

Programming Languages and Techniques (CIS120)

Lecture 10

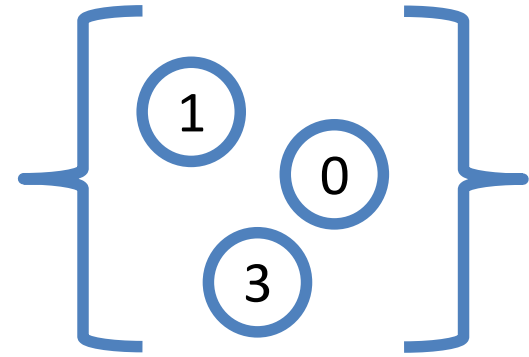
Abstract types: Sets

Chapter 10

Sets as Abstract Types

Mathematical Sets

Mathematical sets are *collections* of things:



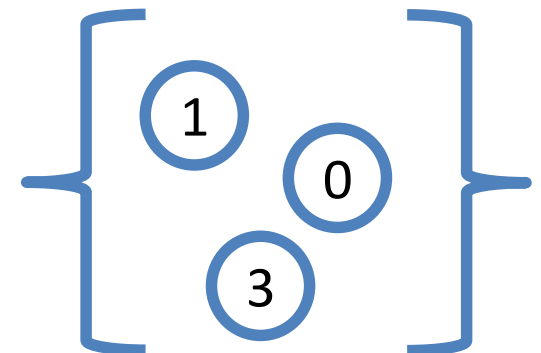
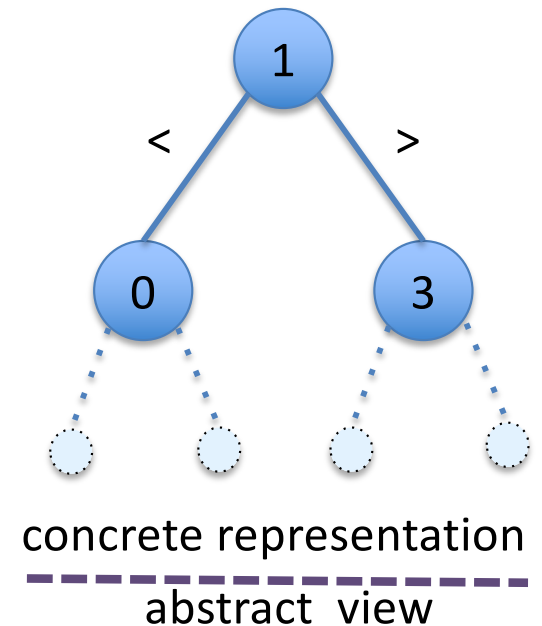
Empty Set:	\emptyset	no things
Nonempty Sets:	$\{0, 1, 2, 3\}$	four integers
	$\{(0,1), (2,3)\}$	two points in the plane
	$\{\text{true}, \text{false}\}$	two Boolean values
Set operations:		
	$S \cup T$	union
	$S \cap T$	intersection
Predicates:	$x \in S$	“x is a member of set S”

A *set* is a Collection

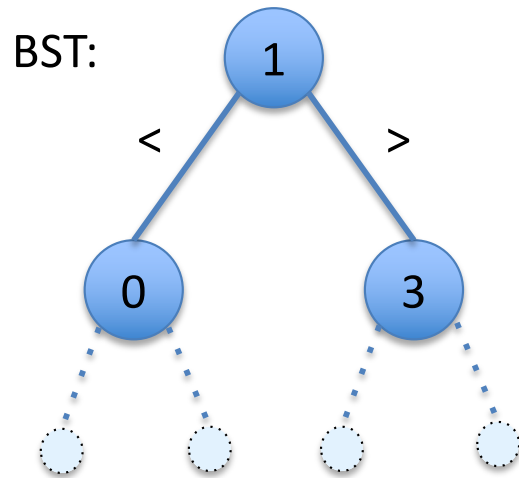
- A set is a *collection* of elements
 - we have operations for creating sets of elements
 - we can ask whether elements are in a set
 - Sets show up frequently in code
 - Examples: set of students in a class, set of coordinates in a graph, set of answers to a survey, set of data samples from an experiment, ...
 - A set is a lot like a list, except:
 - Order doesn't matter
 - Duplicates don't matter
 - *It isn't built into OCaml*
- An element's *presence* or *absence* in the set is all that matters...

