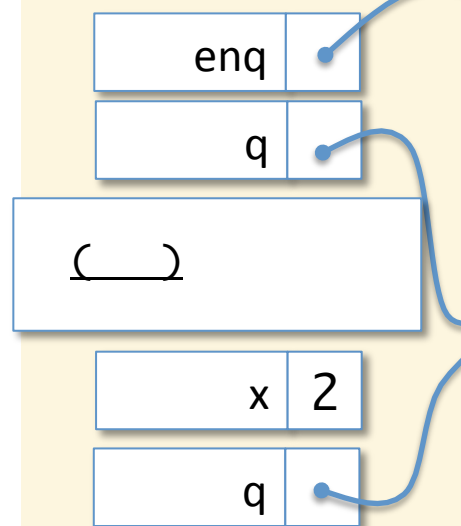


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## Workspace

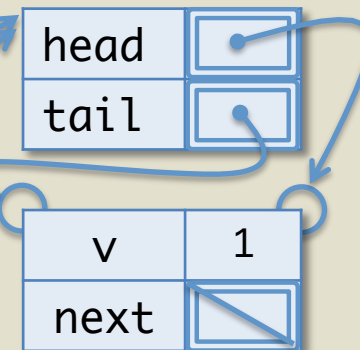
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let newnode = {v=x; next=None} in
begin match q.tail with
| None ->
  q.head <- Some newnode;
  q.tail <- Some newnode
| Some n ->
  n.next <- Some newnode;
  q.tail <- Some newnode
end
```

## Stack



## Heap

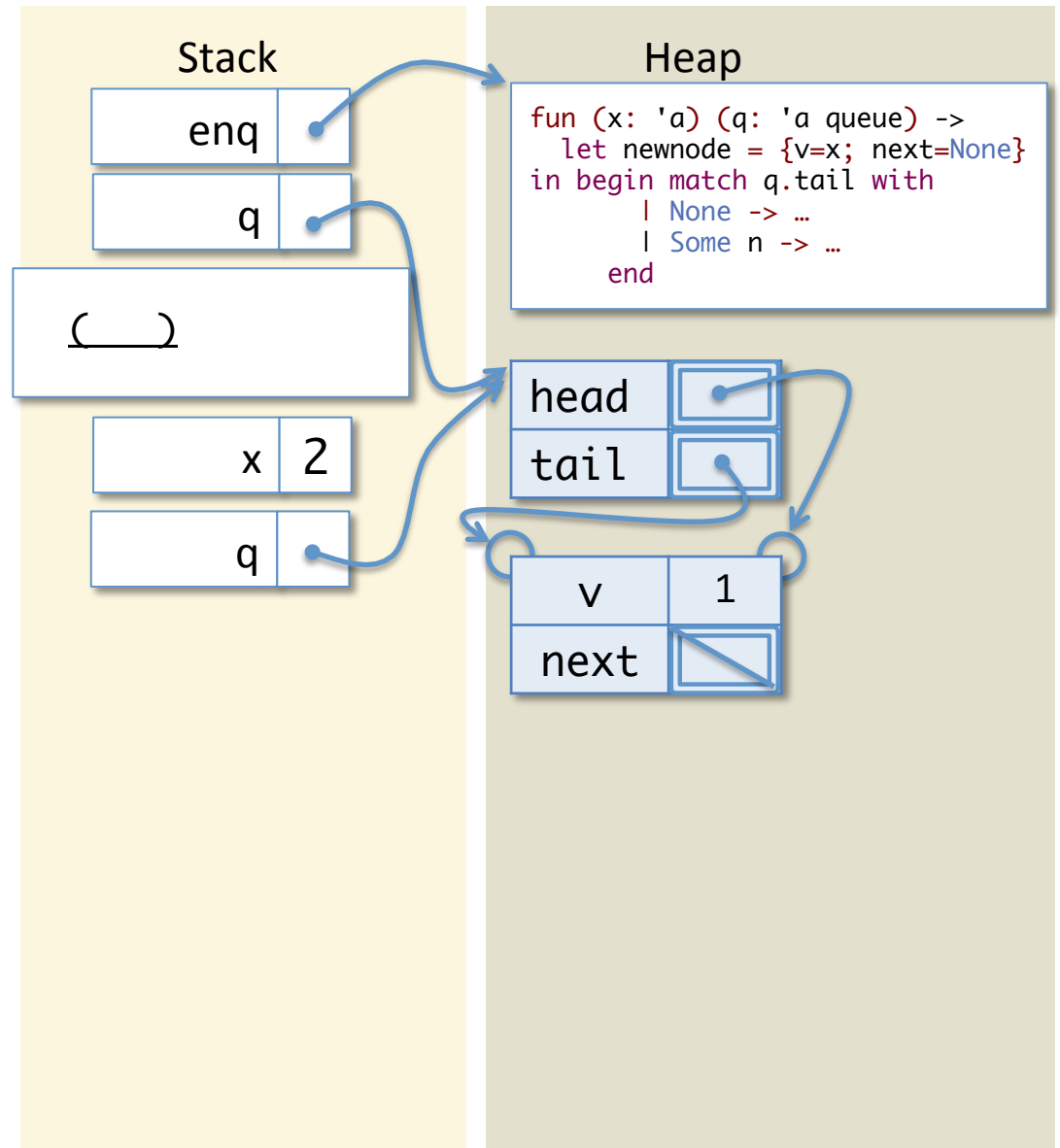
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fun (x: 'a) (q: 'a queue) ->
  let newnode = {v=x; next=None} in
  begin match q.tail with
  | None -> ...
  | Some n -> ...
  end
```



# Calling Enq on a non-empty queue

## Workspace

```
let newnode = {v=2; next=None} in
begin match q.tail with
| None ->
  q.head <- Some newnode;
  q.tail <- Some newnode
| Some n ->
  n.next <- Some newnode;
  q.tail <- Some newnode
end
```

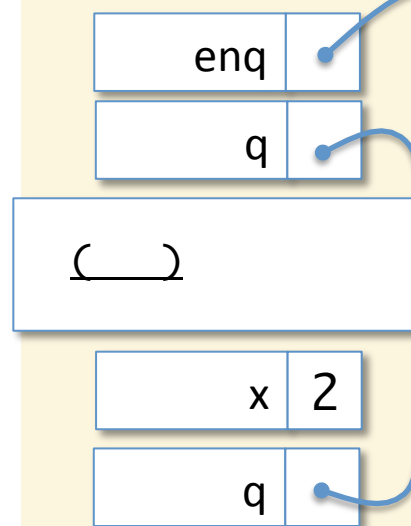


# Calling Enq on a non-empty queue

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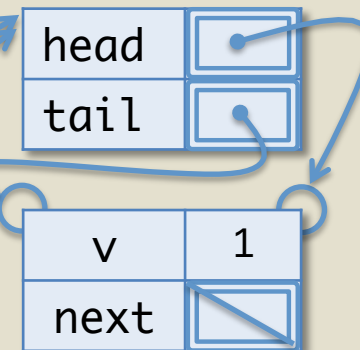
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let newnode = {v=2; next=None} in
begin match q.tail with
| None ->
  q.head <- Some newnode;
  q.tail <- Some newnode
| Some n ->
  n.next <- Some newnode;
  q.tail <- Some newnode
end
```

## Stack

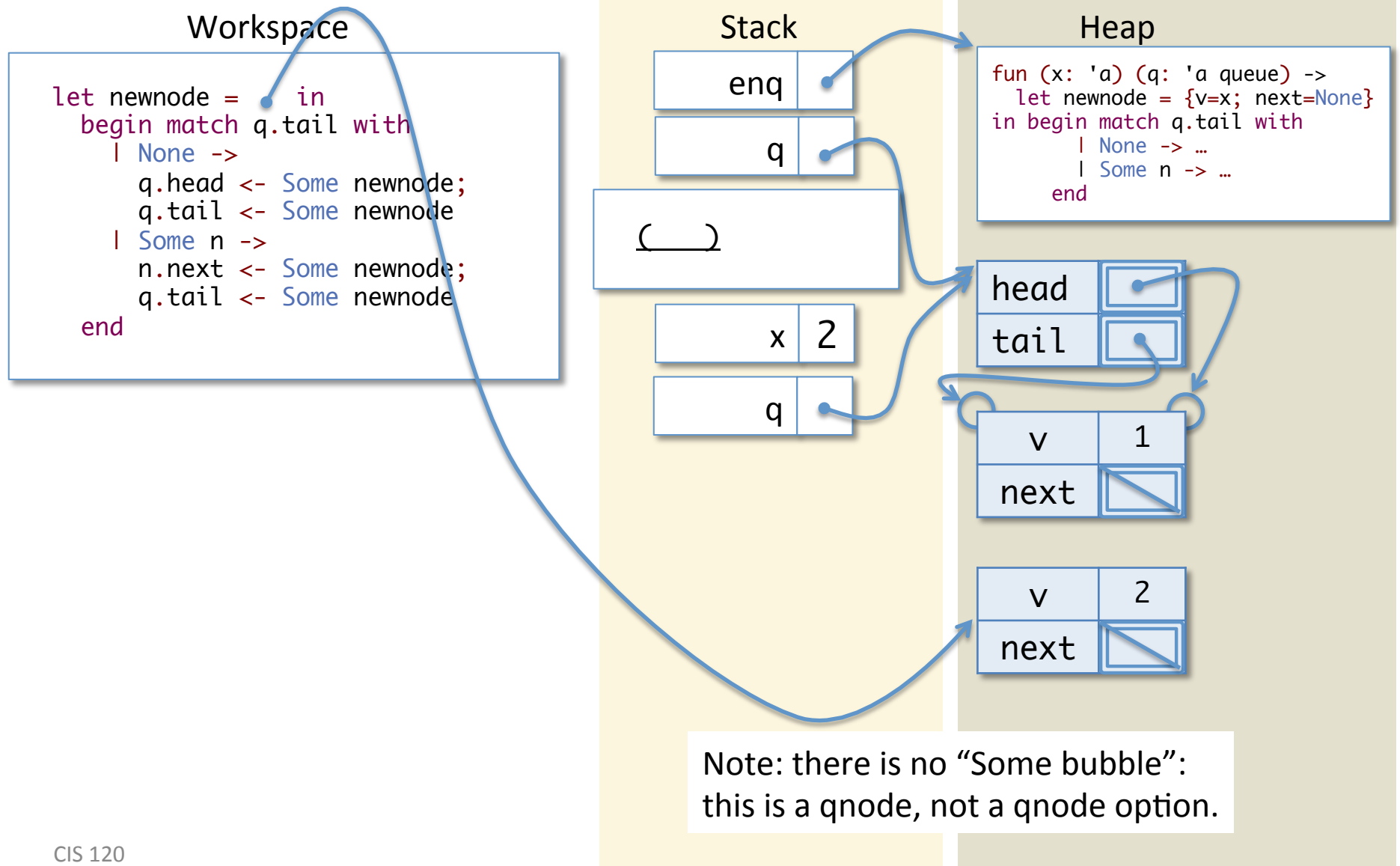


## Heap

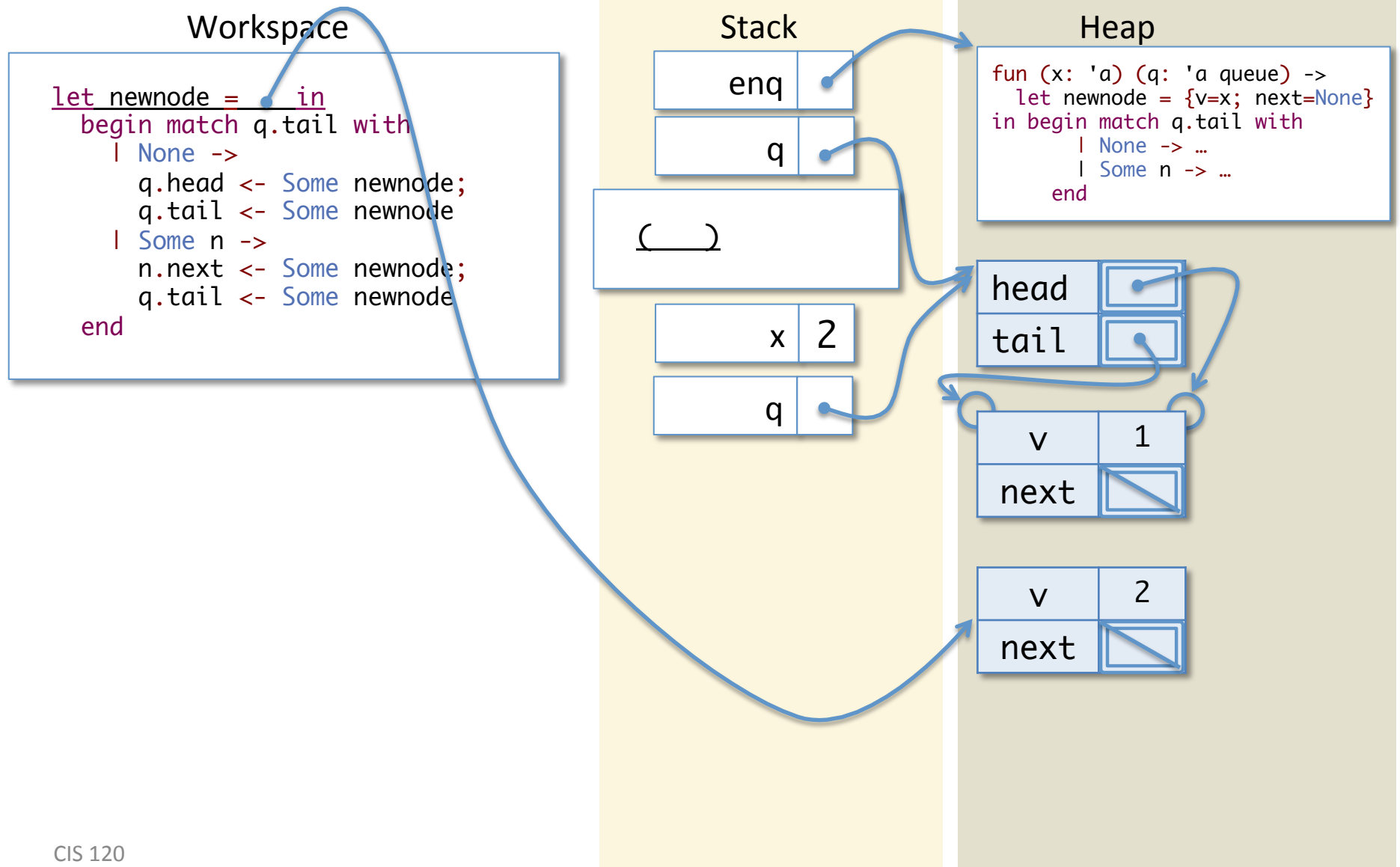
```
fun (x: 'a) (q: 'a queue) ->
  let newnode = {v=x; next=None}
  in begin match q.tail with
    | None -> ...
    | Some n -> ...
  end
```



# Calling Enq on a non-empty queue



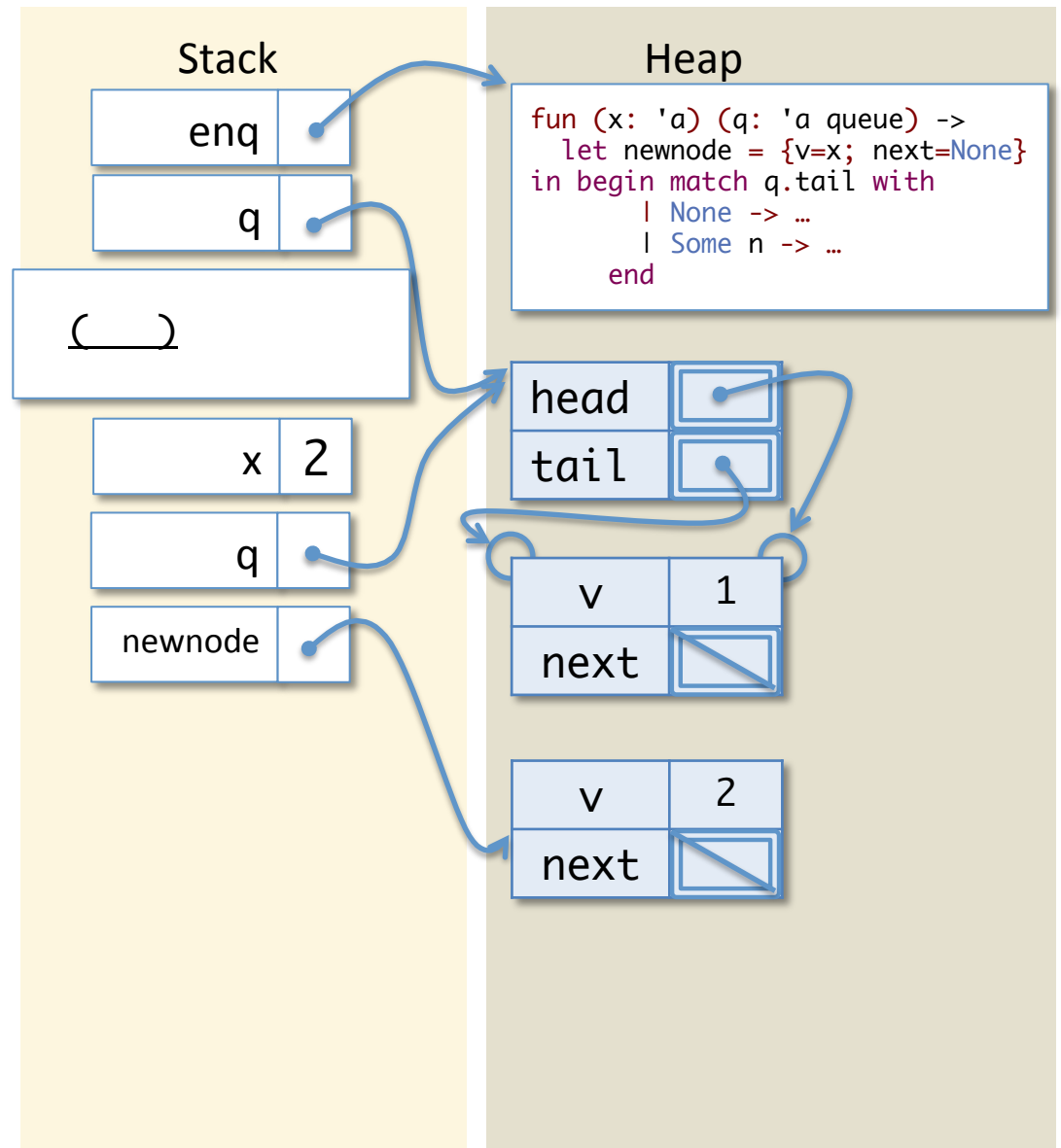
# Calling Enq on a non-empty queue



# Calling Enq on a non-empty queue

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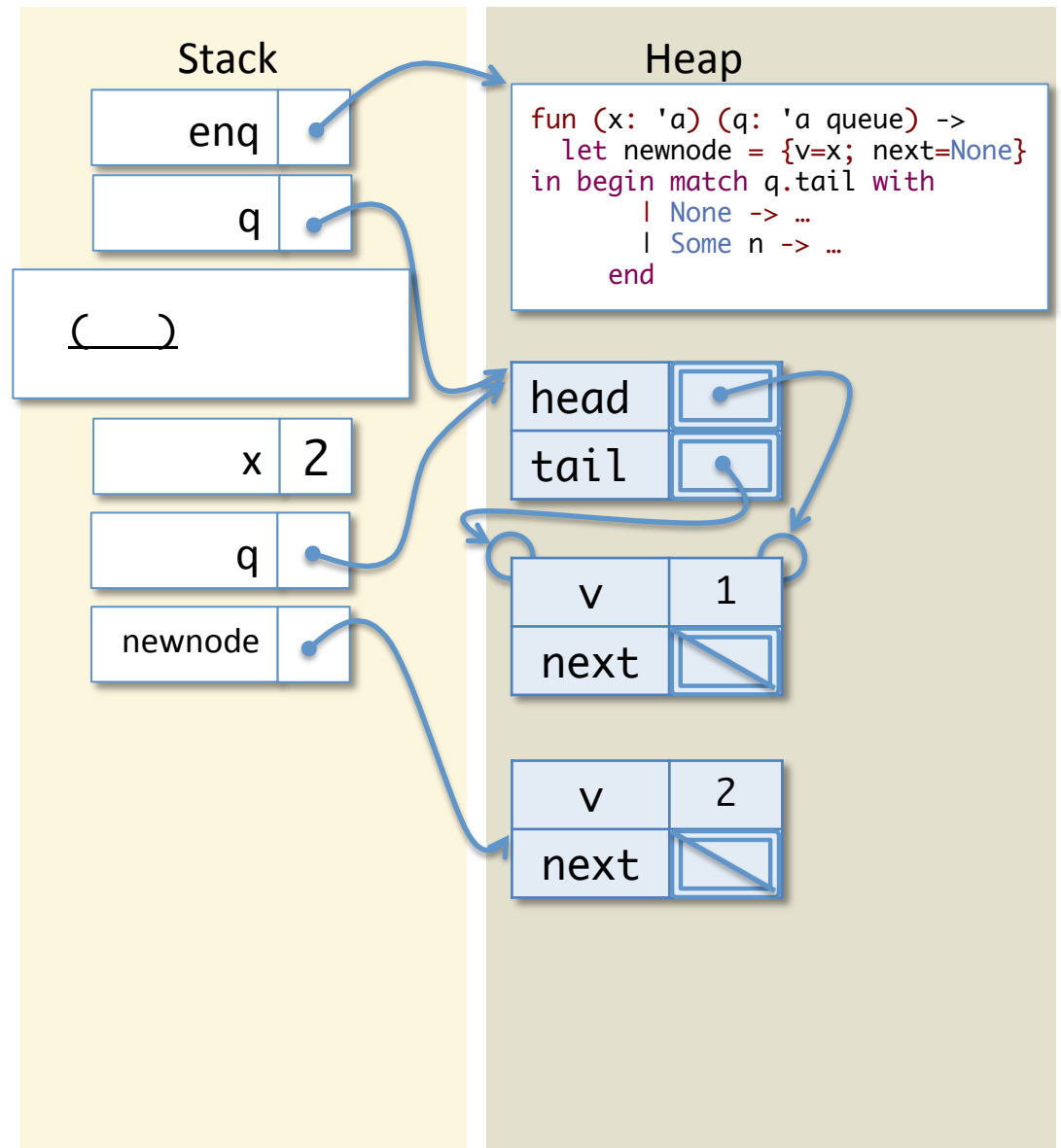
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  q.tail <- Some newnode
end
```



