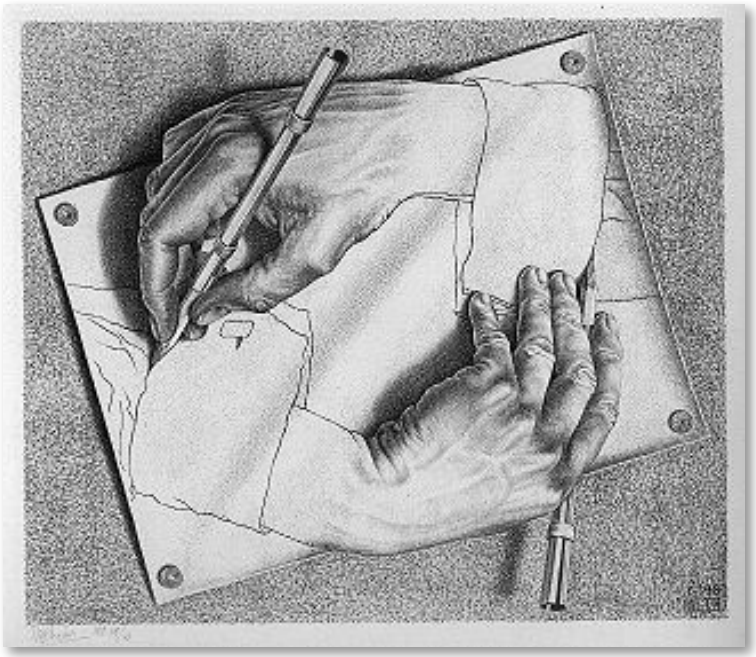
Programming Languages and Techniques (CIS120)

Bonus Lecture

April 25, 2016

"Code is Data"

Code is Data



M.C. Escher, Drawing Hands, 1948

Code *is* Data

- A Java source file is just a sequence of characters.
- We can represent programs with Strings!

```
String p_0 = "class C { public static void main(String args[])
String p_1 = "class C { public static void main(String args[])
String p_2 = "class C { public static void main(String args[])
String p_3 = "class C { public static void main(String args[])
String p_13 = "class C { public static void main(String args[])
 {System.out.println(\"Hello, world!\");}}";
```

• • •

```
String p_12012023423123123O = /* Mushroom of Doom! */
      "class Game { public static void main(String args[]) {…}}";
```

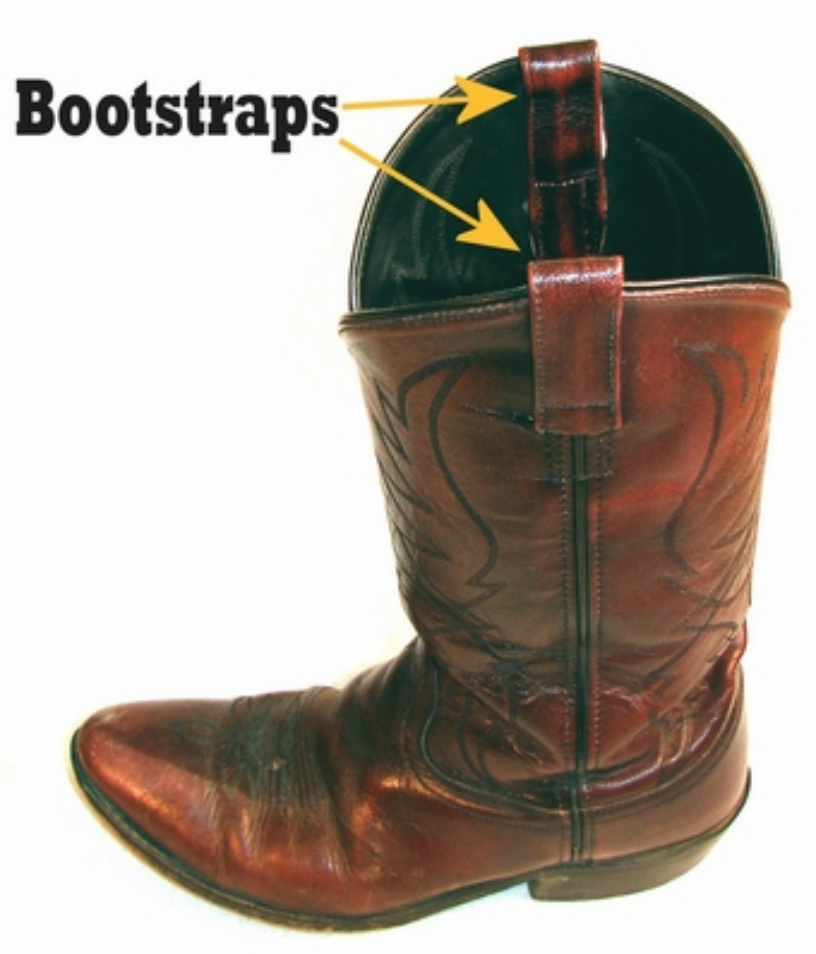
• • •

```
String p_99993249000902300239400823407O234 = /* Minecraft! */
      "class Minecraft { public static void main(String args[]) {…}}";
```