A Set is an Abstraction

- A BST can *implement (represent)* a *set*
 - there is a way to represent an empty set (*Empty*)
 - there is a way to list all elements contained in the set (*inorder*)
 - there is a way to test membership (*lookup*)
 - Can define union/intersection (with insert and delete)
- *BSTs do not have to represent sets*
- *BSTs are **not the only** way to implement sets*

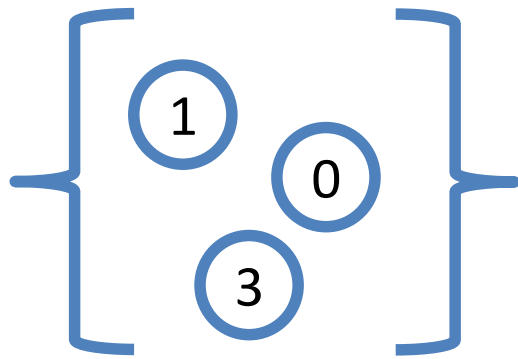


Three Example Representations of Sets



concrete representation

abstract view

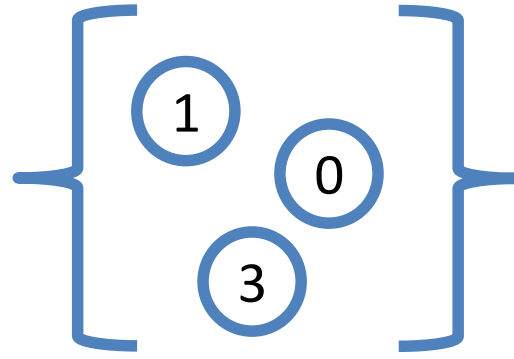


Alternate representation:
unsorted linked list.

3::0::1::□

concrete representation

abstract view

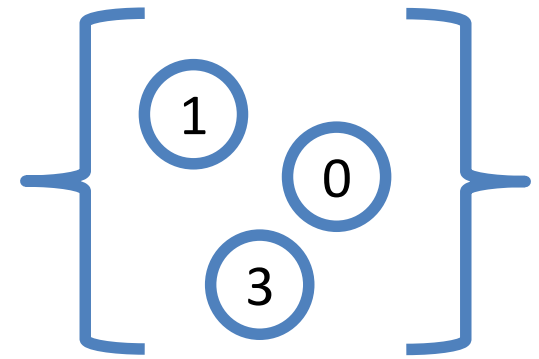


Alternate representation:
reverse sorted array with
Index of next slot.



concrete representation

abstract view



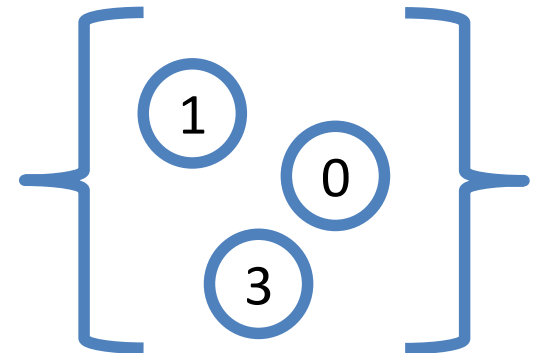
A Set is an Abstract Type

- An abstract type is defined by its *interface* and its *properties*, not its representation
- **Interface:** defines the type and operations
 - There is a type of sets
 - There is an empty set
 - There is a way to add elements to make a bigger set
 - There is a way to list all elements in a set
 - There is a way to test membership
- **Properties:** define how the operations interact with each other
 - Elements that were added can be found in the set
 - Adding an element a second time doesn't change the listing of elements
 - Adding elements in a different order doesn't change the listing of elements
- *When we use a set, we can forget about the representation!*



concrete representation

abstract view



Sets in OCaml

OCaml directly supports the declaration of abstract types via *signatures*

Set Signature

The name of the signature

The **sig** keyword indicates an interface declaration

```
module type SET = sig
```

```
  type 'a set
```