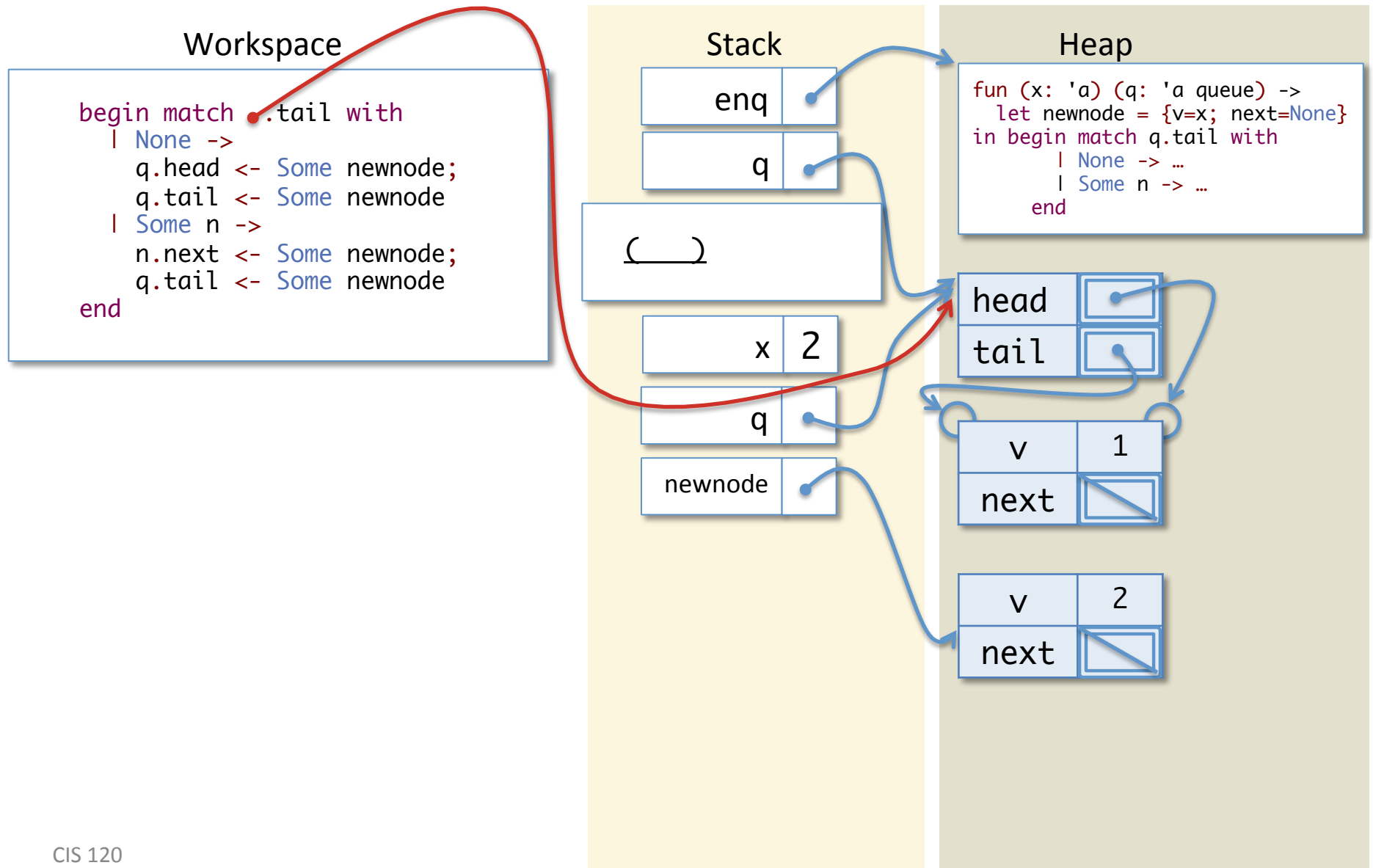
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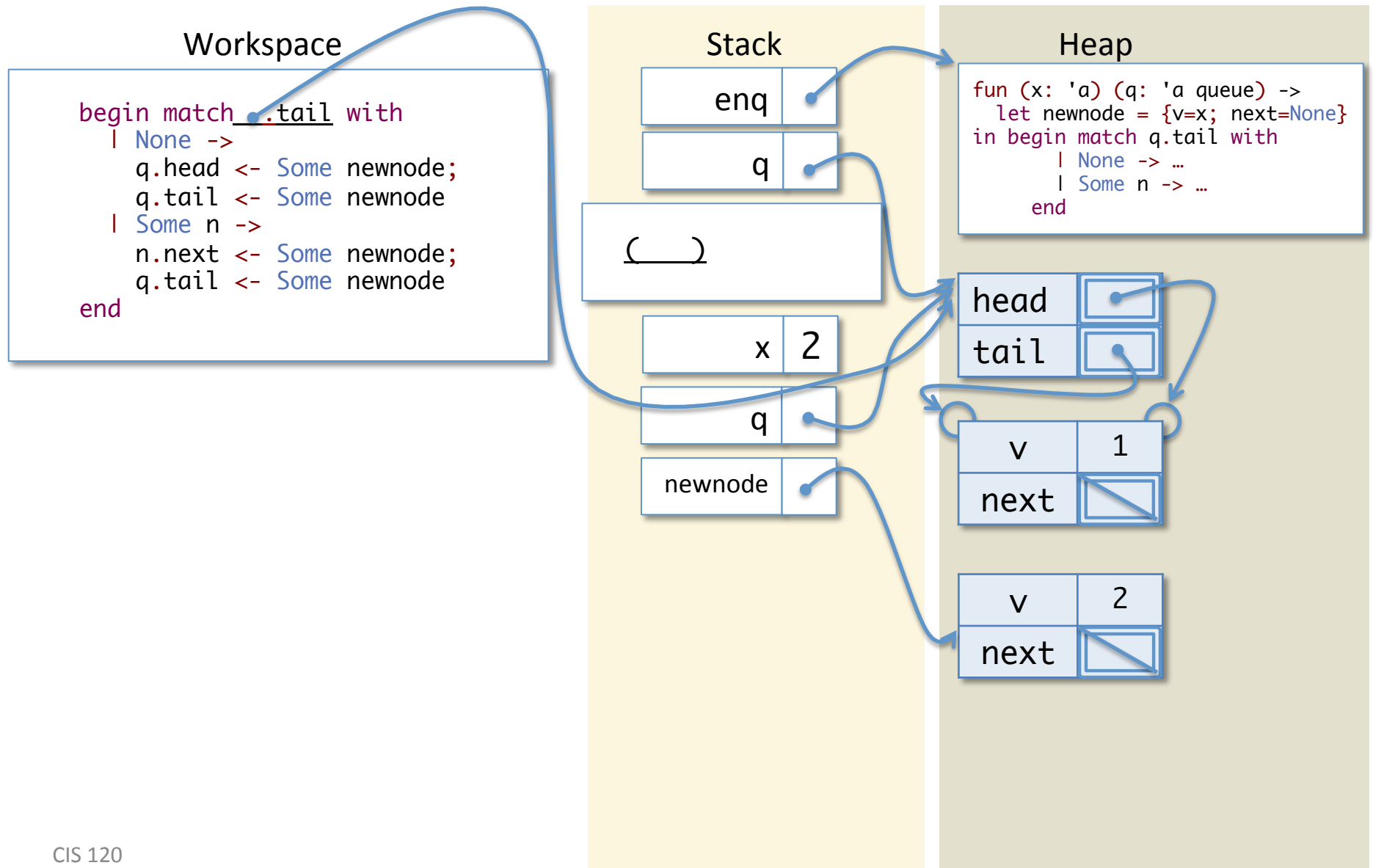
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  q.head <- Some newnode;
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| Some n ->
  n.next <- Some newnode;
  q.tail <- Some newnode
end
```



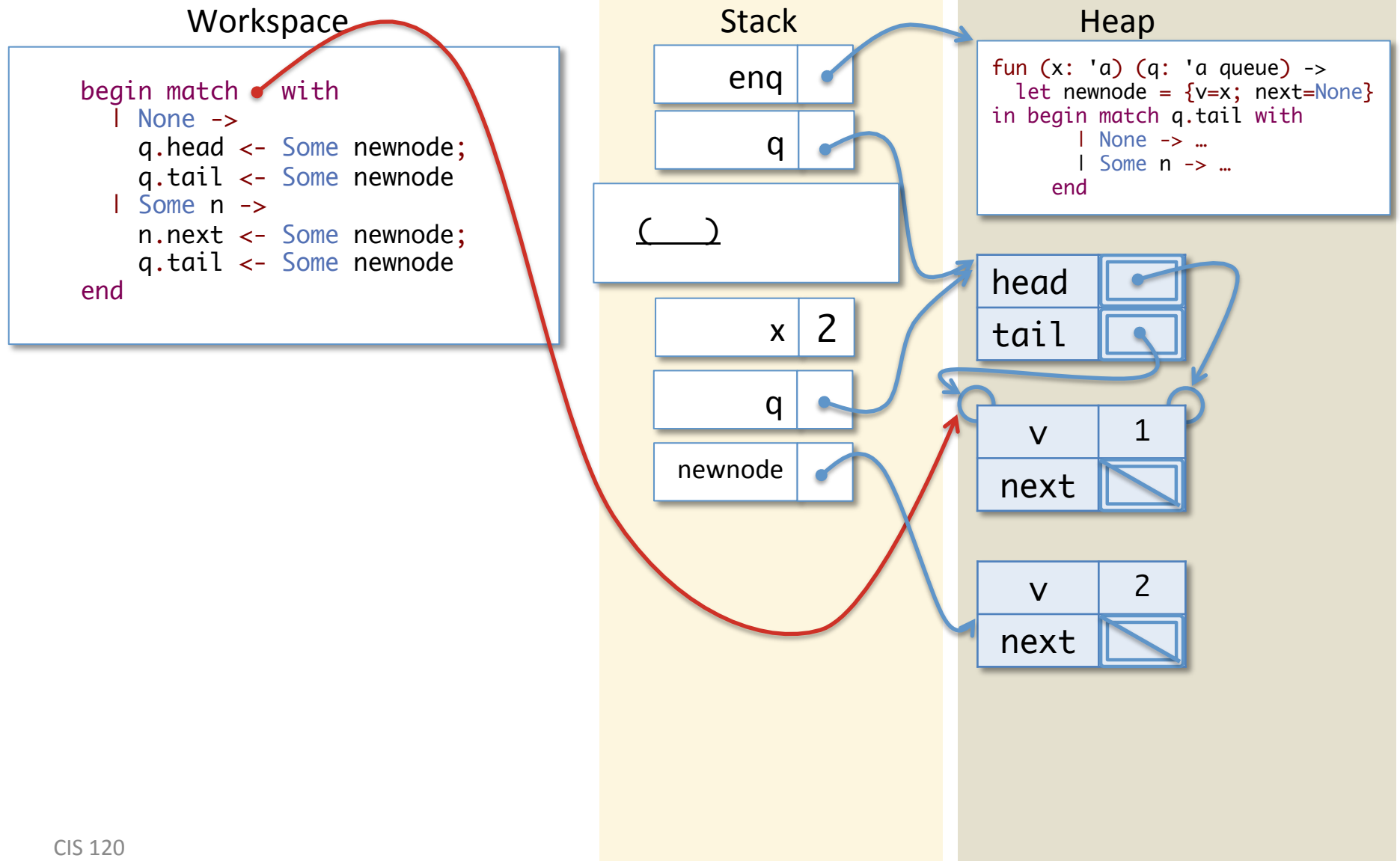
# Calling Enq on a non-empty queue



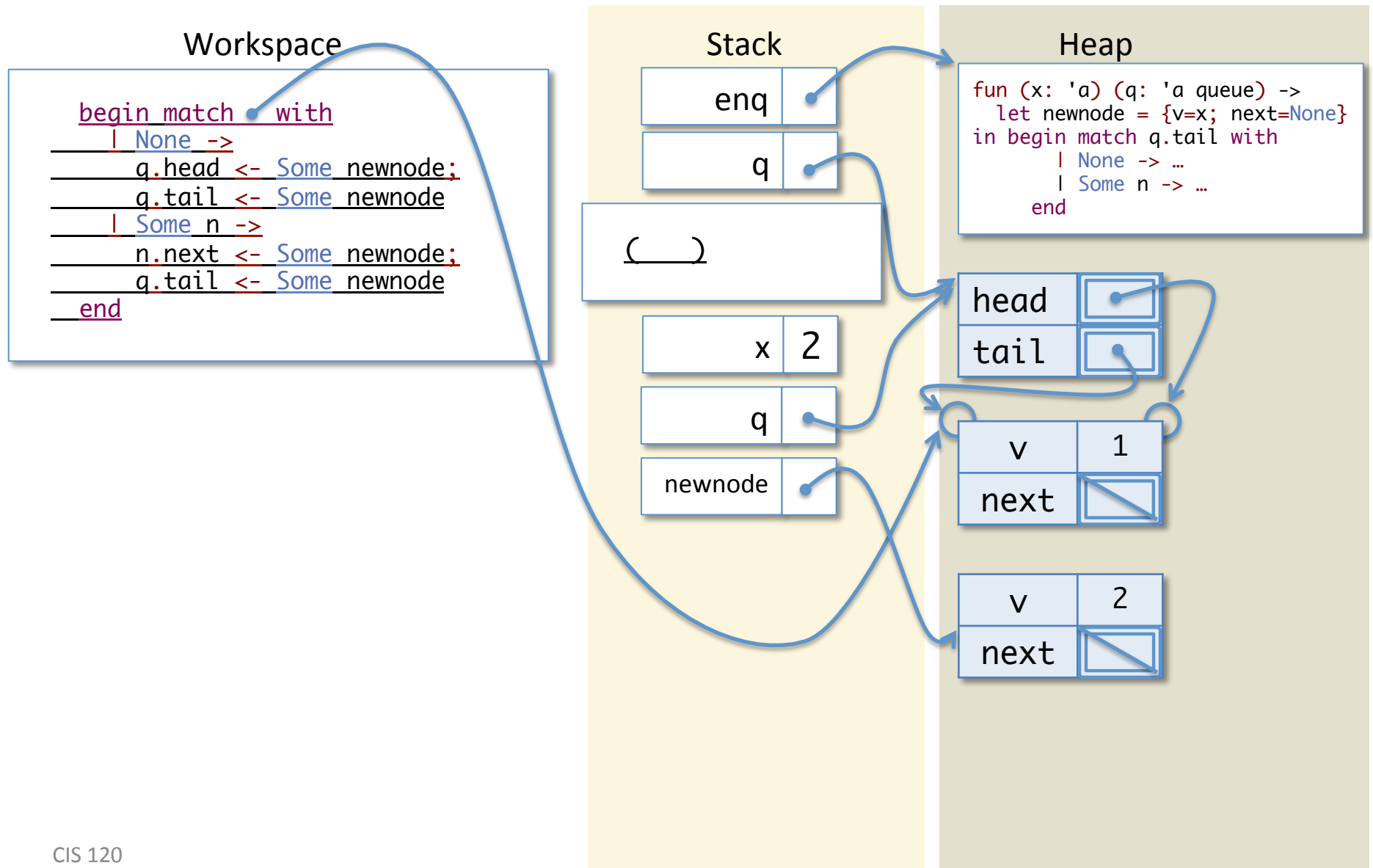