• • •

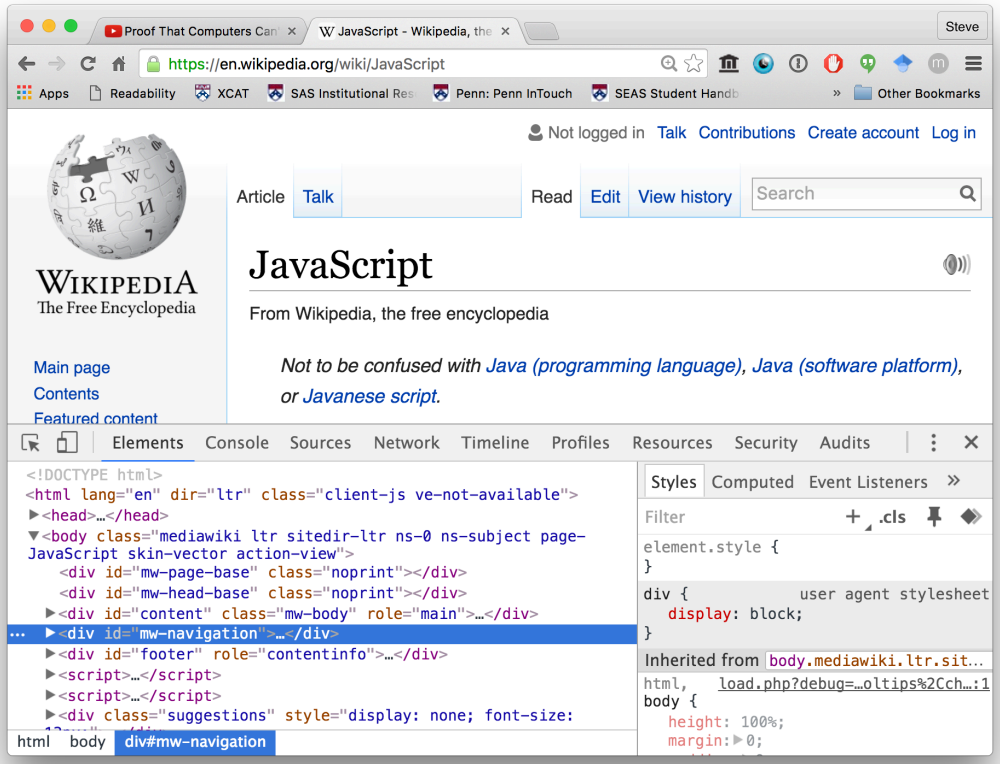
```
String p_9923499234239999932490002342823407345O234534 = /* Eclipse! */
      "class Eclipse { public static void main(String args[]) {…}}";
```

