

CMOS Design, Gates & PLAs

Introduction to Computer Systems, Fall 2022

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Logistics

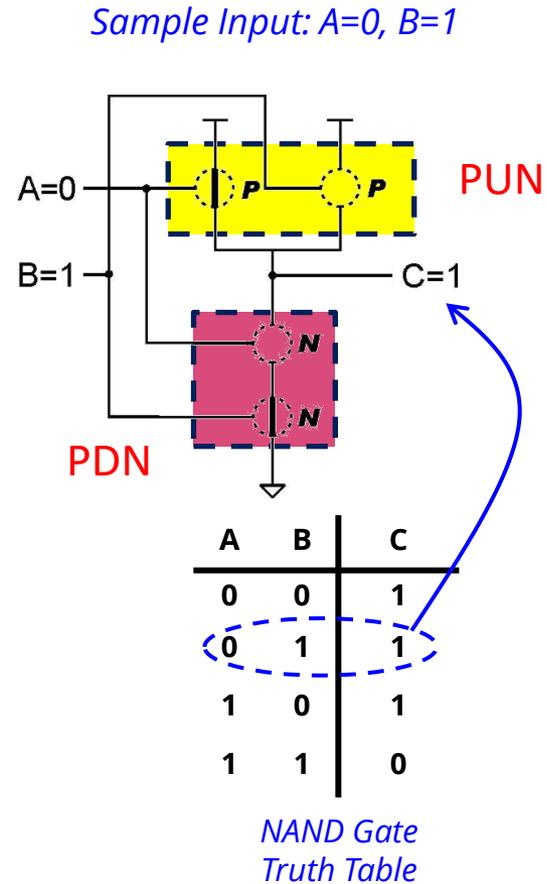
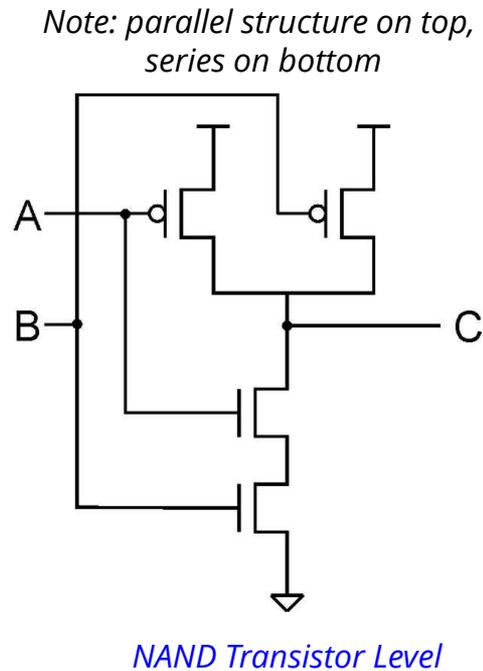
- ❖ HW01 bits.c: **This Friday** 9/16 @ 11:59 pm
 - Will require VM setup
 - Has you “program” in C
 - Should have everything you need
 - Terminal & starting demo in Recitations this week

- ❖ HW02 Combinational Logic: to be released this week

Lecture Outline

- ❖ **CMOS Circuit Design**
- ❖ Gates & PLAs
- ❖ Transistors in the Real World

The NAND (NOT-AND) Circuit



Rules & Suggestions For Design

❖ Rules:

- You cannot have pMOS transistors in the PDN
- You cannot have nMOS in the PUN
- Exactly one of PDN/PUN must be “on” at a time
 - Cannot have neither connected to output
 - Cannot have both connected to output
- Every transistor in the PDN must have a complimentary transistor in the PUN
 - When any transistor in PDN is ON, its compliment transistor is OFF

❖ Suggestions

- Start with the PDN and then do the PUN (most find it easier)
- Simplify logic before you design any part of the circuit

