

Readings

- [Lecture Notes Chapter 24: Tries](#)

Review

A trie is a tree-based data structure that stores strings to support information retrieval. Tries are primarily useful when we need to repeatedly query a fixed text because it allows us to pre-process this text such that each subsequent query is fast, offsetting this initial cost of building the trie.

Standard Trie: In a standard trie, each root-to-leaf path corresponds to some string inserted into the trie. If the total length of all strings inserted into the trie is n , then a standard trie takes $O(n)$ time to build (using an incremental algorithm) and uses $O(n)$ space as well.

Patricia/Compressed Trie: A compressed trie is a trie where we guarantee that every internal node has at least two children by compressing branches/chains of single-child nodes into a supernode. If the total length of all strings is n and we have s strings, then a compressed trie takes $O(n)$ time to build but only uses $O(s)$ space, since the tree is now at least as full as a full binary tree (which has $O(s)$ nodes if it has s leaves).

Suffix Trie: A suffix trie is a trie where the strings are all the suffixes of a string S . Using an incremental algorithm, we can build a suffix trie in $O(|S|^2)$ time, but we can actually also do it in $O(|S|)$ time using Ukkonen's Algorithm; however, the details behind how this works are outside the scope of CIS 121. A compressed suffix trie uses $O(|S|)$ space.

Problems

Problem 1

Given a set of N strings, design an efficient algorithm to find the longest common prefix between any two strings. What is the running time of your algorithm?

Problem 2

Given some string S , design an efficient algorithm to find the longest repeated substring. What is the running time of your algorithm?