Consequence 1: Programs that manipulate programs



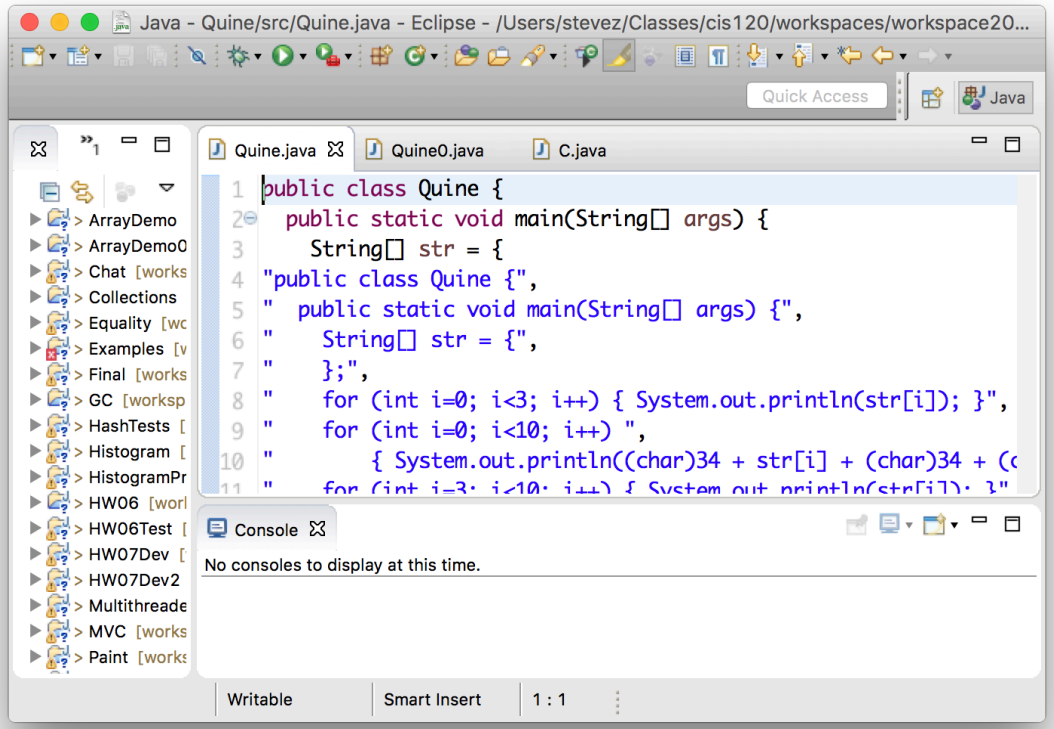
Interpreters

- We can create *programs* that manipulate *programs*
- An *interpreter* is a program that executes other programs
- `interpret ("3 + 4")` ➔ `7`
- Example 1: JavaScript



IDEs and Compilers

- Example 2: Eclipse
  - Note that Eclipse manipulates a *representation* of Java programs
  - Eclipse itself is written in Java
  - So you could use Eclipse to edit the code for Eclipse… ?!

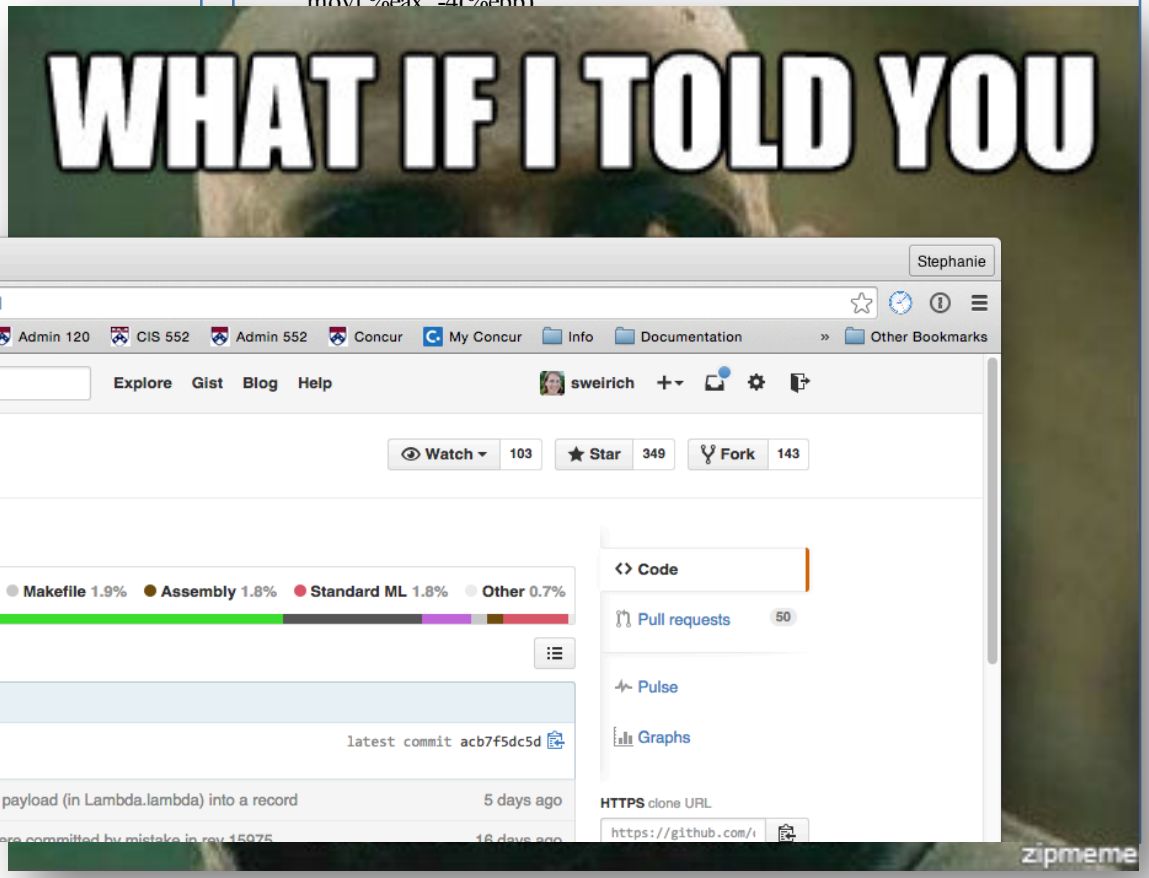


- Example 3: Compiler
  - The Java compiler takes a representation of a Java program
  - It outputs a "low-level" representation of the program as a .class file (i.e. Java byte code)
  - Can also compile to other representations, e.g. x86 "machine code"

Example Compilation: Java to X86

```
class Point {
   int x;
   int y;
   Point move(int
   int dy) {
      x = x + dx;
   }
}
```

```
.globl __fun__Point.move
__fun__Point.move:
      pushl %ebp
      movl %esp, %ebp
      subl $4, %esp
__5:
      movl 8(%ebp), %eax
      movl 4(%eax), %eax
      movl %eax, -4(%ebp)
```



Consequence 2: Malware

Rene Magritte, The Human Condition, 1933

Consequence 2: Malware

- Why does Java do array bounds checking?
- *Unsafe* language like C and C++ don't do that checking;
  - They will happily let you write a program that "writes past" the end of an array.
- Result:
  - viruses, worms, "jailbreaking" mobile phones, Spam, botnets, …
- Fundamental issue:
  - Code is data.
  - Why?

