CIS 120 Midterm I October 4, 2013

Name (printed): _____ Pennkey (login id): _____

My signature below certifies that I have complied with the University of Pennsylvania's Code of Academic Integrity in completing this examination.

 Signature:
 Date:

1	/20
2	/10
3	/16
4	/14
5	/20
6	/20
Total	/100

- Do not begin the exam until you are told to do so.
- You have 50 minutes to complete the exam.
- There are 100 total points.
- Make sure your name and Pennkey (a.k.a. username) is on the top of this page.
- Be sure to allow enough time for all the problems—skim the entire exam first to get a sense of what there is to do.

1. List Processing (20 points)

For each of the following programs, write the value computed for r:

```
a.
     let x : int = 42
     let f (y : int) : int = y + x
     let x : int = 120
     let r : int = f 0
  Answer: r =
b.
     let rec f (l : 'a list) : ('a list * 'a list) =
       begin match 1 with
       | [] -> ([], [])
       | [x] -> (1 , [])
       | u::v::w ->
          let (y,z) : ('a list * 'a list) = f w in
          (u::y, v::z)
       end
     let r : ('a list * 'a list) = f [1;2;3;4;5]
  Answer: r =
c.
     let rec f (x : ('a -> 'b) list) (y : 'a) : 'b list =
       begin match x with
       | h::t -> h y :: f t y
       | _ -> []
       end
     let r : int list = f [(fun x -> - x); (fun x -> x * (x + 1))] 6
  Answer: r =
d.
     type foo = {mutable bar : int}
     let f (x : foo) (y : foo) : int * int * int =
       let z = x in
       x.bar <- y.bar + 1;
       (x.bar, y.bar, z.bar)
     let r : int * int * int = f {bar = 0} {bar = 0}
  Answer: r =
```

2. Types (10 points)

For each OCaml value or function definition below, fill in the blank where the type annotation could go or write "ill typed" if there is a type error. If an expression can have multiple types, give the most generic one. Recall that the @ operator appends two lists together in OCaml. We have done the first one for you. Consider the definitions to be below the following code:

```
module type MAP = sig
 type ('a * 'b) map
 val fromList : ('a * 'b) list -> ('a * 'b) map
end
module LMap : MAP = struct
 type ('a * 'b) map = ('a * 'b) list
 let fromList (l : ('a * 'b) list) = l
end
open LMap;;
let x : _____ int list _____ = [2 + 2]
let a : _____ = 42 ^ " 42"
let b : _____ = [42] :: [[]]
let c : _____ = [42] :: [42]
let d : _____ = [("cis", 120)]
let e : _____ = fromList [(120, 42)]
let f :
fromList ([("benjamin", 42)] @ [("pierce", 120)])
let q : _
 fromList [("benjamin", "cis"); ("pierce", 120)]
let h :
 fromList [(120,fromList [("benjamin", 42)])]
let i : _____
                   _____ = (fun f -> f 42)
let j : _____ = (fun x -> x + "foo")
```

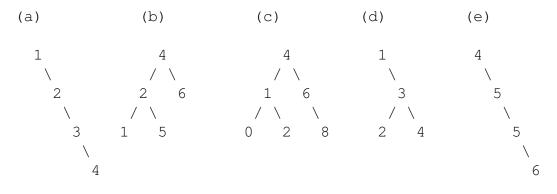
3. Binary Trees and Binary Search Trees (16 points)

Recall the definition of generic binary trees and the BST insert function:

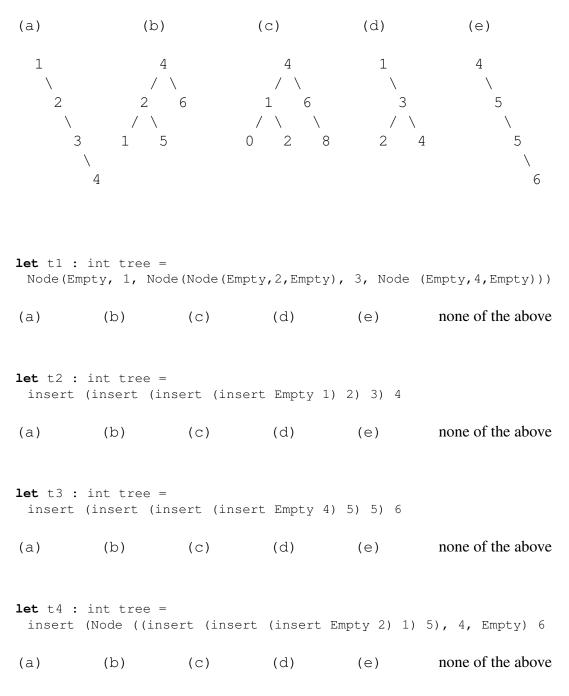
```
type 'a tree =
    | Empty
    | Node of 'a tree * 'a * 'a tree

let rec insert (t:'a tree) (n:'a) : 'a tree =
    begin match t with
    | Empty -> Node(Empty, n, Empty)
    | Node(lt, x, rt) ->
        if x = n then t
        else if n < x then Node (insert lt n, x, rt)
        else Node(lt, x, insert rt n)
    end</pre>
```

a. Circle the trees that satisfy the *binary search tree invariant*. (Note that we have omitted the Empty nodes from these pictures, to reduce clutter.)



b. For each definition below, circle the letter of the tree that it constructs or "none of the above".