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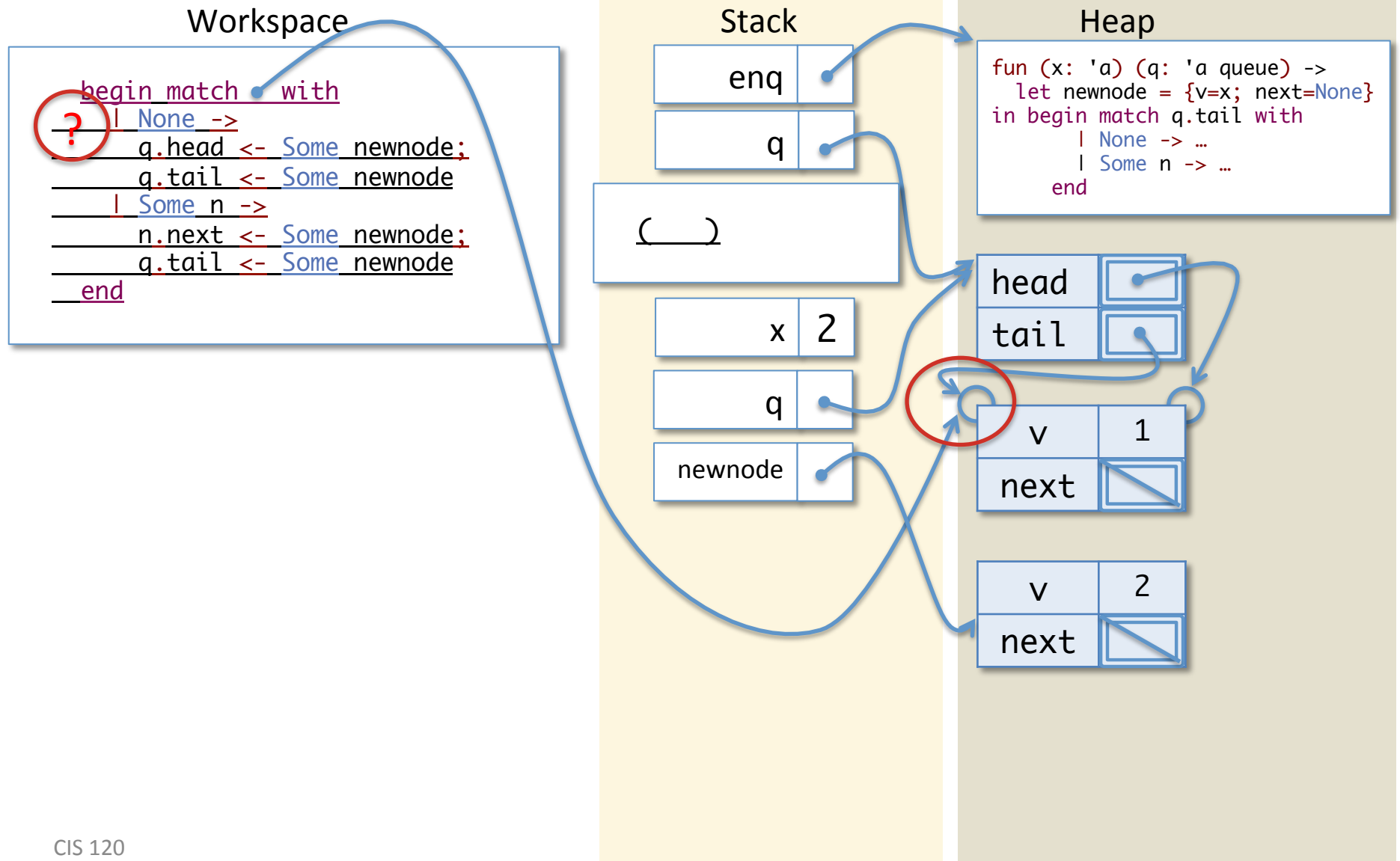
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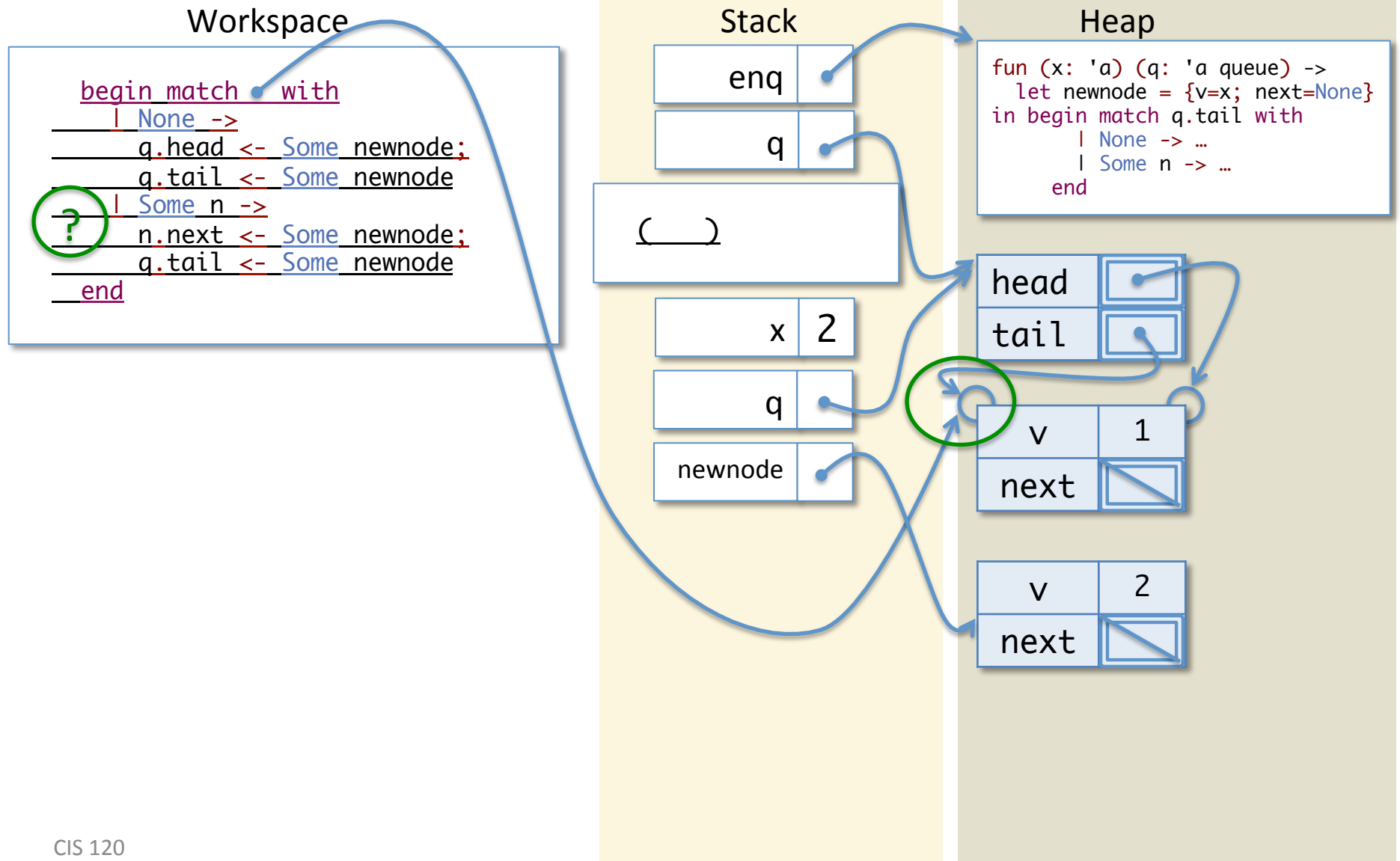
# Calling Enq on a non-empty queue



# Calling Enq on a non-empty queue



# Calling Enq on a non-empty queue

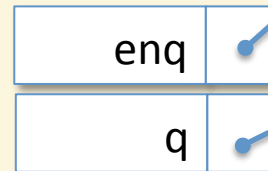


# Calling Enq on a non-empty queue

## Workspace

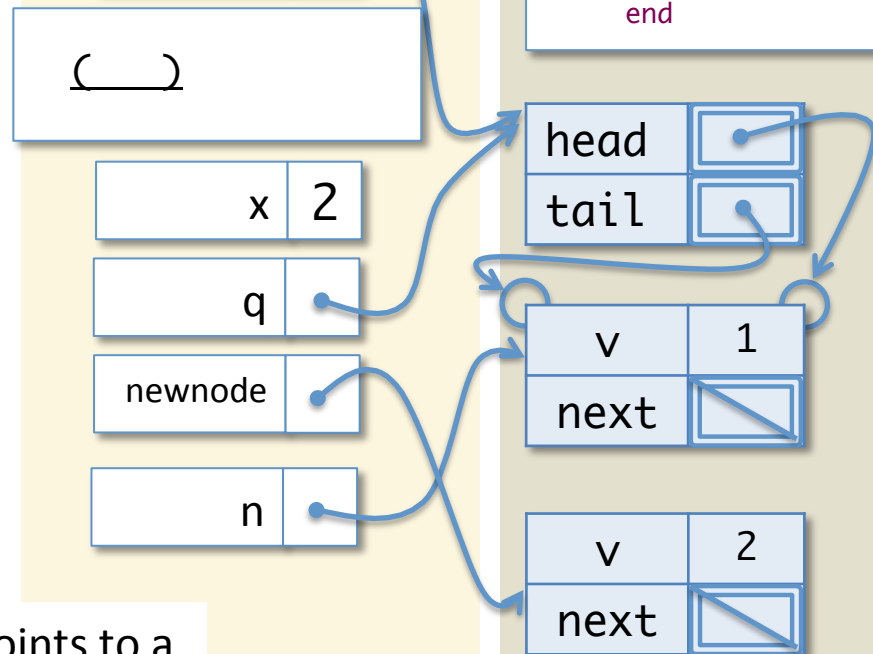
```
n.next <- Some newnode;  
q.tail <- Some newnode
```

## Stack



## Heap