Consider this C Program

```
void m() {
  char[2] buffer;

  char c = read();
  int i = 0;
  while (c != -1) {
    buffer[i] = c;
    c = read();
    i++;
  }
  process(buffer);
}

void main() {
  m();
  // do some more stuff
}
```

Notes:

- C doesn't check array bounds
- Unlike Java, it stores arrays directly on the stack
- What could possibly go wrong?

Abstract Stack Machine

"Stack Smashing Attack"

Abstract Stack Machine

Workspace

Stack

```
m();
// do some more stuff
```

Call to main() to start the program…

Abstract Stack Machine

Workspace

```
char[2] buffer;

char c = read();
int i = 0;
while (c != -1) {
  buffer[i] = c;
  c = read();
  i++;
}
process(buffer);
```

Stack

```
_;
// do some more stuff
```

Push the saved workspace, run m()

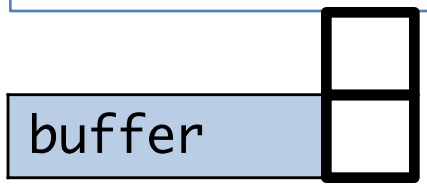
Abstract Stack Machine

Workspace

```
char c = read();
int i = 0;
while (c != -1) {
  buffer[i] = c;
  c = read();
  i++;
}
process(buffer);
```

Stack

```
_;
// do some more stuff
```



Allocate space for buffer on the stack.

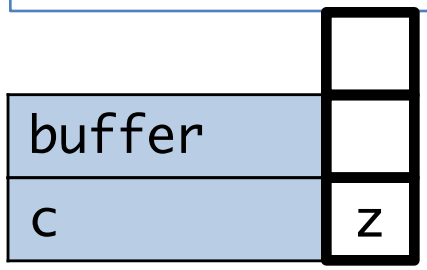
Abstract Stack Machine

Workspace

```
int i = 0;
while (c != -1) {
  buffer[i] = c;
  c = read();
  i++;
}
process(buffer);
```

Stack

```
_;
// do some more stuff
```



Allocate space for c.

Read the first user input… 'z'.

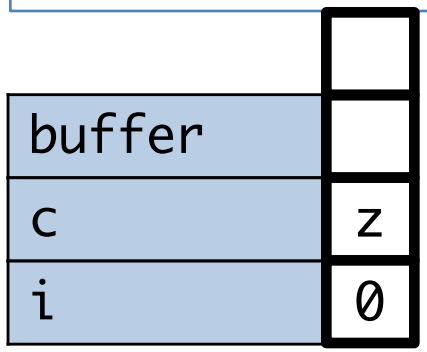
Abstract Stack Machine

Workspace

```
while (c != -1) {
  buffer[i] = c;
  c = read();
  i++;
}
process(buffer);
```

Stack

```
_;
// do some more stuff
```



Allocate space for i.

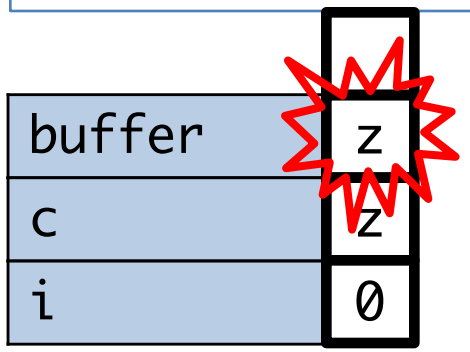
Abstract Stack Machine

Workspace

```
while (c != -1) {
   buffer[i] = c;
   c = read();
   i++;
}
process(buffer);
```

Stack

```
_;
// do some more stuff
```



Copy (contents of) c to buffer[0]

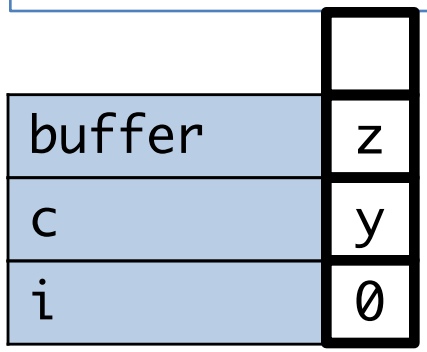
Abstract Stack Machine

Workspace

```
while (c != -1) {
  buffer[i] = c;
  c = read();
  i++;
}
process(buffer);
```