Type declaration has no “right-hand side” – its representation is *abstract*!

```
  val empty      : 'a set
```

```
  val add        : 'a -> 'a set -> 'a set
```

```
  val member     : 'a -> 'a set -> bool
```

```
  val equals     : 'a set -> 'a set -> bool
```

```
  val set_of_list : 'a list -> 'a set
```

```
end
```

The interface members are the (only!) means of manipulating the abstract type.

Signature (a.k.a. interface): defines operations on the type

Implementing sets

- There are many ways to implement sets
 - lists, trees, arrays, etc.
- *How do we choose which implementation?*
 - Depends on the needs of the application...
 - How often is ‘member’ used vs. ‘add’?
 - How big can the sets be?
- *How do we preserve the invariants of the implementation?*
- Many such implementations are of the flavor “a set is a ... with some *invariants*”
 - A set is a *list* with no repeated elements.
 - A set is a *tree* with no repeated elements
 - A set is a *binary search tree*

Invariant: a property that remains unchanged when a specified transformation is applied.

A module implements an interface

- An implementation of the set interface will look like this:

Name of the module

Signature that it implements

The **struct** keyword indicates a module implementation

```
module BSTSet : SET = struct
  ...
  (* implementations of type and operations *)
  ...
end
```

Implement the Set Module

```
module BSTSet : SET = struct
```

```
  type 'a tree =
```

```
    | Empty
```

```
    | Node of 'a tree * 'a * 'a tree
```

```
  type 'a set = 'a tree
```

```
  let empty : 'a set = Empty
```

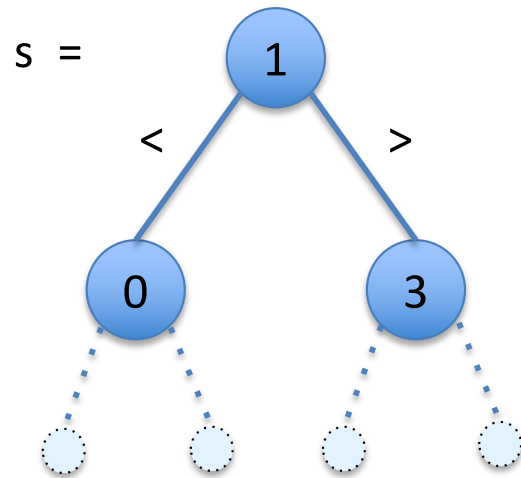
```
  ...
```

```
end
```

Module must define (give a *concrete representation* to) the type declared in the signature

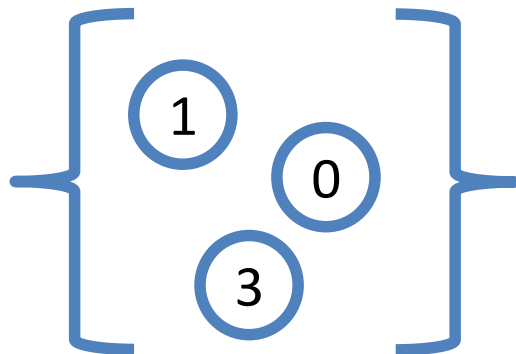
- The implementation must include everything promised by the interface
 - It can contain *more* functions and type definitions (e.g. auxiliary or helper functions) but those cannot be used outside the module
 - The types of the provided implementations must match the signature

Abstract vs. Concrete BSTSet



concrete representation

abstract view



```
module BSTSet : SET = struct
  type 'a tree = ...
  type 'a set = 'a tree
  let empty : 'a set = Empty
  let add (x:'a) (s:'a set) : 'a set =
    ... (* can treat s as a tree *)
```

end

```
-----
[ module type SET = sig
  | type 'a set
  | val empty : 'a set
  | val add : 'a -> 'a set -> 'a set
  | end
-----
```


```
(* A client of the BSTSet module *)
(* Cannot treat a set as a tree *)
;; open BSTSet
```

```
let s : int set
  = add 0 (add 3 (add 1 empty))
```