```
fun (x: 'a) (q: 'a queue) ->  
  let newnode = {v=x; next=None}  
  in begin match q.tail with  
    | None -> ...  
    | Some n -> ...  
  end
```

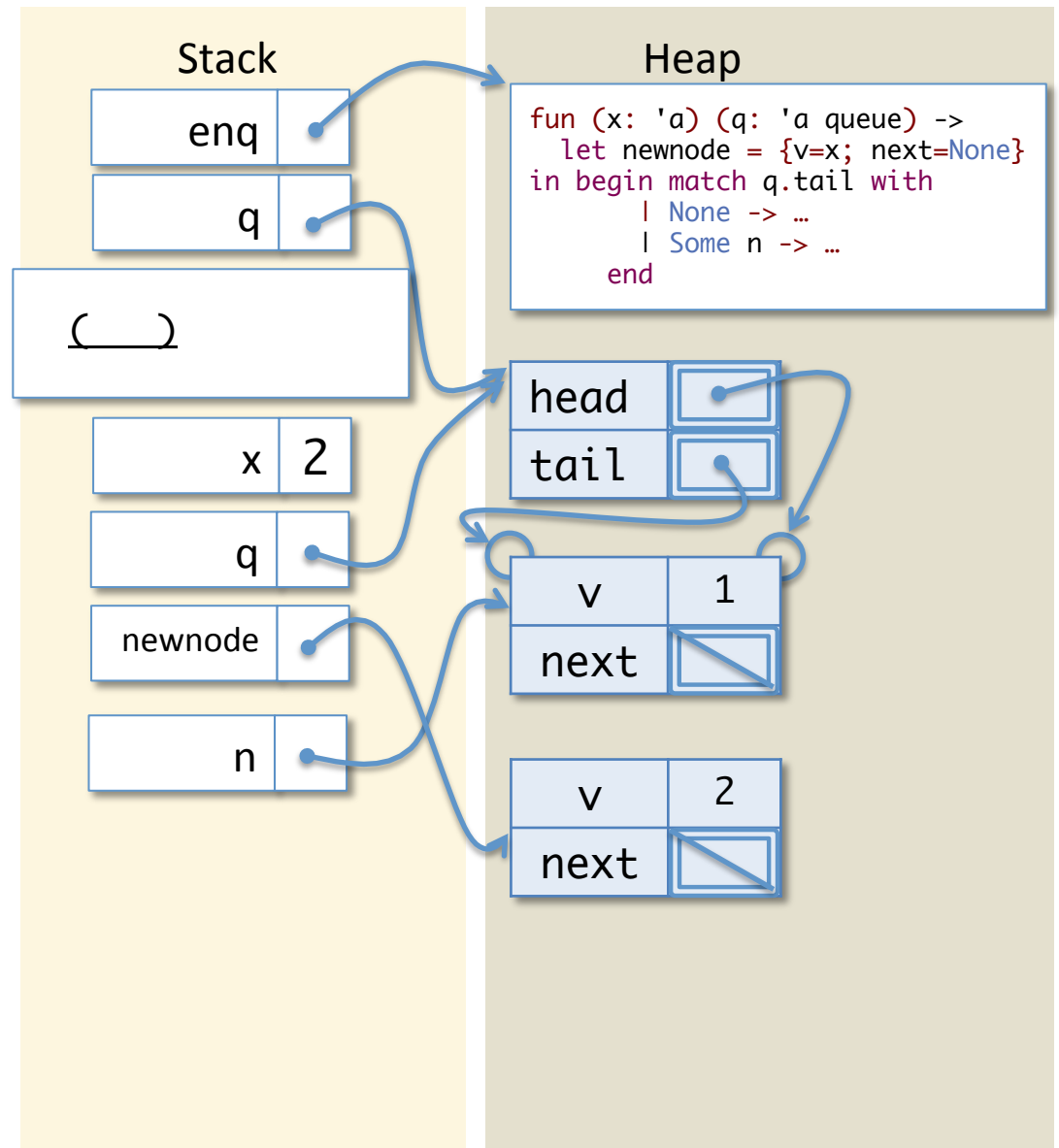


Note: n points to a  
qnode, not a  
qnode option.

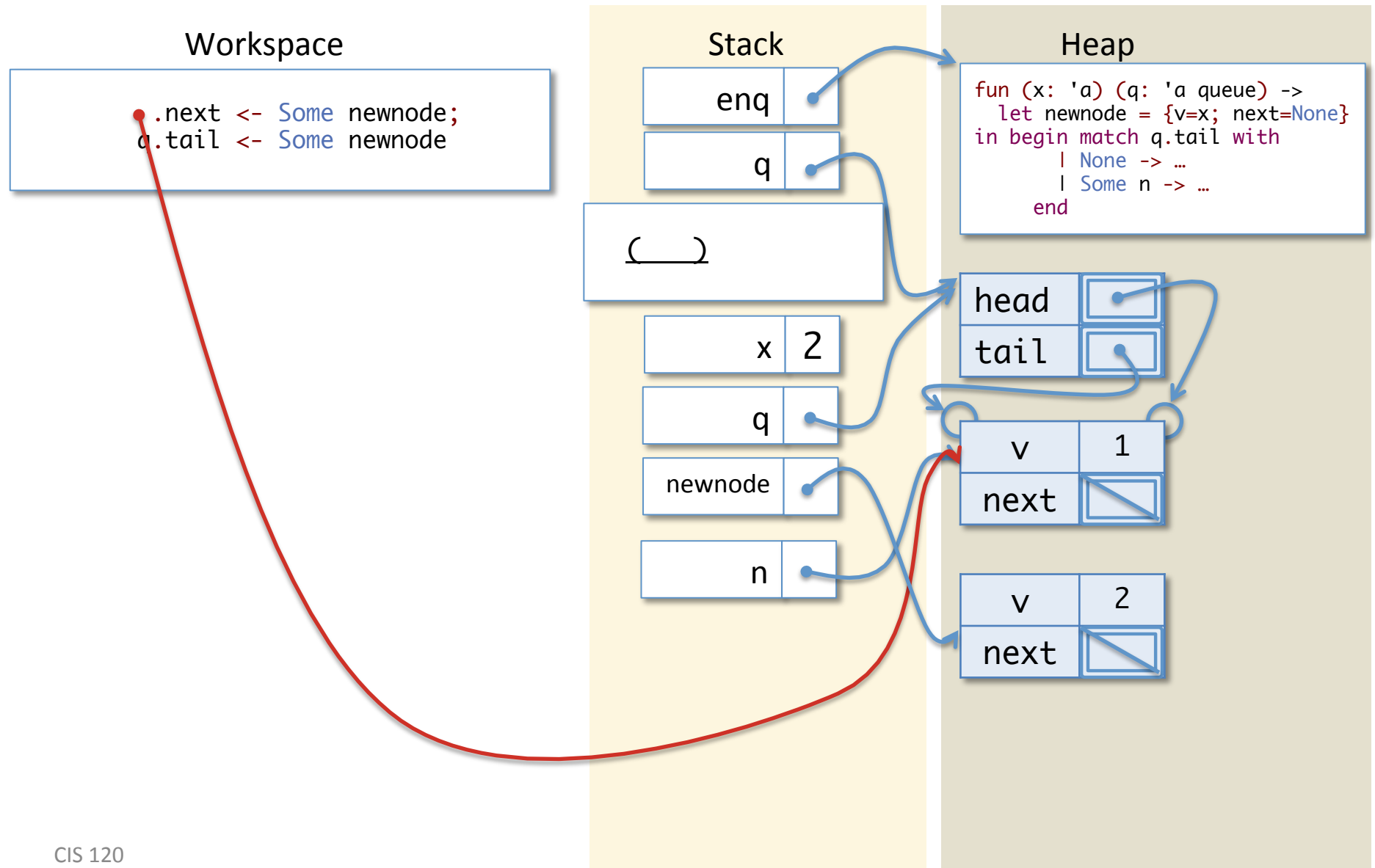
# Calling Enq on a non-empty queue

## Workspace

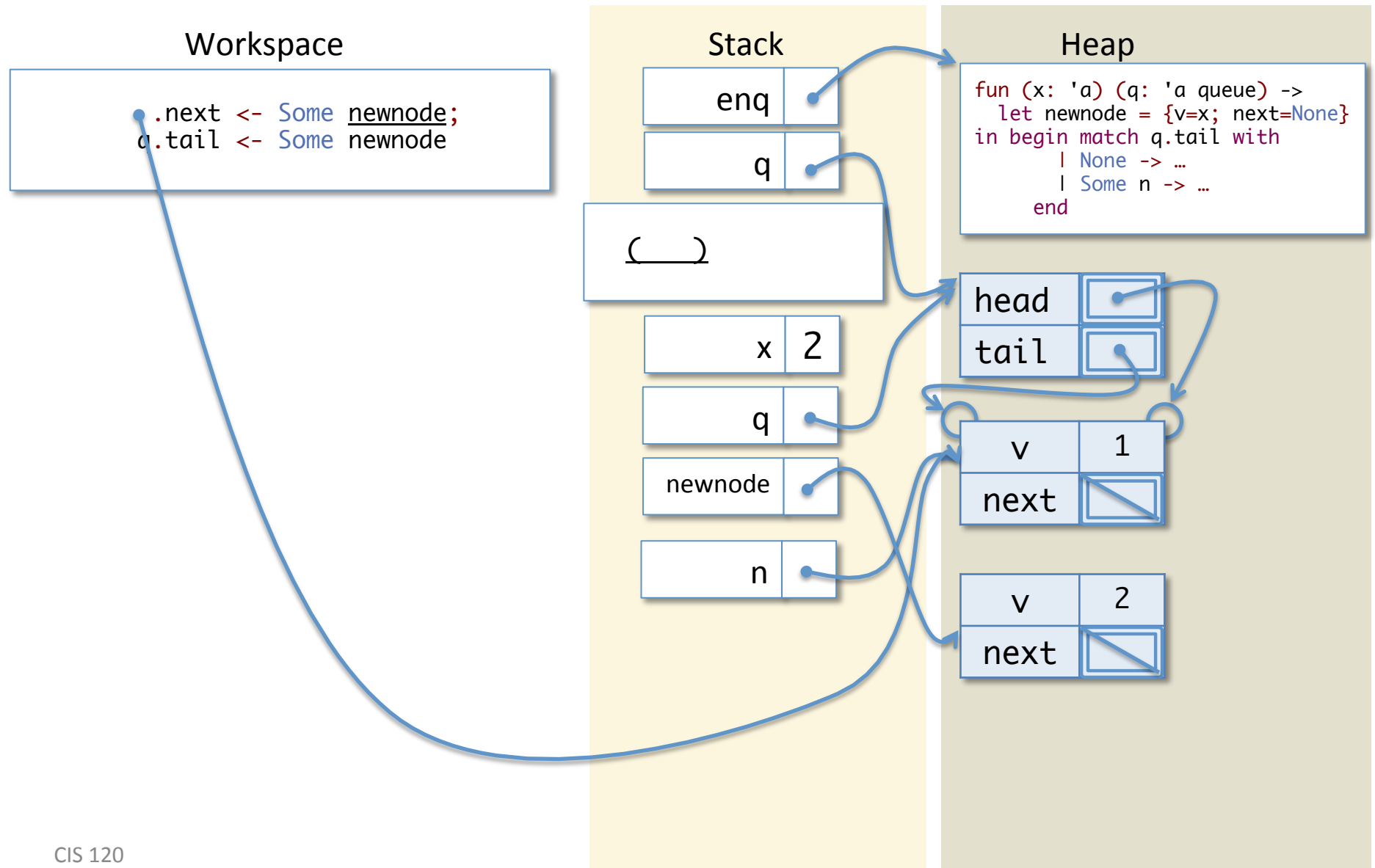
```
n.next <- Some newnode;  
q.tail <- Some newnode
```



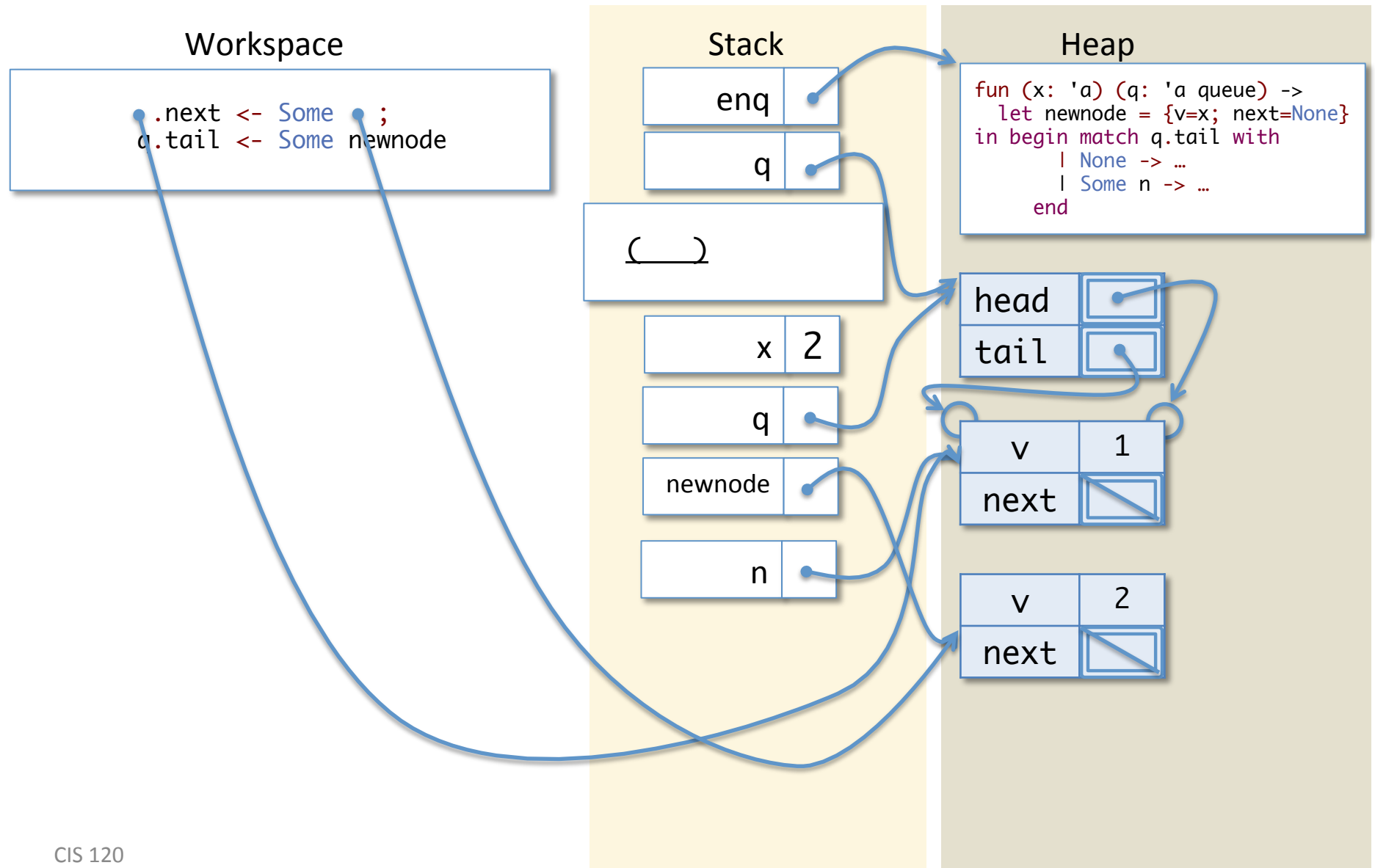