Stack

```
_;
// do some more stuff
```



Read next character … 'y'

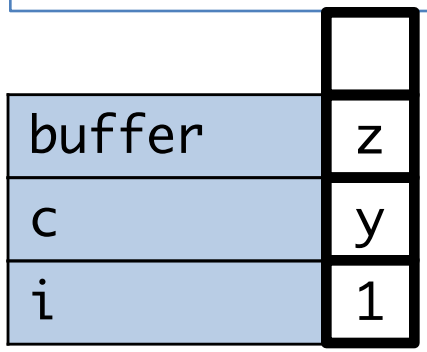
Abstract Stack Machine

Workspace

```
while (c != -1) {
  buffer[i] = c;
  c = read();
  i++;
}
process(buffer);
```

Stack

```
_;
// do some more stuff
```



Increment i

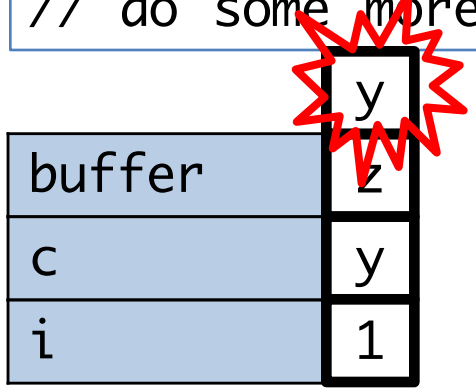
Abstract Stack Machine

Workspace

```
while (c != -1) {
   buffer[i] = c;
   c = read();
   i++;
}
process(buffer);
```

Stack

```
_;
// do some more stuff
```



Copy (contents of) c to buffer[1]

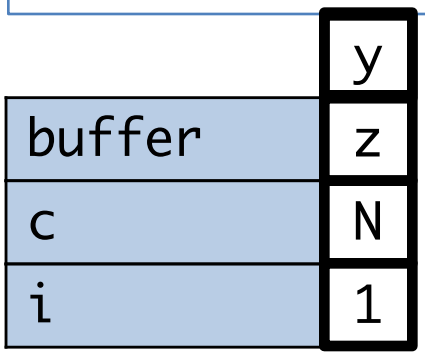
Abstract Stack Machine

Workspace

```
while (c != -1) {
  buffer[i] = c;
  c = read();
  i++;
}
process(buffer);
```

Stack

```
_;
// do some more stuff
```



Read next character … 'N'

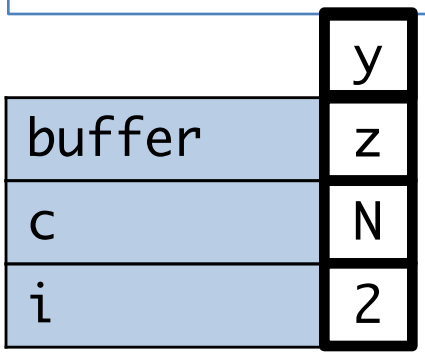
Abstract Stack Machine

Workspace

```
while (c != -1) {
  buffer[i] = c;
  c = read();
  i++;
}
process(buffer);
```

Stack

```
_;
// do some more stuff
```



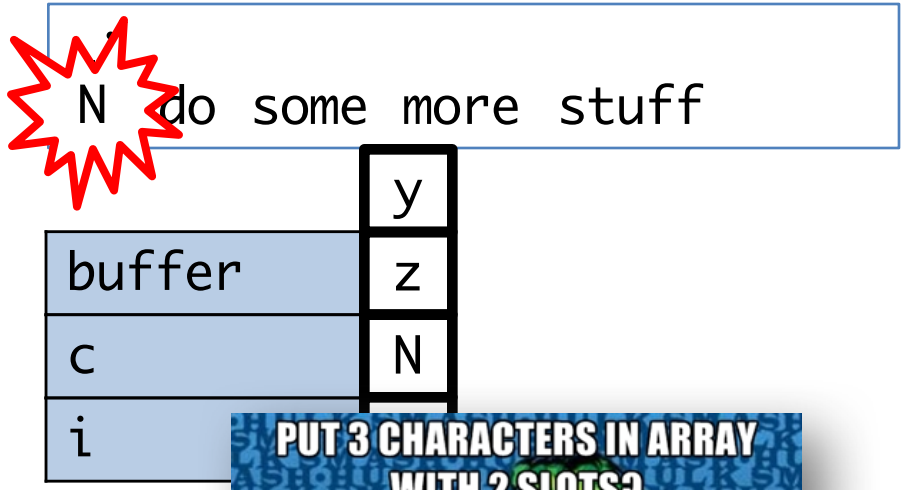
Increment i

Abstract Stack Machine

Workspace

```
while (c != -1) {
   buffer[i] = c;
   c = read();
   i++;
}
process(buffer);
```

Stack



Copy (contents of) c to buffer[2] ?!?

Overwrites the saved workspace!?