5

4. Modules (14 points)

For this question, suppose we've written the following definition of a module signature MOD and two modules ModOne and ModTwo conforming to the signature MOD:

```
module type MOD = sig
 type t
 val x : t
 val f : t -> int
end
module ModOne : MOD = struct
 type t = bool
 let x = false
 let f (b:t) : int = if b then 1 else 0
end
module ModTwo : MOD = struct
 type t = int
 let x = 41
 let y = 16
 let f (b:t) : int = b + 1
end
```

a. Does the following module definition typecheck? If not, *briefly* (one sentence) explain why not.

```
module MyMod3 : MOD = struct
type t = int
let x = 16
end
```

b. Does the following client code typecheck? If so, what value does it print? If not, briefly explain why not.

```
;; print_int (ModOne.f ModOne.x)
```

c. Does the following client code typecheck? If so, what value does it print? If not, briefly explain why not.

```
;; print_int (ModOne.f true)
```

d. Does the following client code typecheck? If so, what value does it print? If not, briefly explain why not.

```
;; print_int (ModTwo.f ModTwo.y)
```

e. Does the following client code typecheck? If so, what value does it print? If not, briefly explain why not.

```
;; print_int (ModOne.f ModTwo.x)
```

f. Does the following client code typecheck? If so, what value does it print? If not, briefly explain why not.

```
;; print_int (ModOne.f ModOne.x + ModTwo.f ModTwo.x)
```

g. List *all* the values that could possibly be printed when we run the following client code (for all possible ways of filling in the body of *z*).

let z : ModTwo.t = ...
;; print_int (ModTwo.f z)

5. Program Design (20 points)

Use the four-step design methodology to implement a function called rotations that computes all the cyclic permutation of a list, ie. given a list $[v_1, v_2, \ldots, v_n]$ of n values (where $n \ge 0$), returns the list

$$[[v_1; v_2; \ldots; v_n]; [v_2; \ldots; v_n; v_1]; [v_n; v_1; \ldots; v_{n-1}]]$$

also of length n.

For example, rotations [1;2;3;4] should yield the list:

```
[ [1;2;3;4]; [2;3;4;1]; [3;4;1;2]; [4;1;2;3] ]
```

- **a.** Step 1 is *understanding the problem*. You don't have to write anything for this part—your answers below will demonstrate whether or not you succeeded with Step 1.
- **b.** Step 2 is *formalizing the interface*. Write down the *type* of the rotations function as you might find it in a .mli file or module interface.

val rotations:

c. Step 3 is *writing test cases*. Complete the following three tests with the expected behavior. We have done the first one for you, based on the problem description.

Note that some test cases are better than others, and credit will be assigned accordingly: make sure your tests cover a sufficiently broad range of "interesting" input numbers and lists. Fill in the description string of the run_test function with a short explanation of *why* the test case is interesting. Your description should not just restate the test case, e.g. "rotations [1;2;3]".

<pre>iii. let test () : bool =</pre>		
(rotations) =	
;; run_test "		" test

d. Step 4 is *implementing the program*. Fill in the body of the rotations function to complete the design. You can use @ or any of the higher order functions in the appendix in your answer. Hint : You can also define an auxilliary function.

let rotations (l : _____) : _____ =

6. Higher-Order Functions (20 points)

In this problem you will be guided into coding a function cartesian, which computes the socalled *cartesian product* of two lists 11 and 12 — that is, the list of all tuples (a, b) where a comes from 11 and b comes from 12.

You should use a single higher-order function — one of transform, filter, or fold — in your solution to each of the subproblems (except the last).

For reference, the definitions of these functions are given in the appendix.

a. Write a function make_product_list that, given a list 1 and a value a, returns a list of tuples whose first element is a and whose second elements ranges over all the elements of 1. For example, make_product_list [3;4;5] 1 = [(1,3); (1,4); (1,5)]. Your answer should be a call of a single higher-order function with some argument
let make_product_list (1 : 'b list) (a : 'a): ('a * 'b) list =

b. Write a function cartesian_helper that given a list [a1; a2; ...; an] and another [b1; b2; ...; bm], where $m, n \ge 0$, produces the following list of lists:

For example,

cartesian_helper [1;2] [3;4] = [[(1, 3); (1, 4)]; [(2, 3); (2, 4)]].

Hint: You should use a higher-order function and make_product_list.

let cartesian_helper (l1 : 'a list) (l2 : 'b list) : ('a * 'b) list list =

c. Write a function concat that given a list of lists produces a single list that is the concatenation of all the elements of the original list. For example:

concat [[(1,3);(1,4)];[(2,3);(2,4)]] = [(1,3);(1,4);(2,3);(2,4)]
Again, you answer should be a call to a single higher-order function with some argument.
let concat (l : 'a list list) : 'a list =

d. Finally, use the functions you defined above to write the function cartesian, defined in the beginning of this problem. For example,

cartesian [1;2] [3;4;5] = [(1,3);(1,4);(1,5);(2,3);(2,4);(2,5)]Hint: Do not use any higher-order functions or recursion here. A simple combination of the functions you have coded above should be enough!

let cartesian (l1 : 'a list) (l2 : 'b list) : ('a * 'b) list =

Appendix

```
let rec transform (f: 'a -> 'b) (x: 'a list): 'b list =
begin match x with
       [] -> []
            h :: t -> (f h) :: (transform f t)
    end
let rec fold (combine: 'a -> 'b -> 'b) (base: 'b) (l: 'a list): 'b =
begin match l with
       [] -> base
            h :: t -> combine h (fold combine base t)
    end
let rec filter (f: 'a -> bool) (l: 'a list) : 'a list =
begin match l with
            [] -> []
            h ::t -> if f h then h :: filter f t else filter f t
        end
```