# Calling Enq on a non-empty queue



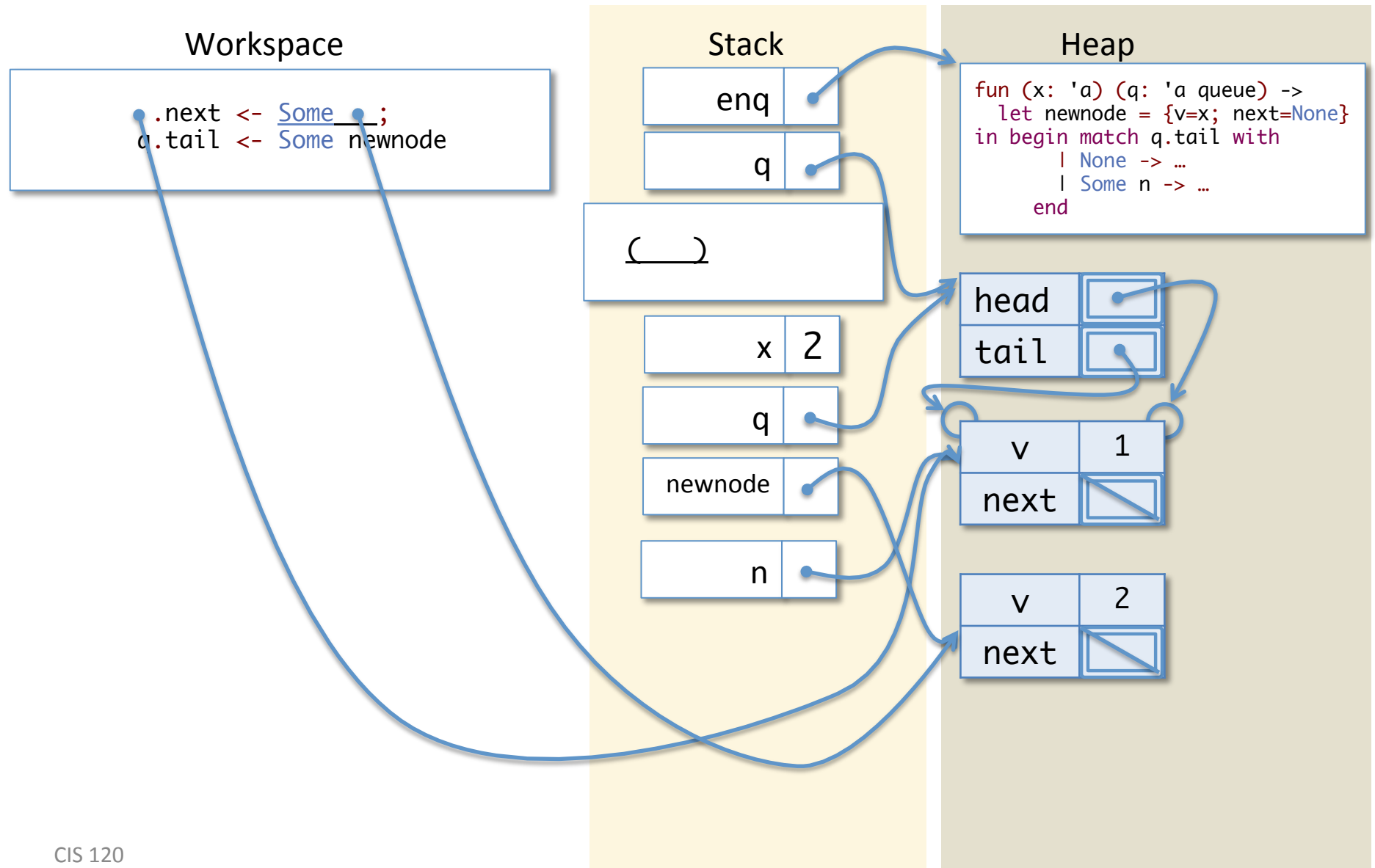
# Calling Enq on a non-empty queue



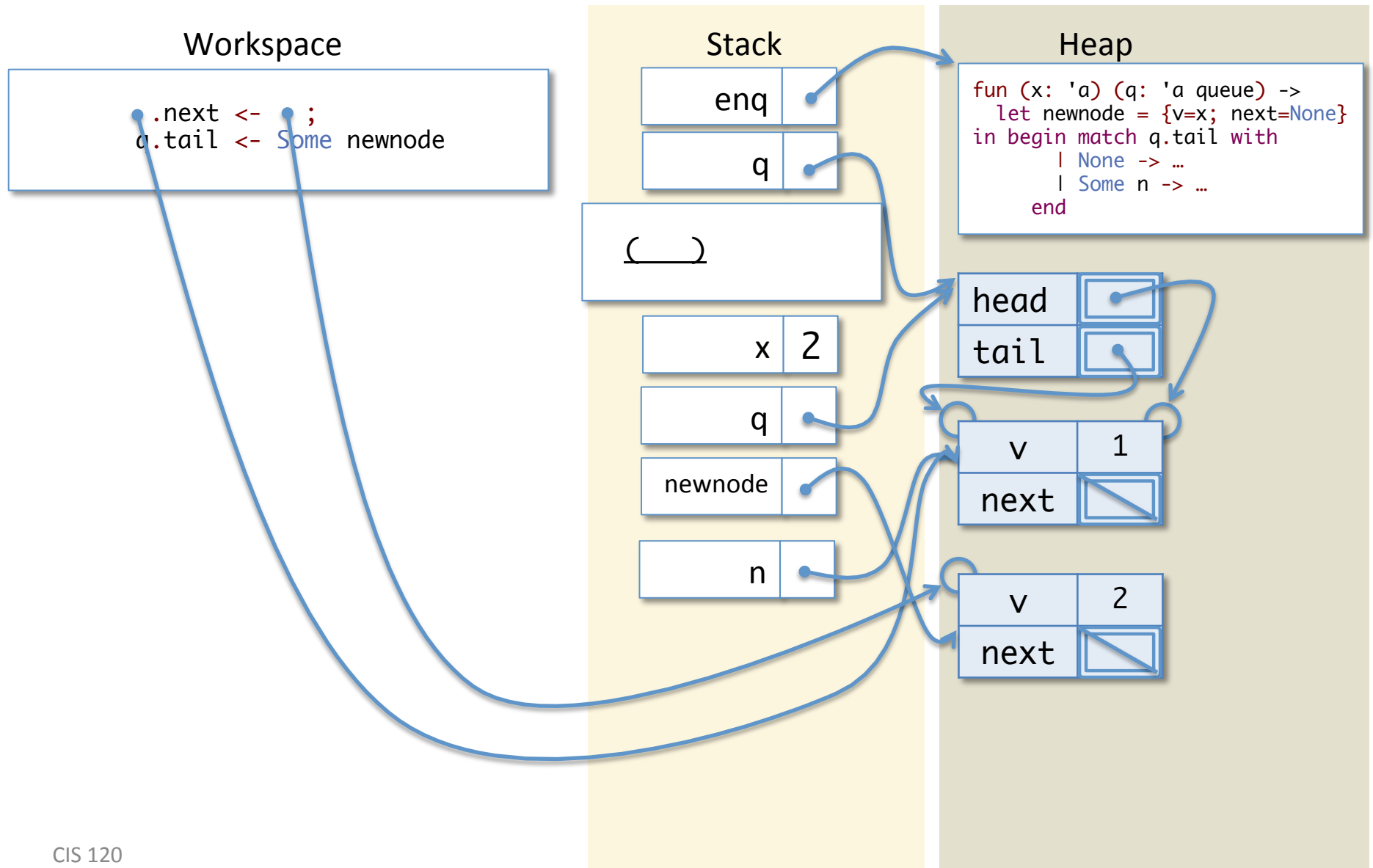
# Calling Enq on a non-empty queue



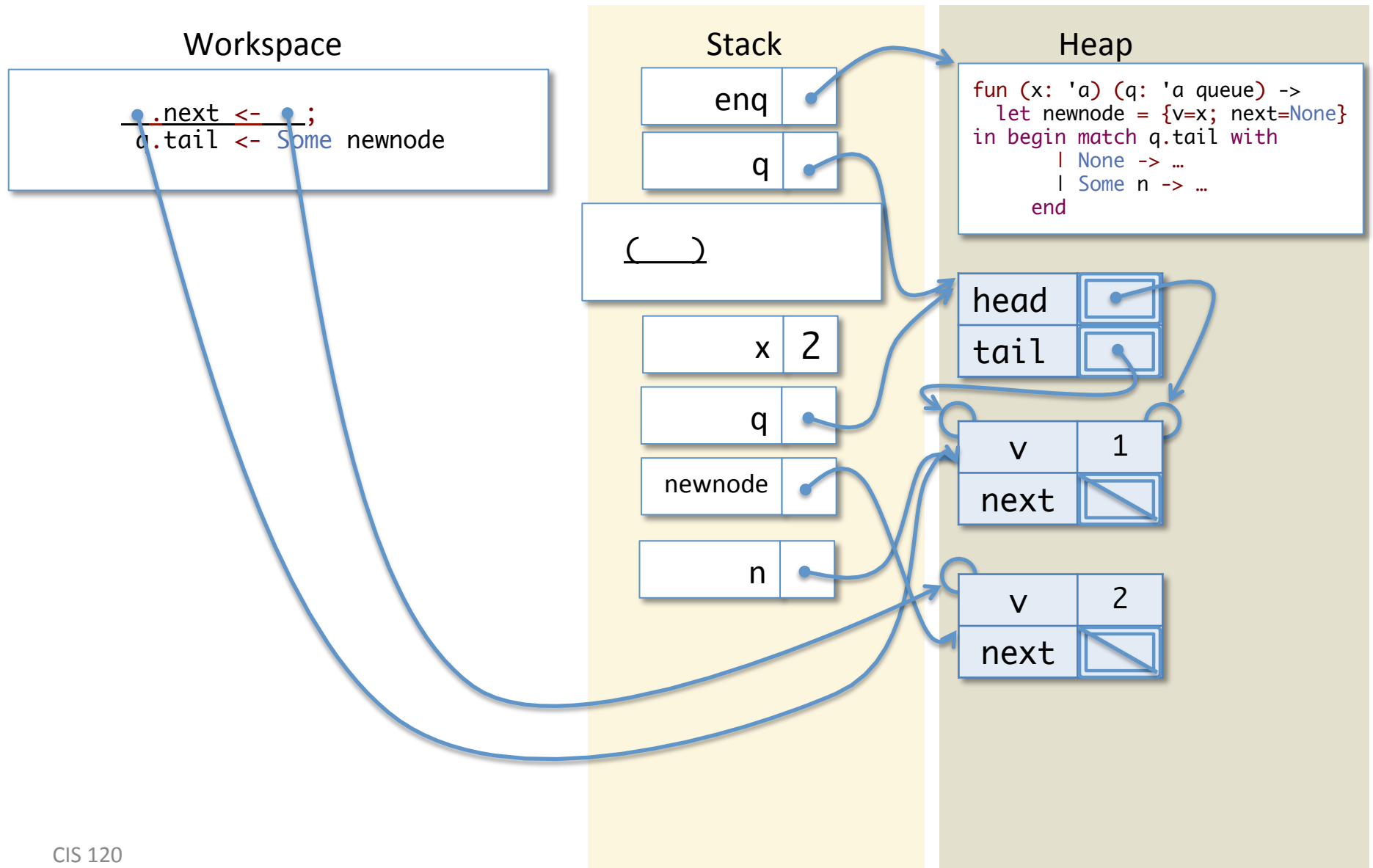
# Calling Enq on a non-empty queue



# Calling Enq on a non-empty queue



# Calling Enq on a non-empty queue

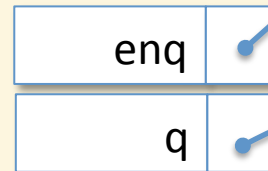


# Calling Enq on a non-empty queue

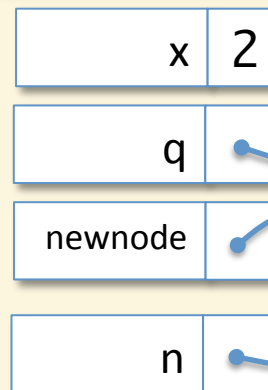
## Workspace