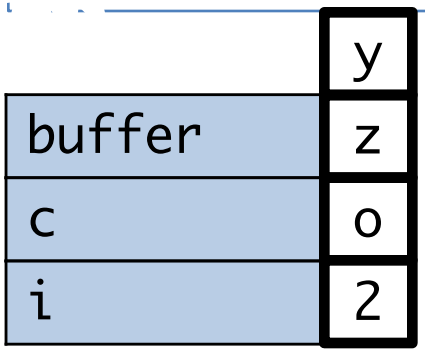
Abstract Stack Machine

Workspace

```
while (c != -1) {
   buffer[i] = c;
   c = read();
   i++;
}
process(buffer);
```

Stack

```
N  do some more stuff
```



Keep going… read 'o'…

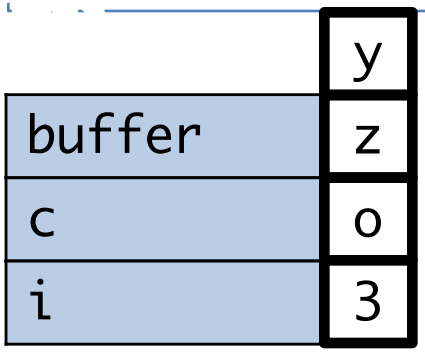
Abstract Stack Machine

Workspace

```
while (c != -1) {
   buffer[i] = c;
   c = read();
   i++;
}
process(buffer);
```

Stack

```
N  do some more stuff
```



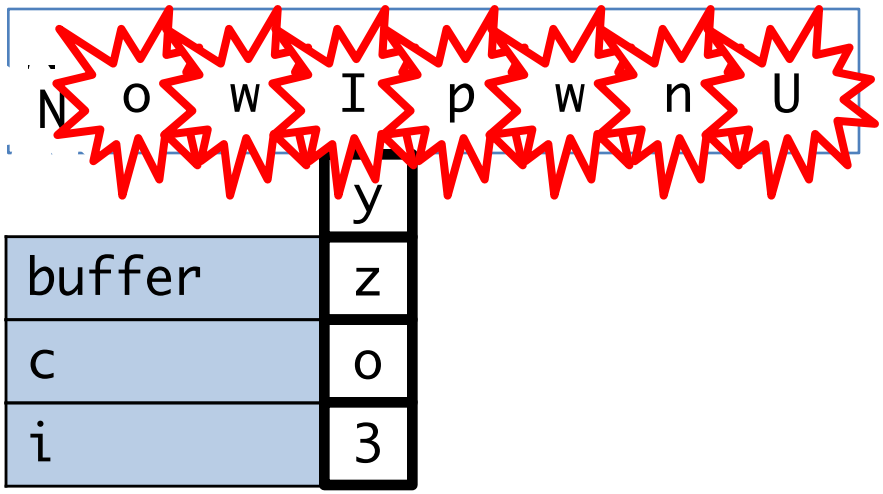
Keep going… read 'o'…increment i…

Abstract Stack Machine

Workspace

```
while (c != -1) {
   buffer[i] = c;
   c = read();
   i++;
}
process(buffer);
```

Stack



Keep going… read 'o'…increment i…write 'o' into saved workspace…

Abstract Stack Machine

Workspace

Stack

Now I pwn U!!!!

buffer	z
c	o
i	3

Later…

Abstract Stack Machine

Workspace

Stack

```
Now I pwn U!!!!
```

The stack smashing attack successfully wrote *arbitrary* code into the program's workspace…

RogueWave SOFTWARE
CodeBuzz

The Top Five Cyber Security Vulnerabilities

POSTED IN GENERAL SECURITY, INCIDENT RESPONSE ON JULY 2, 2015

US-CERT
UNITED STATES COMPUTER EMERGENCY READINESS TEAM

Buffer overflows are the top software security vulnerability of the past 25 years