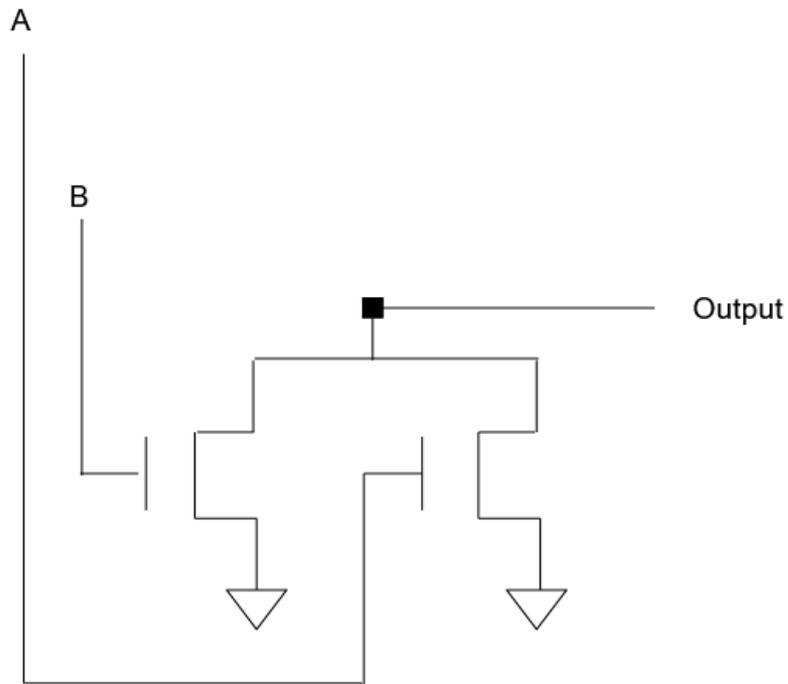
NOR Circuit Design Walkthrough

- ❖ First, start with the Boolean Expression
 - NOR is NOT-OR, which is
$$\sim(A \mid B) \leftarrow \text{statement for when output is 1 (True)}$$
- ❖ Since I like to start with PDN, find the cases when output is FALSE (output connected to GND), then simplify
 - This can be done by negating the original expression
 - $\sim\sim(A \mid B) \leftarrow \text{statement for when output is 0 (False)}$
 - $(A \mid B) \leftarrow \text{simplified by identity property}$

NOR Circuit Design Walkthrough

- ❖ With the expression for the PDN (when output is 0), translate it to a circuit diagram for the PDN
 - OR's become transistors in parallel
 - AND's become transistors in series

$(A \mid B)$ becomes

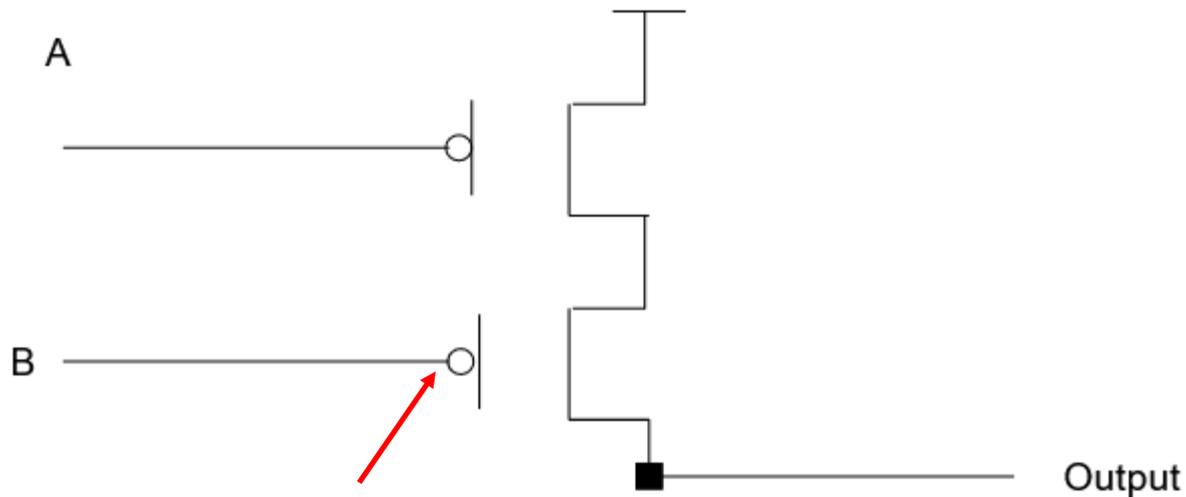


NOR Circuit PUN (Strategy 1)

- ❖ With the simplified expression for PDN, negate it to get the expression for PUN
 - $(A \mid B)$ becomes $\sim(A \mid B)$
 - Simplify
 - $\sim(A \mid B)$
 - $(\sim A \ \& \ \sim B)$ // by De Morgan's Law
- In this example, initial PUN expression is the same as the original expression. This is not always the case*
- ❖ From here, we can convert $(\sim A \ \& \ \sim B)$ into a PUN

NOR Circuit PUN (Strategy 1)

- ❖ From here, we can convert ($\sim A$ & $\sim B$) into a PUN
 - OR's become transistors in parallel
 - AND's become transistors in series
 - pMOS transistors have an implicit "NOT"
 - pMOS only turns on when input is low (0)