```
();  
q.tail <- Some newnode
```

## Stack

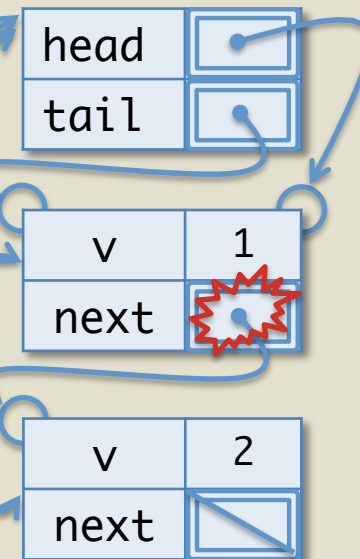


( )



## Heap

```
fun (x: 'a) (q: 'a queue) ->  
  let newnode = {v=x; next=None}  
  in begin match q.tail with  
    | None -> ...  
    | Some n -> ...  
  end
```



# Calling Enq on a non-empty queue

Workspace

```
Q:  
q.tail <- Some newnode
```

Stack

enq

q

( )

x

2

q

newnode

n

Heap

```
fun (x: 'a) (q: 'a queue) ->  
  let newnode = {v=x; next=None}  
  in begin match q.tail with  
    | None -> ...  
    | Some n -> ...  
  end
```

head

tail

v

1

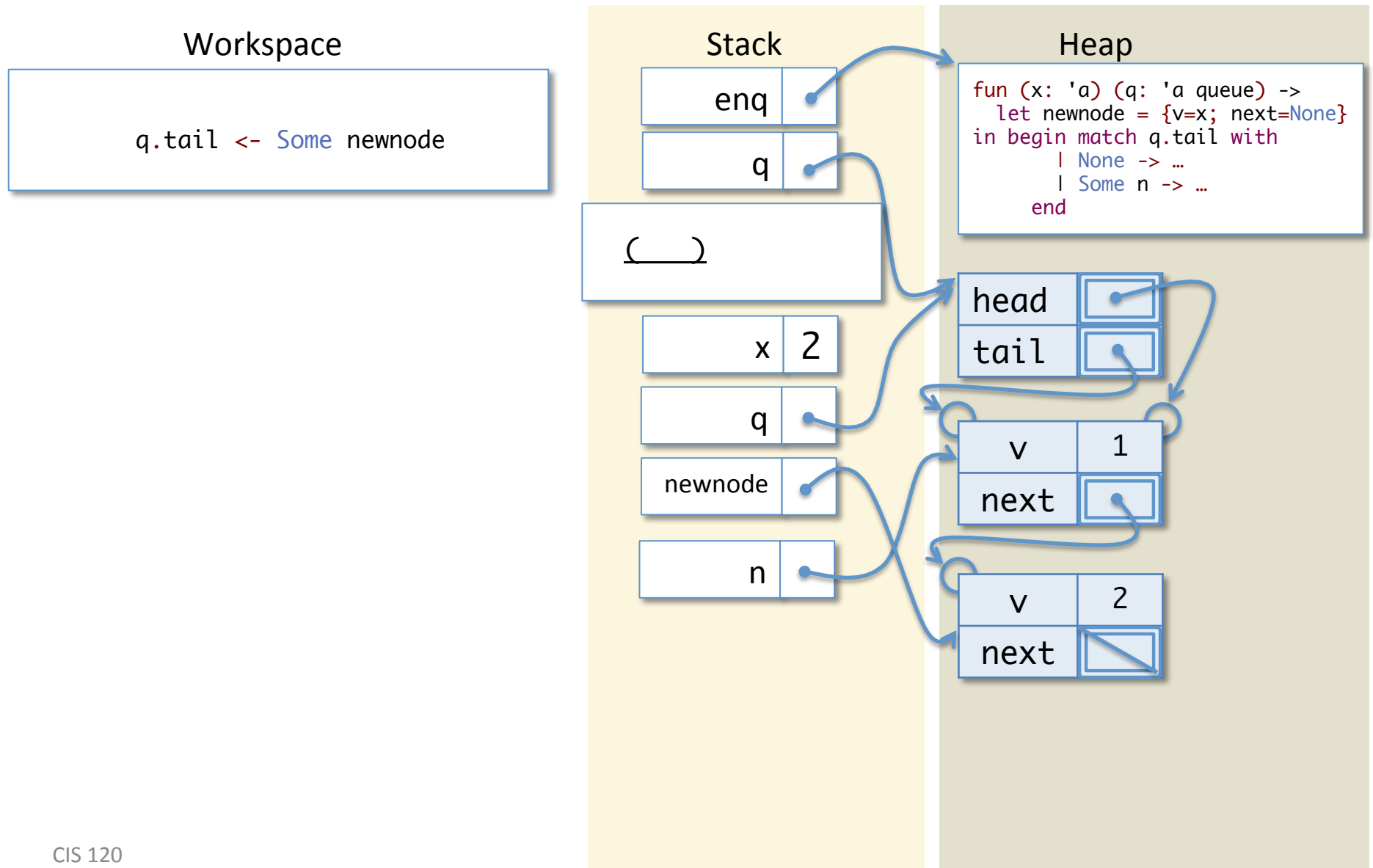
next

v

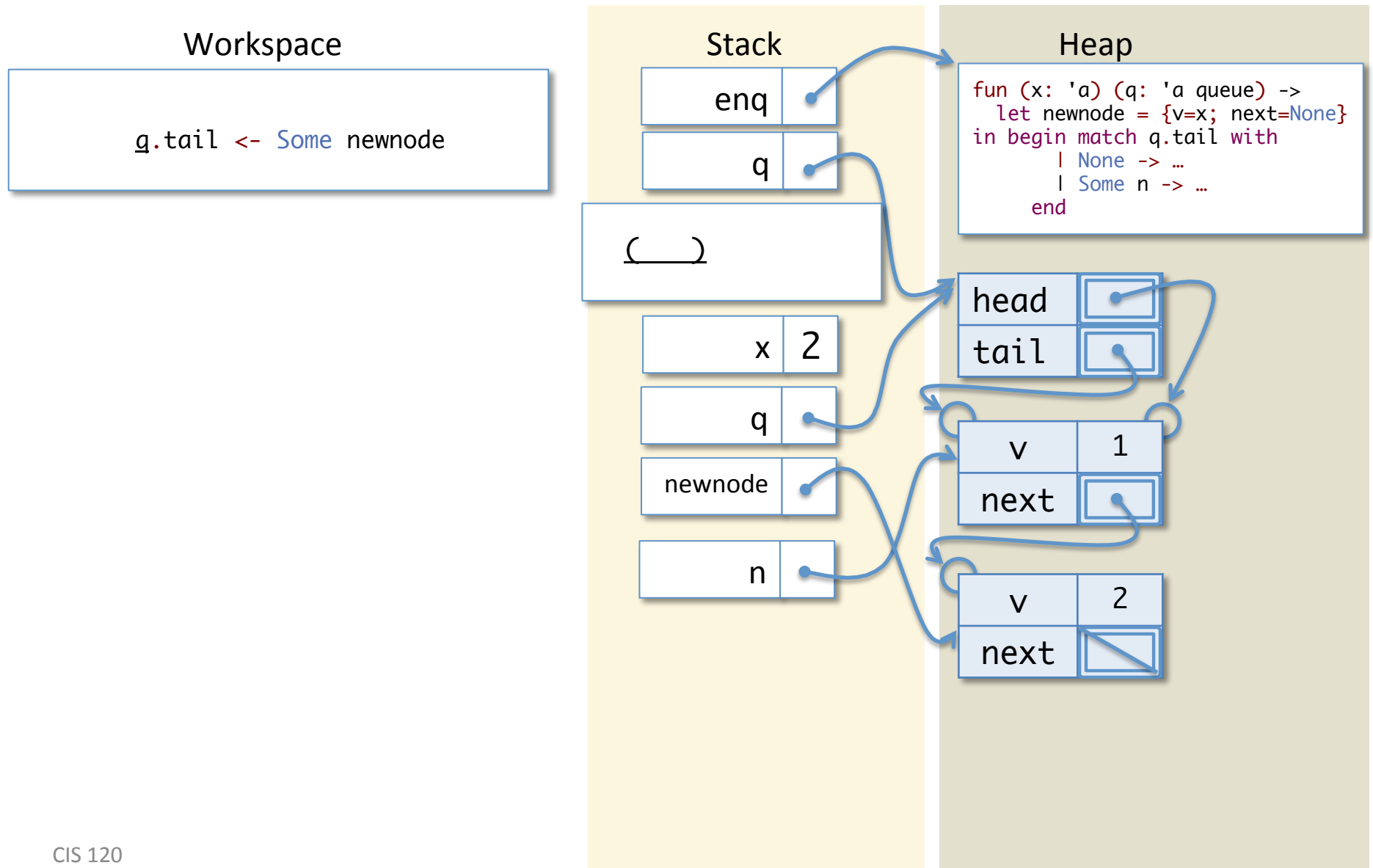
2

next

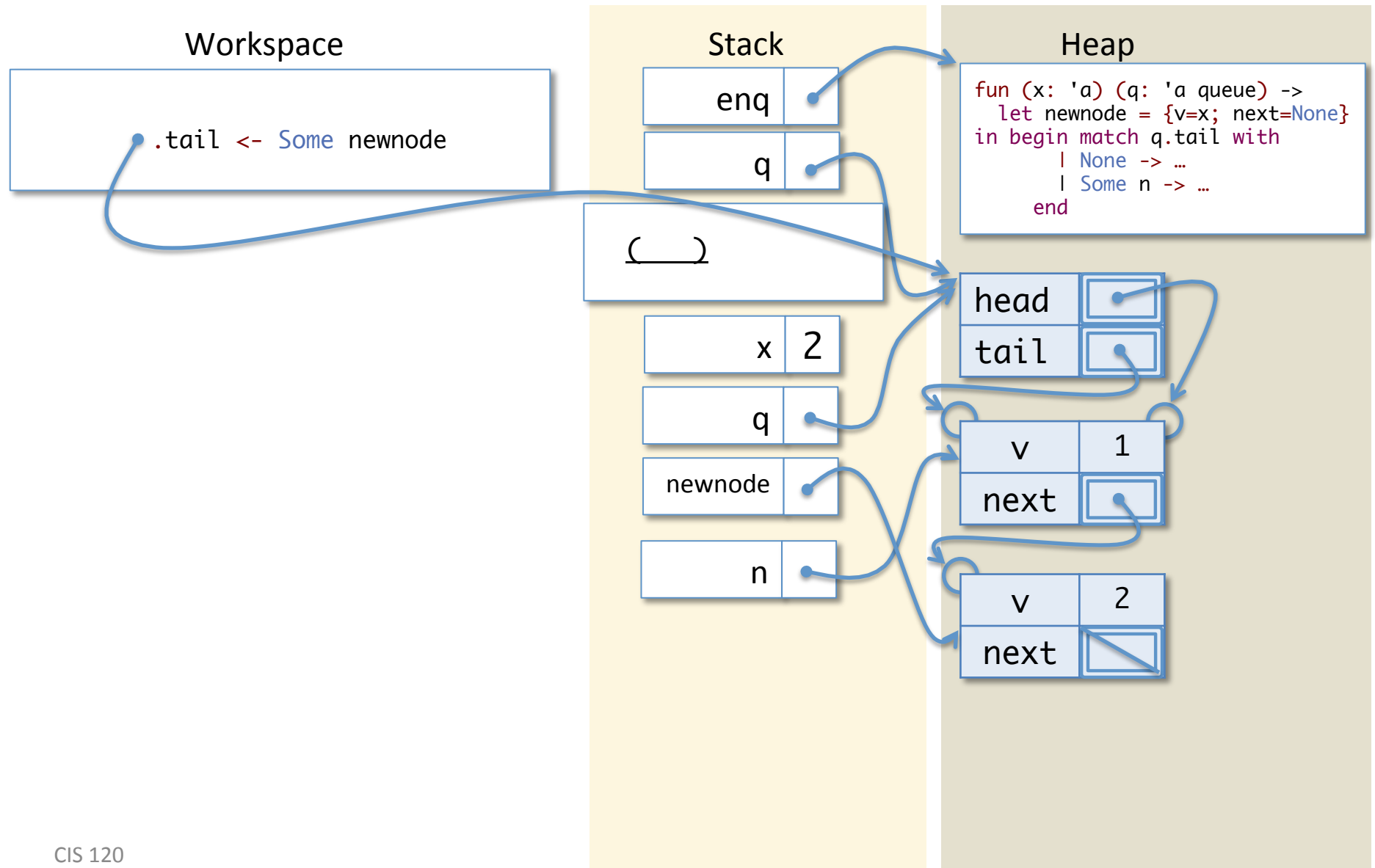
# Calling Enq on a non-empty queue



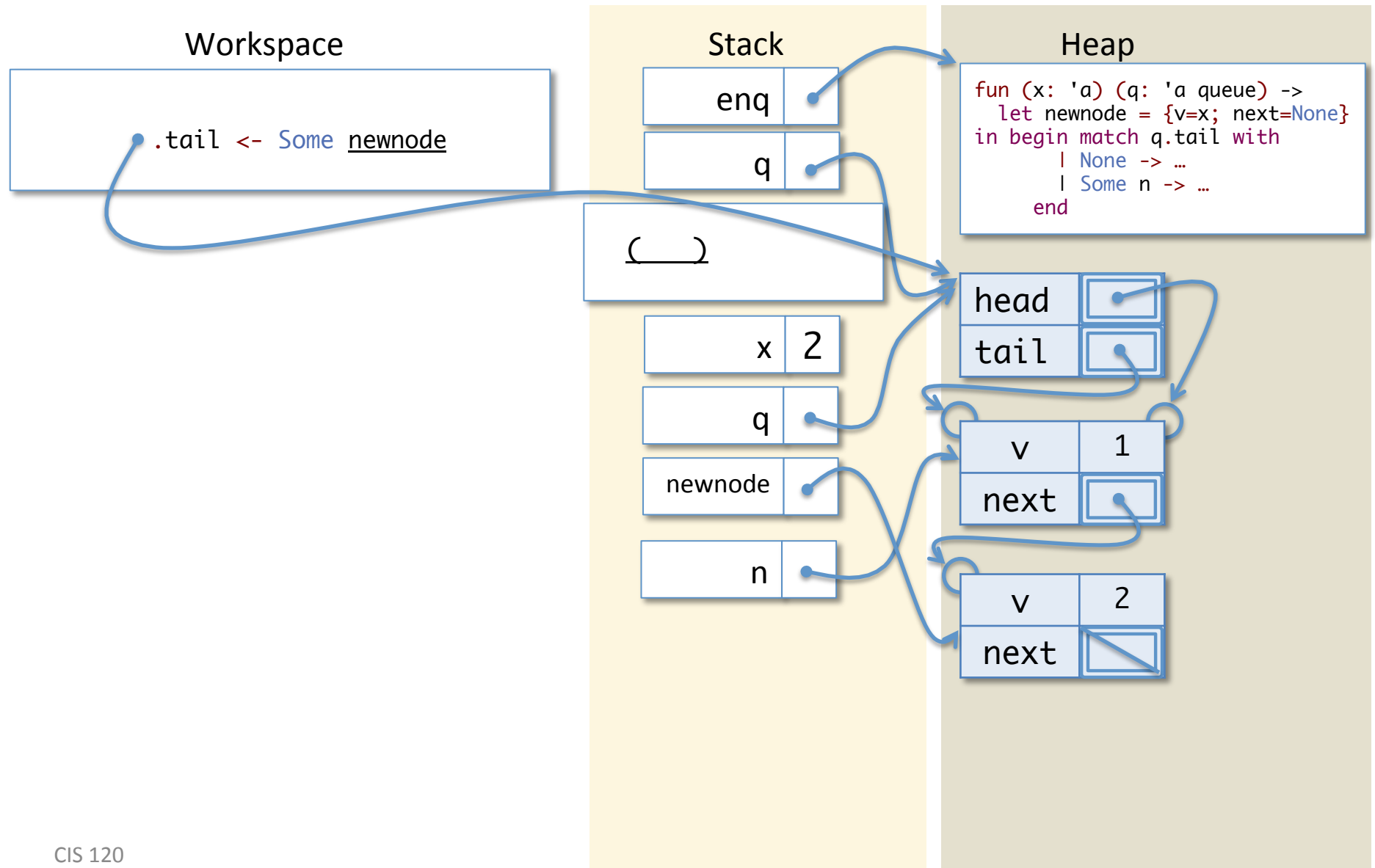