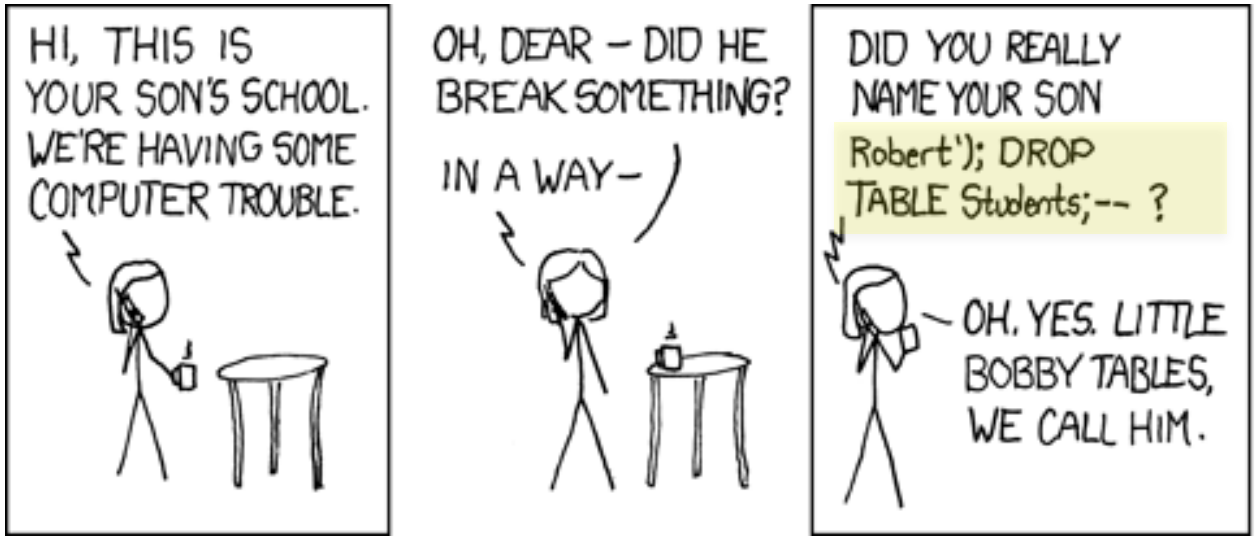
ON MAR 11, 13 • BY CHRIS BUBINAS • WITH 2 COMMENTS

In a report analyzing the entire CVE and NVD databases, which date back to 1988, Sourcefire

Other Code Injection Attacks

```
void registerStudent() {
  print("Welcome to student registration.");
  print("Please enter your name:");
  String name = readLine();
  evalSQL("INSERT INTO Students('" + name + "')" );
}
```

```
"INSERT INTO Students('Robert'); DROP TABLE Students; --')"
    + "Robert'); DROP TABLE Students; --" + "')"
```



http://xkcd.com/327/

Consequence 3: Undecidability

Undecidability Theorem

Theorem: It is impossible to write a method

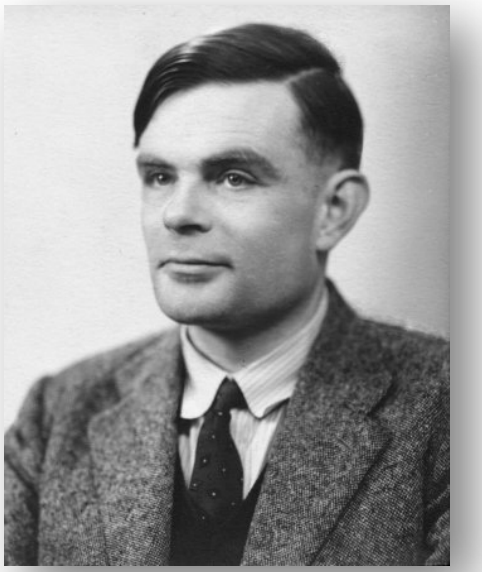
```
boolean halts(String prog)
```

such that, for any valid Java program P represented as a string `p_P`*,*

```
halts(p_P)
```

returns true exactly when the program P halts, and false otherwise.

Alonzo Church, April 1936



Alan Turing, May 1936

Halt Detector

- Suppose we could write such a program:

```
class HaltDetector {
  public static boolean halts(String javaProgram) {
    // ...do some super-clever analysis...
    // return true if javaProgram halts
    // return false if javaProgram does not
  }
}
```

- A correct implementation of HaltDetector.halts(p) always returns either true or false
  - i.e., it never raises an exception or loops
- HaltDetector.halts(p) $\Rightarrow$ true means "p halts"
- HaltDetector.halts(p) $\Rightarrow$ false means "p loops forever"

Do these methods halt?

“boolean m(){ return false; }”

⇒ YES

“boolean m(){ return m(); }”

⇒ NO (assuming infinite stack space)

```
“boolean m(){
    if (“abc”.length() == 3) return true;
    else return m(); }”
```

⇒ YES

```
“boolean m(){
    String x = “”;
    while (true) {
        if (x.length() == 3) return true;
        x = x + \“a\”;
     }
   return false;
}”
```

⇒ YES

Do these methods halt?

```
boolean m(){ return false; }
```

⇒ YES

```
boolean m(){ return m(); }
```

⇒ NO (assuming infinite stack space)

```
boolean m() {
    if ("abc".length() == 3 ) return true;
    else return m();
}
```

⇒ YES

Does this method halt for *all* n?

```
boolean m(int n) {
  if (n<=1) return true;
  else if ((n%2) == 0) return m(n/2);
  else return m(3*n + 1);
}
```

Assuming infinite amount of stack space and arbitrarily large integers, it is *unknown* whether this program halts for all (strictly) positive integers!

Collatz Conjecture proposed in 1937!

Consider this Program called Q:

```
class HaltDetector {
  public static boolean halts(String javaProgram) {
    // ...do some super-clever analysis...
    // return true if javaProgram halts
    // return false if javaProgram does not
  }
}

class Main {
  public static void Q() {
     String p_Q = ???;   // string representing Q
     if (HaltDetector.halts(p_Q)) {
       while (true) {}  // infinite loop!
     }
  }
}
```

What happens when we run Q?