Another Implementation

```
module ULSet : SET =  
  struct  
    type 'a set = 'a list  
    let empty : 'a set = []  
    ...  
  end
```

A different definition for
the type set

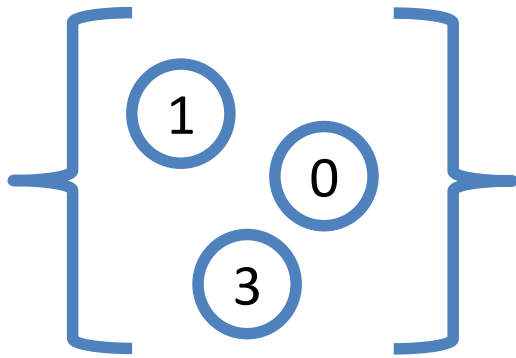


Abstract vs. Concrete ULSet

s = 0::3::1::[]

concrete representation

abstract view



```
module ULSet : SET = struct
  type 'a set = 'a list
  let empty : 'a set = []
  let add (x:'a) (s:'a set) : 'a set =
    x::s (* can treat s as a list *)
end
```

```
module type SET = sig
  type 'a set
  val empty : 'a set
  val add : 'a -> 'a set -> 'a set
end
```

```
(* A client of the ULSet module *)
(* Cannot treat a set as a list *)
;; open ULSet
```

```
let s : int set
= add 0 (add 3 (add 1 empty))
```

Client code doesn't change!

Testing (and using) sets

- Use “open” to bring all names defined in the interface into scope
- Any names that were already in scope are shadowed

```
;; open ULSet
```

```
let s1 = add 3 empty
```

```
let s2 = add 4 empty
```

```
let s3 = add 4 s1
```

```
let test () : bool = (member 3 s1)
```

```
;; run_test "ULSet.member 3 s1" test
```

```
let test () : bool = (member 4 s3)
```

```
;; run_test "ULSet.member 4 s3" test
```


Testing (and using) sets

- Alternatively, use the “dot” syntax:
`ULSet.<member>`
- Note: Module names must be capitalized in OCaml
- Useful when two modules define the same operations

```
let s1 = ULSet.add 3 ULSet.empty
let s2 = ULSet.add 4 ULSet.empty
let s3 = ULSet.add 4 s1
```

```
let test () : bool = (ULSet.member 3 s1)
;; run_test "ULSet.member 3 s1" test
```

```
let test () : bool = (ULSet.member 4 s3)
;; run_test "ULSet.member 4 s3" test
```

Does this code type check?

```
;; open BSTSet
let s1 : int set = add 1 empty
```

1. yes
2. no

```
module type SET = sig
  type 'a set
  val empty : 'a set
  val add    : 'a -> 'a set -> 'a set
end

module BSTSet : SET = struct
  type 'a tree =
    | Empty
    | Node of 'a tree * 'a * 'a tree
  type 'a set = 'a tree
  let empty : 'a set = Empty
  ...
end
```

Answer: yes

Does this code type check?

```
;; open BSTSet
let s1 = add 1 empty
let i1 = begin match s1 with
            | Node (_,k,_) -> k
            | Empty -> failwith "impossible"
          end
```

1. yes
2. no

```
module type SET = sig
  type 'a set
  val empty : 'a set
  val add    : 'a -> 'a set -> 'a set
end

module BSTSet : SET = struct
  type 'a tree =
    | Empty
    | Node of 'a tree * 'a * 'a tree
  type 'a set = 'a tree
  let empty : 'a set = Empty
  ...
end
```

Answer: no, add constructs a set, not a tree

Does this code type check?

```
;; open BSTSet
let s1 = add 1 empty
let i1 = size s1
```