# Calling Enq on a non-empty queue



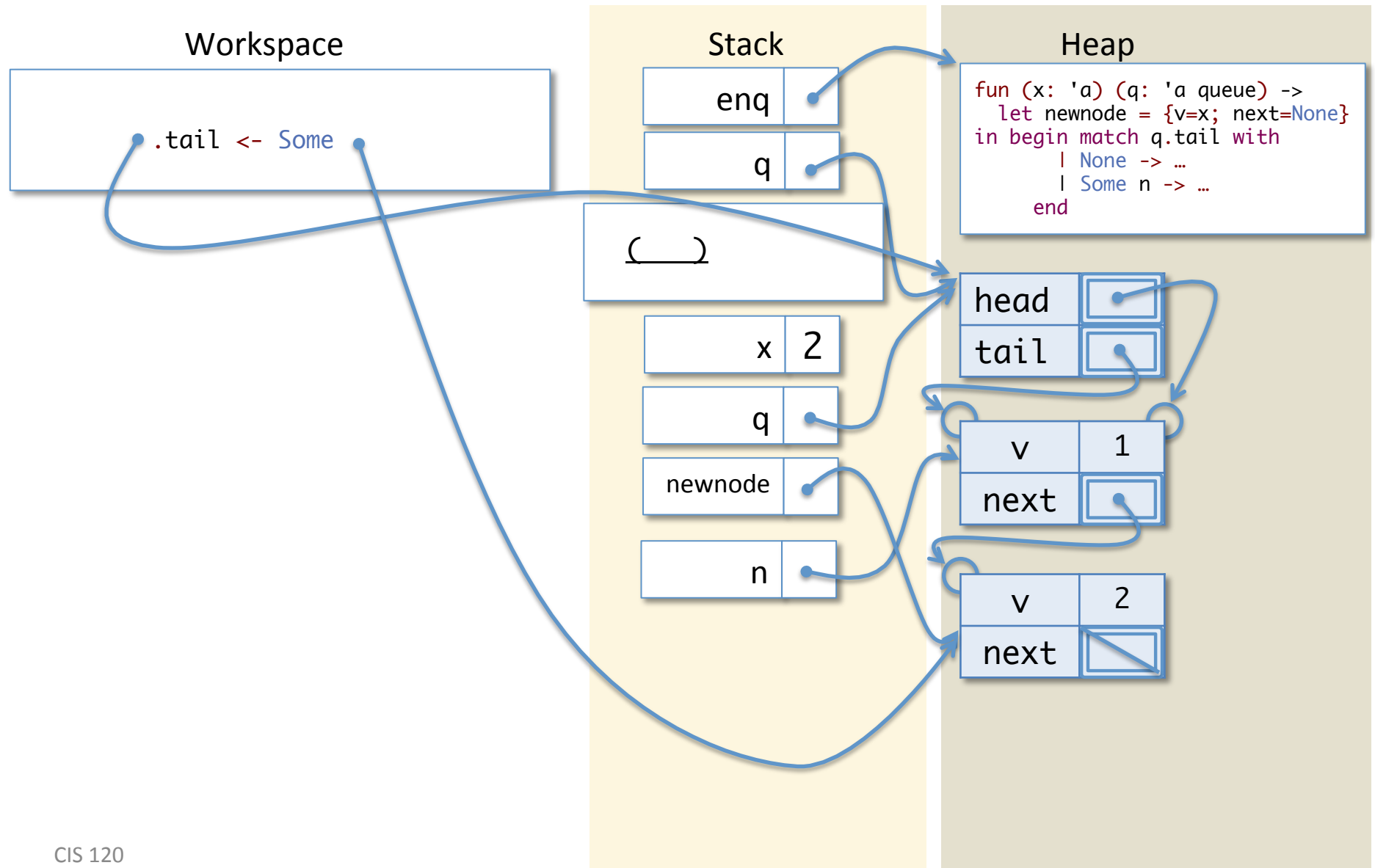
# Calling Enq on a non-empty queue



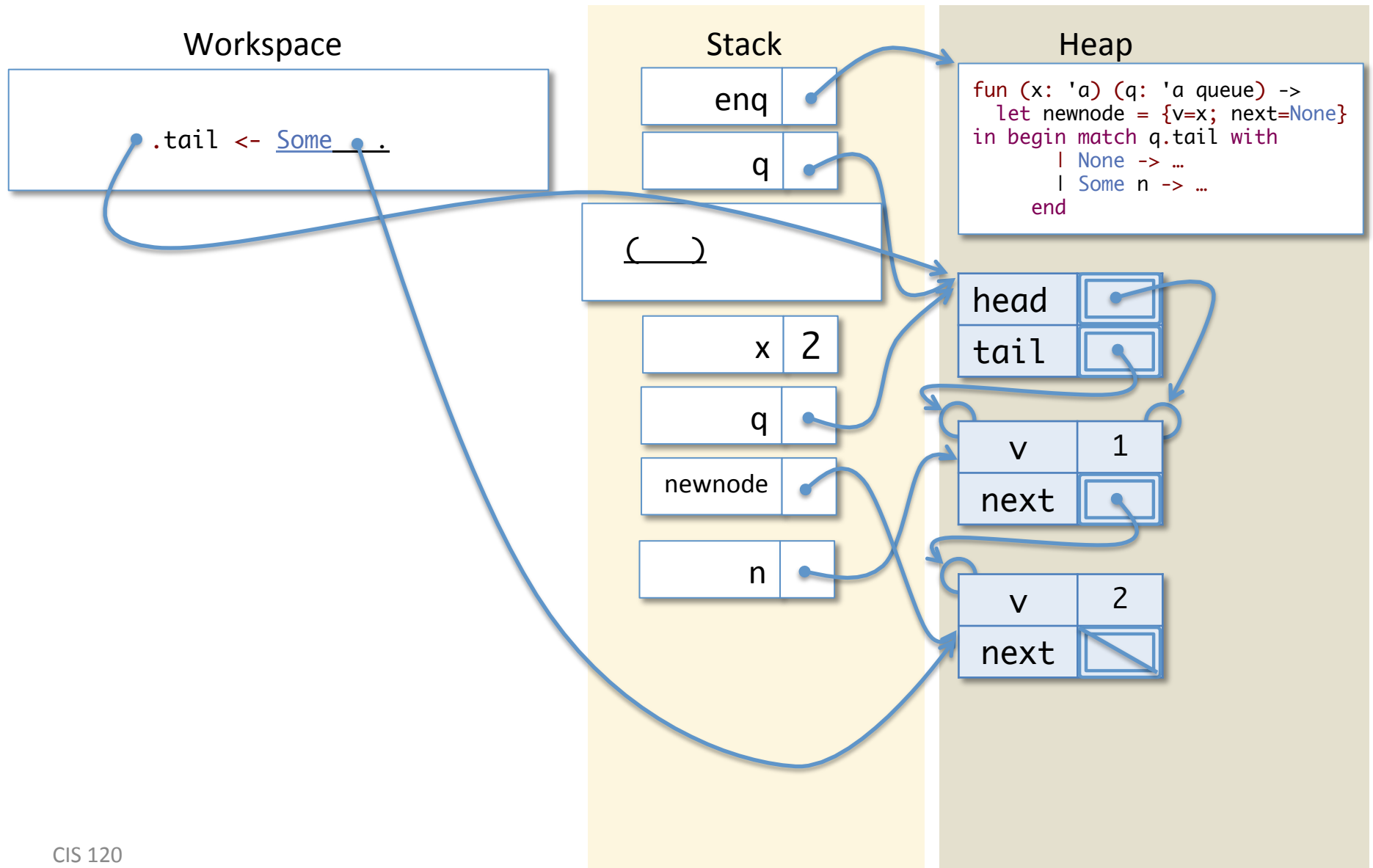
# Calling Enq on a non-empty queue



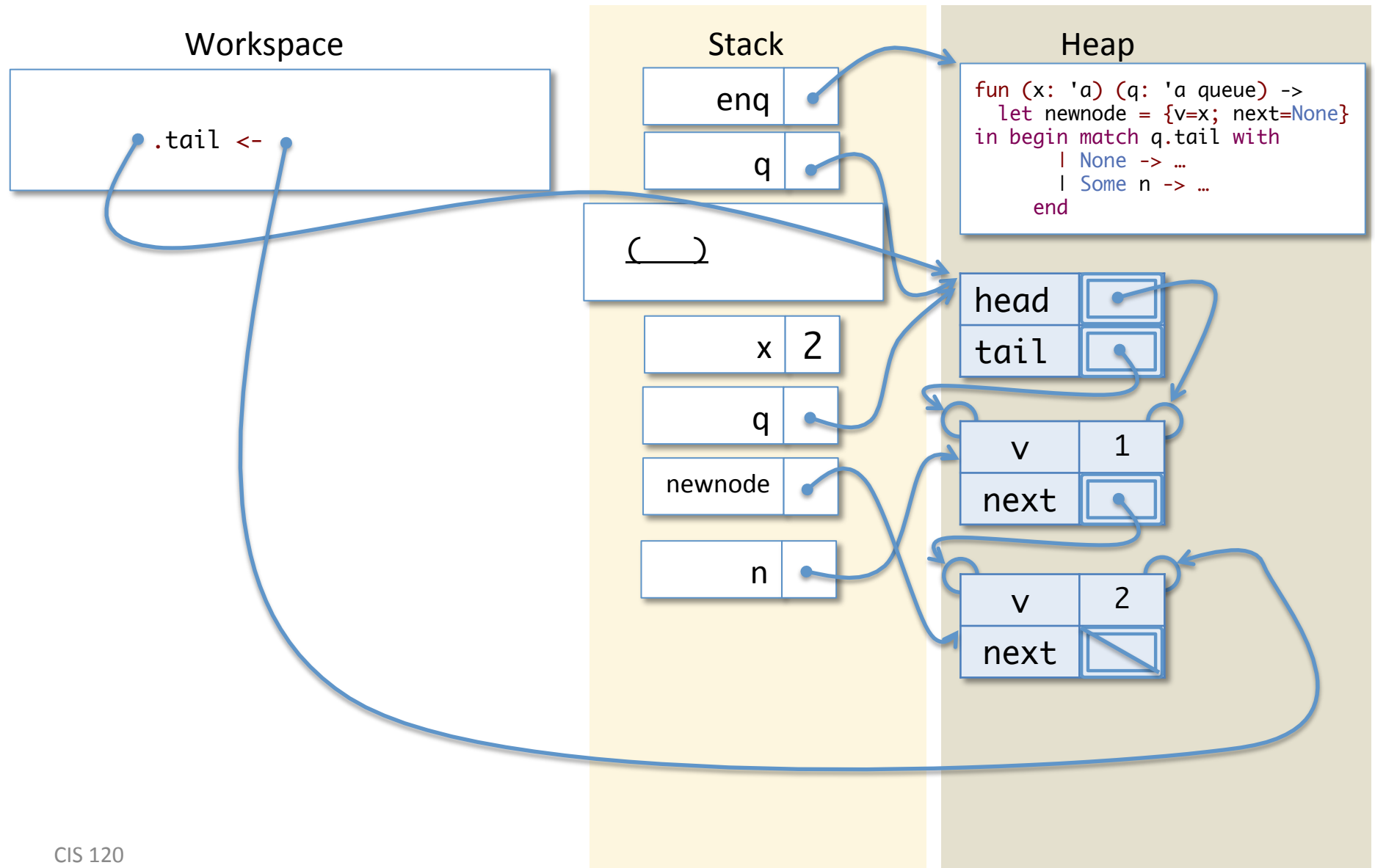
# Calling Enq on a non-empty queue



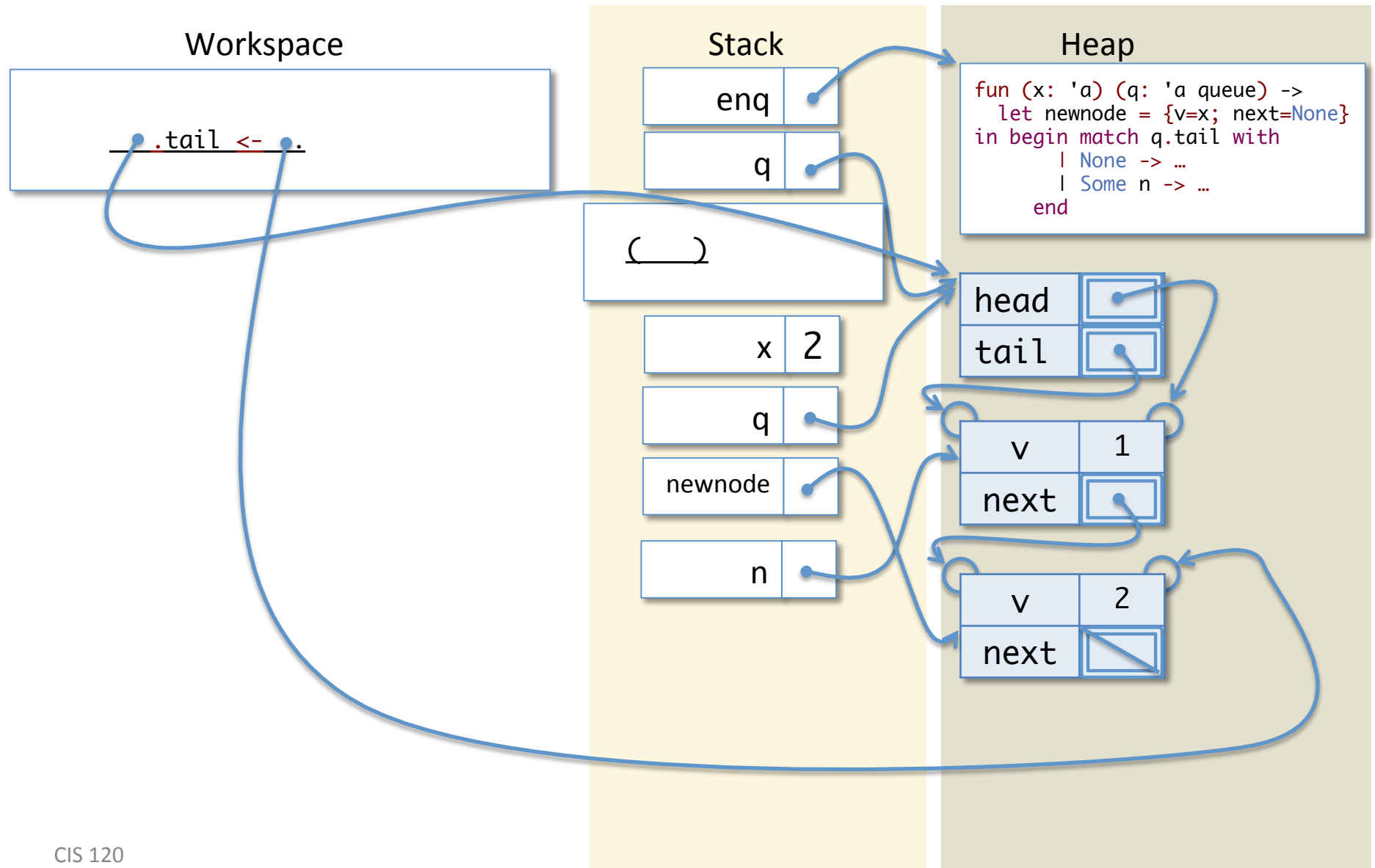
# Calling Enq on a non-empty queue



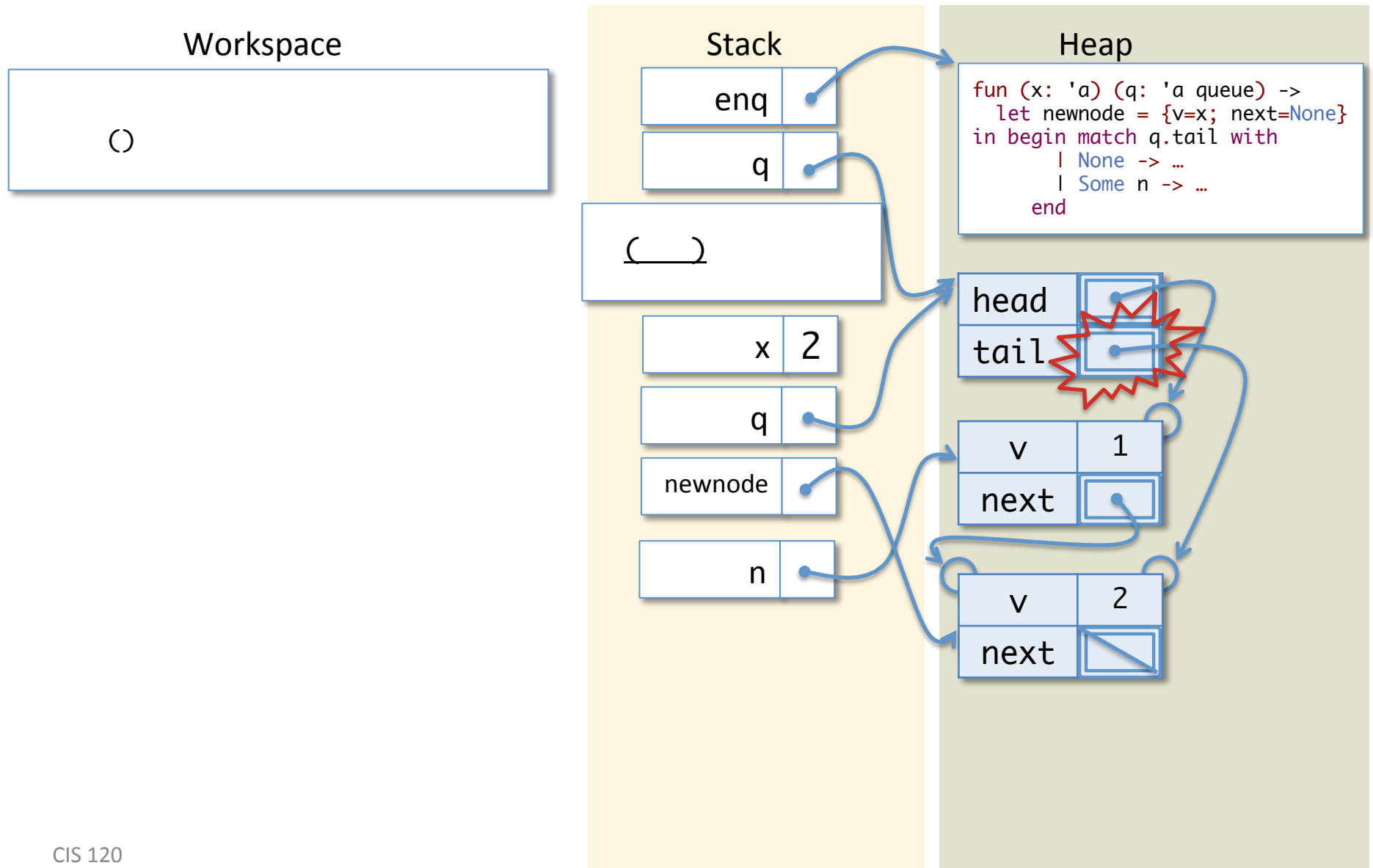
# Calling Enq on a non-empty queue



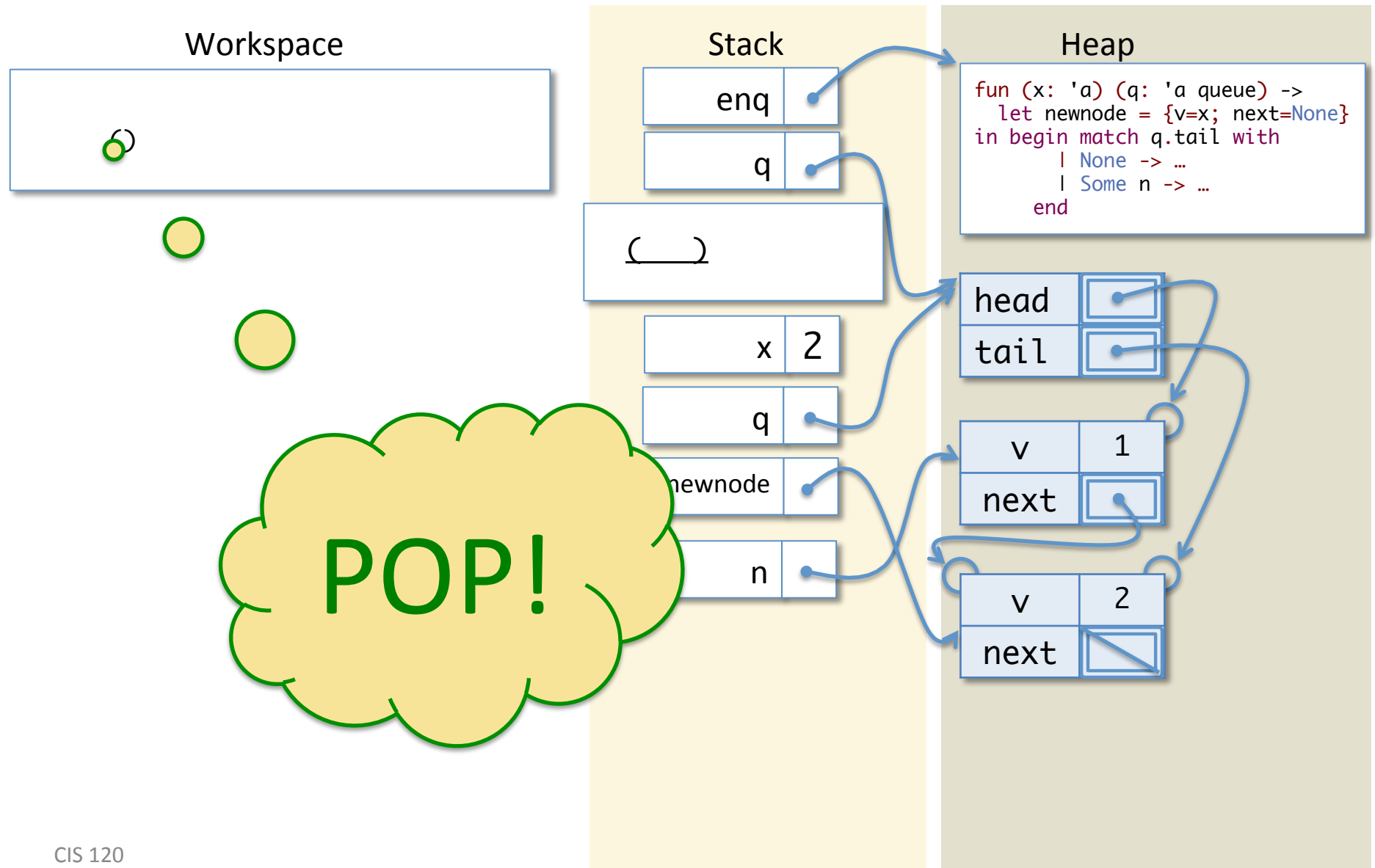
# Calling Enq on a non-empty queue



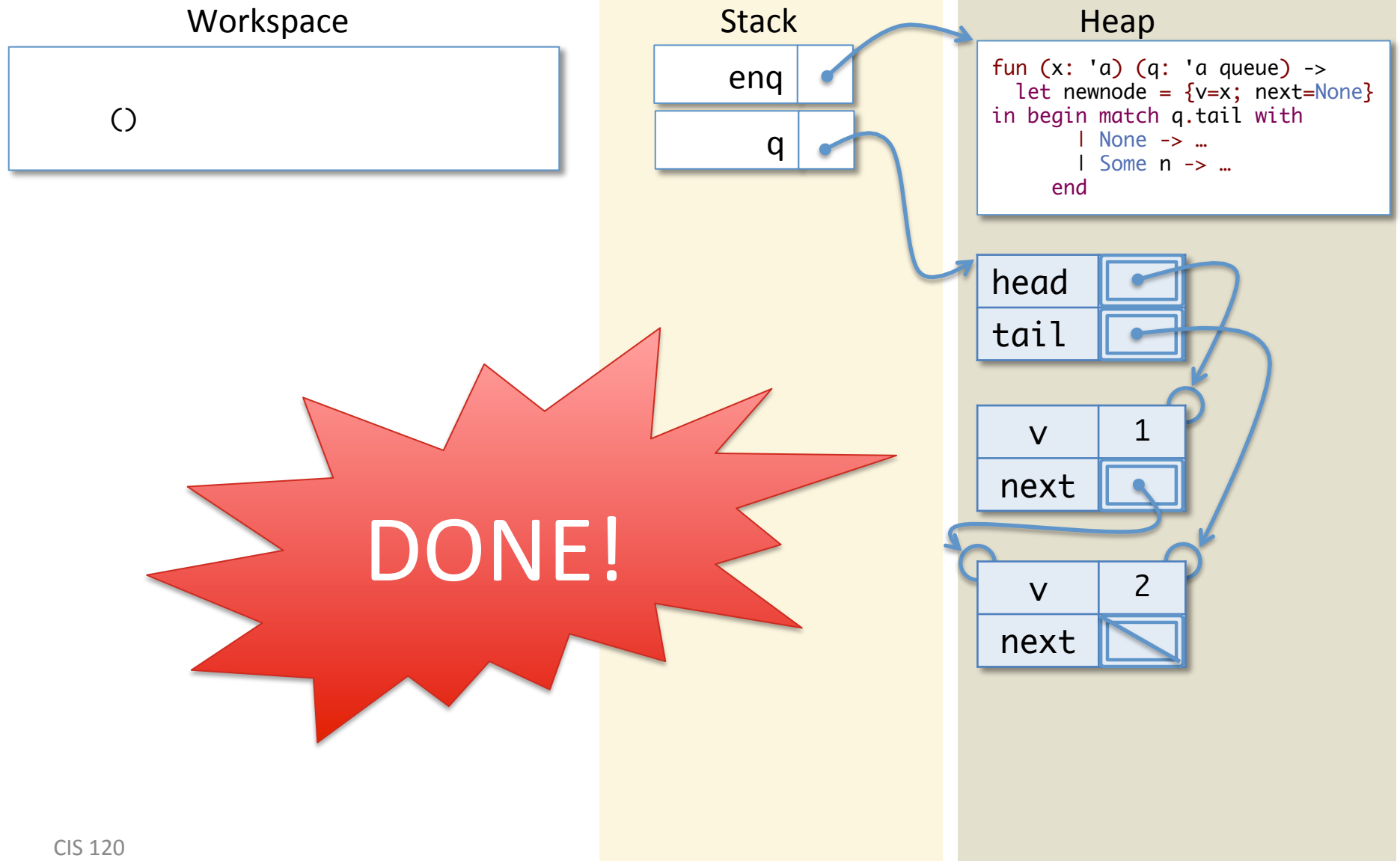
# Calling Enq on a non-empty queue



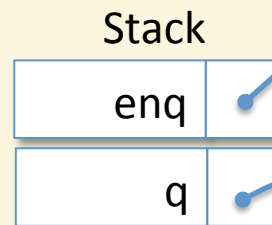
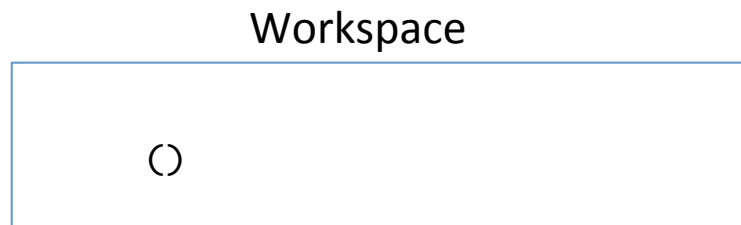
# Calling Enq on a non-empty queue



# Calling Enq on a non-empty queue

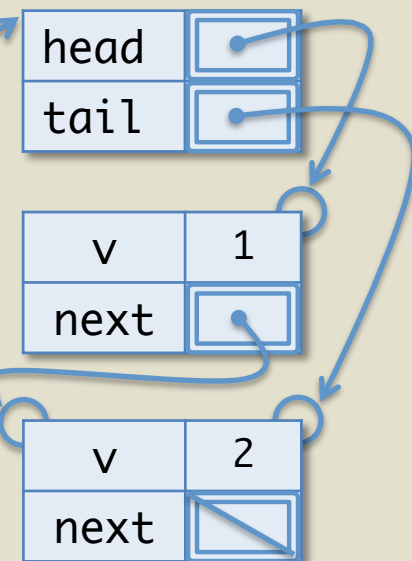


# Calling Enq on a non-empty queue



Heap

```
fun (x: 'a) (q: 'a queue) ->  
  let newnode = {v=x; next=None}  
  in begin match q.tail with  
    | None -> ...  
    | Some n -> ...  
  end
```



Notes:

- the enq function imperatively updated the structure of q
- the new structure still satisfies the queue invariants

# deq

```
(* remove an element from the head of the queue *)
let deq (q: 'a queue) : 'a =
  begin match q.head with
    | None ->
      failwith "deq called on empty queue"
    | Some n ->
      q.head <- n.next;
      if n.next = None then q.tail <- None;
      n.v
  end
```

- The code for `deq` must also “patch pointers” to maintain the queue invariant:
  - The head pointer is always updated to the next element in the queue.
  - If the removed node was the last one in the queue, the tail pointer must be updated to `None`