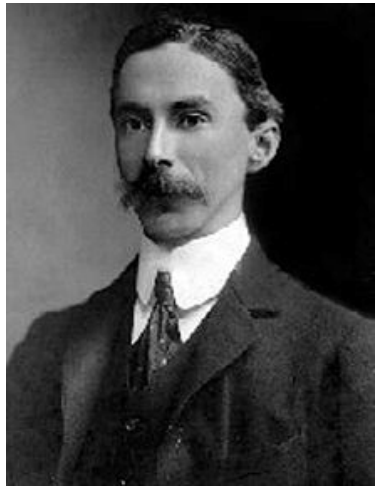
```
public static void Q() {
      String p_Q = ???;    // string representing Q
      if (HaltDetector.halts(p_Q)) {
        while (true) {}  // infinite loop!
      }
  }
```

if HaltDetector.halts(p_Q) ⇒ true then Q ⇒ infinite loop

if HaltDetector.halts(p_Q) ⇒ false then Q ⇒ halts

Contradiction!

- Russell's Paradox (1901)
- Gödel's Incompleteness Theorem (1931)
- Both rely on *self reference*

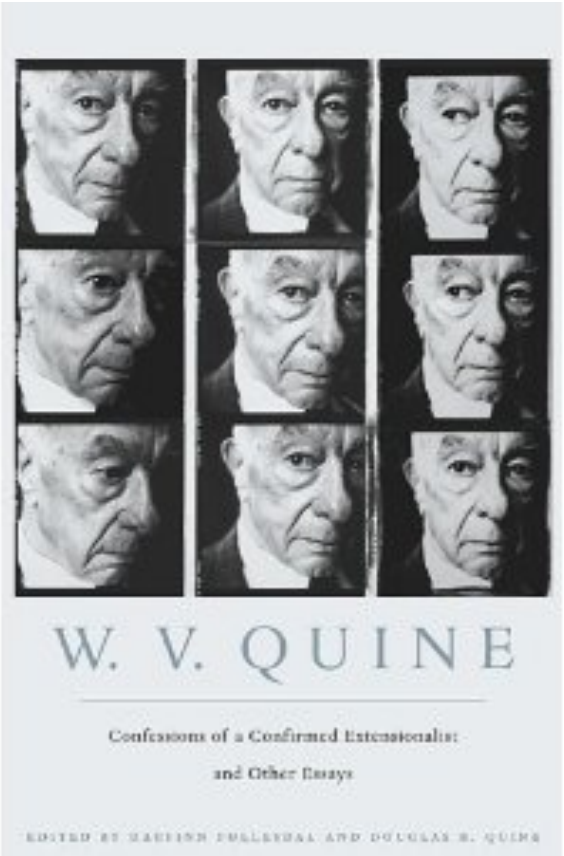


Bertrand Russell, 1901

Kurt Gödel, 1931

Potential Hole in the Proof

- What about the ??? in the program Q?
- It is supposed to be a String representing the program Q itself.
- How can that be possible?
- Answer: code is data!
  - And there's more than one representation for the same data.
- See Quine.java



Profound Consequences

- The "halting problem" is *undecidable*
  - *There are problems that cannot be solved by a computer program!*

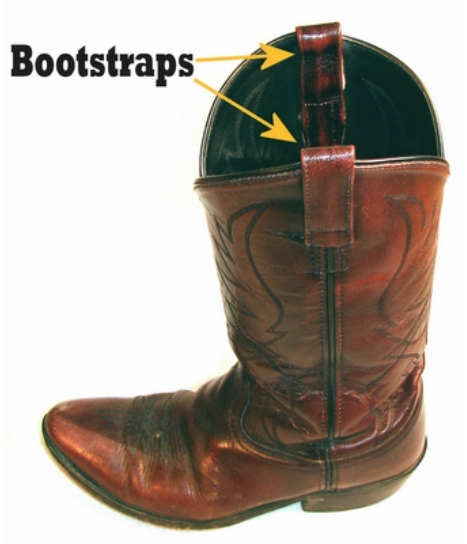
- Rice's Theorem:
  - Every "interesting" property about computer programs is undecidable!

- You can't write a perfect virus detector!
  (whether a program is a virus is certainly interesting)
  1. virus detector might go into an infinite loop
  2. it gives you false positives (i.e. says something is a virus when it isn't)
  3. it gives you false negatives (i.e. it says a program is not a virus when it is)

- Also: You can't write a perfect autograder!
  (whether a program is correct is certainly interesting)

Recommended Courses

- Programs that manipulate Programs
  - CIS 341: Compilers and interpreters
- Malware
  - CIS 331: Intro to Networks and Security
- Undecidability
  - CIS 262: Automata, Computability and Complexity



Recommended Reading