1. yes
2. no

```
module type SET = sig
  type 'a set
  val empty : 'a set
  val add    : 'a -> 'a set -> 'a set
end

module BSTSet : SET = struct
  type 'a tree =
    | Empty
    | Node of 'a tree * 'a * 'a tree
  type 'a set = 'a tree
  let empty : 'a set = Empty
  let size (t : 'a tree) : int = ...
  ...
end
```

Answer: no, cannot access helper functions outside the module

Does this code type check?

```
;; open BSTSet  
let s1 : int set = Empty
```

1. yes
2. no

```
module type SET = sig  
  type 'a set  
  val empty : 'a set  
  val add    : 'a -> 'a set -> 'a set  
end  
  
module BSTSet : SET = struct  
  type 'a tree =  
    | Empty  
    | Node of 'a tree * 'a * 'a tree  
  type 'a set = 'a tree  
  let empty : 'a set = Empty  
  ...  
end
```

Answer: no, the Empty data constructor is not available outside the module

If a client module works correctly and starts with:

```
;; open ULSet
```

will it continue to work if we change that line to:

```
;; open BSTSet
```

assuming that ULSet and BSTSet both implement SET and satisfy all of the set properties?

1. yes
2. no

Answer: yes (though performance may be different)

```

module type SET = sig
  type 'a set
  val empty : 'a set
  val add    : 'a -> 'a set -> 'a set
  val member : 'a -> 'a set -> bool
end

module BSTSet : SET = struct
  type 'a tree =
    | Empty
    | Node of 'a tree * 'a * 'a tree
  type 'a set = 'a tree
  let empty : 'a set = Empty
  ...
end

```

Is it possible for a client to call **member** with a tree that is not a BST?

1. yes
2. no

No: the BSTSet operations preserve the BST invariants. there is no way to construct a non-BST tree using the interface.

Completing ULSet

See [sets.ml](#)

What Should You Test?

- **Interface:** defines operations on the type
- **Properties:** define how the operations interact
 - Elements that were added can be found in the set
 - Adding an element a second time doesn't change the elements of a set
 - Adding in a different order doesn't change the elements of a set



Test the properties!

A *property* is a general statement about the behavior of the interface: For *any* set *S* and *any* element *X*:

$$\text{member } x \text{ (add } x \text{ } s) = \text{true}$$

A (good) test case checks a specific instance of the property:

```
let s1 = add 3 empty
let test () : bool = (member 3 s1)
;; run_test "ULSet.member 3 s1" test
```

Property-based Testing

1. Translate informal requirements into general statements about the interface.

Example: “Order doesn’t matter” becomes

For *any* set s and *any* elements x and y ,
 $\text{add } x (\text{add } y s) \text{ equals } \text{add } y (\text{add } x s)$

2. Write tests for the “interesting” instances of the general statement.

Example. “interesting” choices:

$s = \text{empty}$, $s = \text{nonempty}$,

$x = y$, $x \neq y$

one or both of x, y already in s

Notes:

- one can’t (usually) exhaustively test all possibilities (too many!)
so instead, cover the “interesting” possibilities
- be careful with equality! `ULSet.equals` is *not* the same as `=`.

Abstract types: **BIG IDEA**

Hide the *concrete representation* of a type behind an *abstract interface* to preserve invariants

- The interface **restricts** how other parts of the program can interact with the data
 - Type checking ensures that the **only** way to create a set is with the operations in the interface
 - If all operations preserve invariants, then *all* sets in the program must satisfy invariants
 - Example: all BST-implemented sets must satisfy the BST invariant, therefore the lookup function can assume that its input is a BST
- Benefits
 - **Safety:** The other parts of the program can't violate invariants, which would cause bugs
 - **Modularity:** It is possible to change the implementation without changing the rest of the program

Summary: Abstract Types

- Different programming languages have different ways of letting you define abstract types
- At a minimum, this means providing:
 - A way to specify (write down) an interface
 - A means of hiding implementation details (*encapsulation*)
- In OCaml:
 - Interfaces are specified using a *signature* or *interface*
 - Encapsulation because the interface can *omit* information
 - type definitions
 - names and types of auxiliary functions
 - Clients *cannot* mention values or types not named in the interface