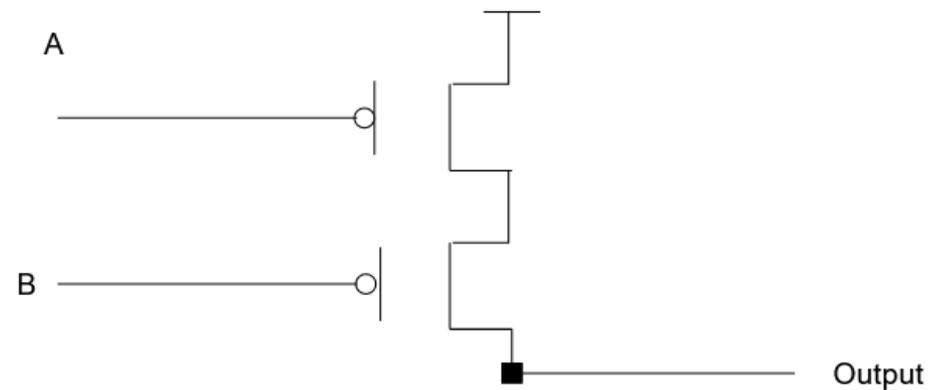
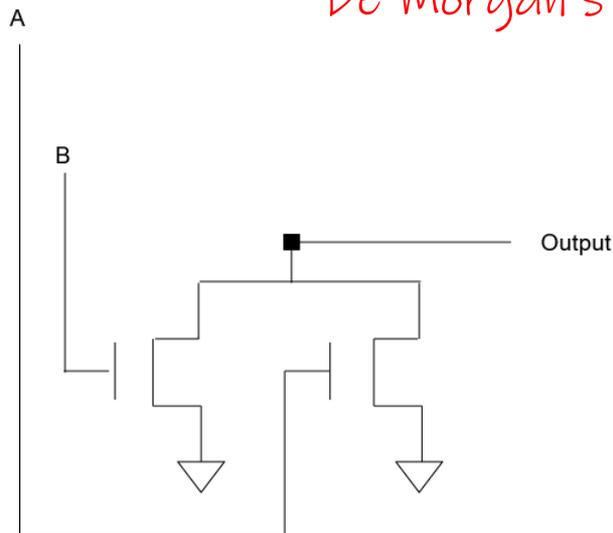


The reason for the circle in pMOS:
 Circle represents a "Not" or negation

NOR Circuit PUN (Strategy 2)

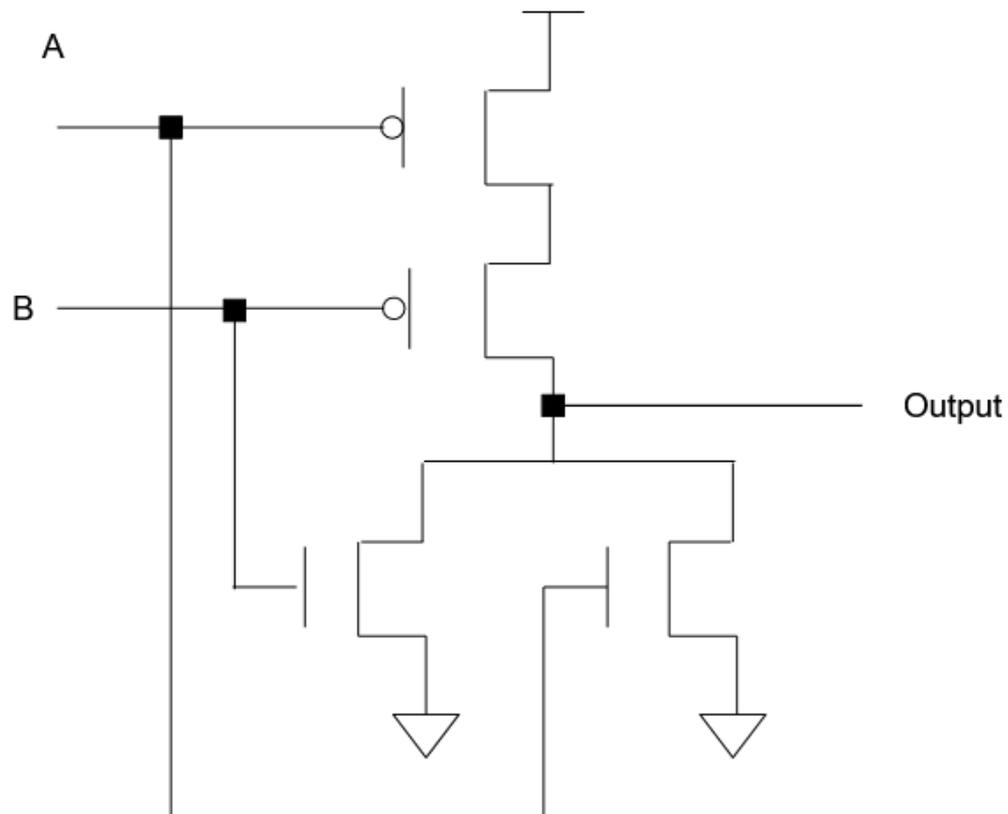
- ❖ Alternative way to get the PUN from the PDN
 - Series relations in PDN become Parallel relations in PUN
 - Parallel relations in PDN become Series relations in PUN
 - nMOS transistors become pMOS transistors

This is just a "short-cut" for applying De Morgan's to the PDN



Completed NOR Circuit

- ❖ Finally, all you have to do is combine the PDN and PUN
 - To be safe, check your work and create a truth table



A	B	Output
0	0	1
0	1	0
1	0	0
1	1	0

Complex Practice

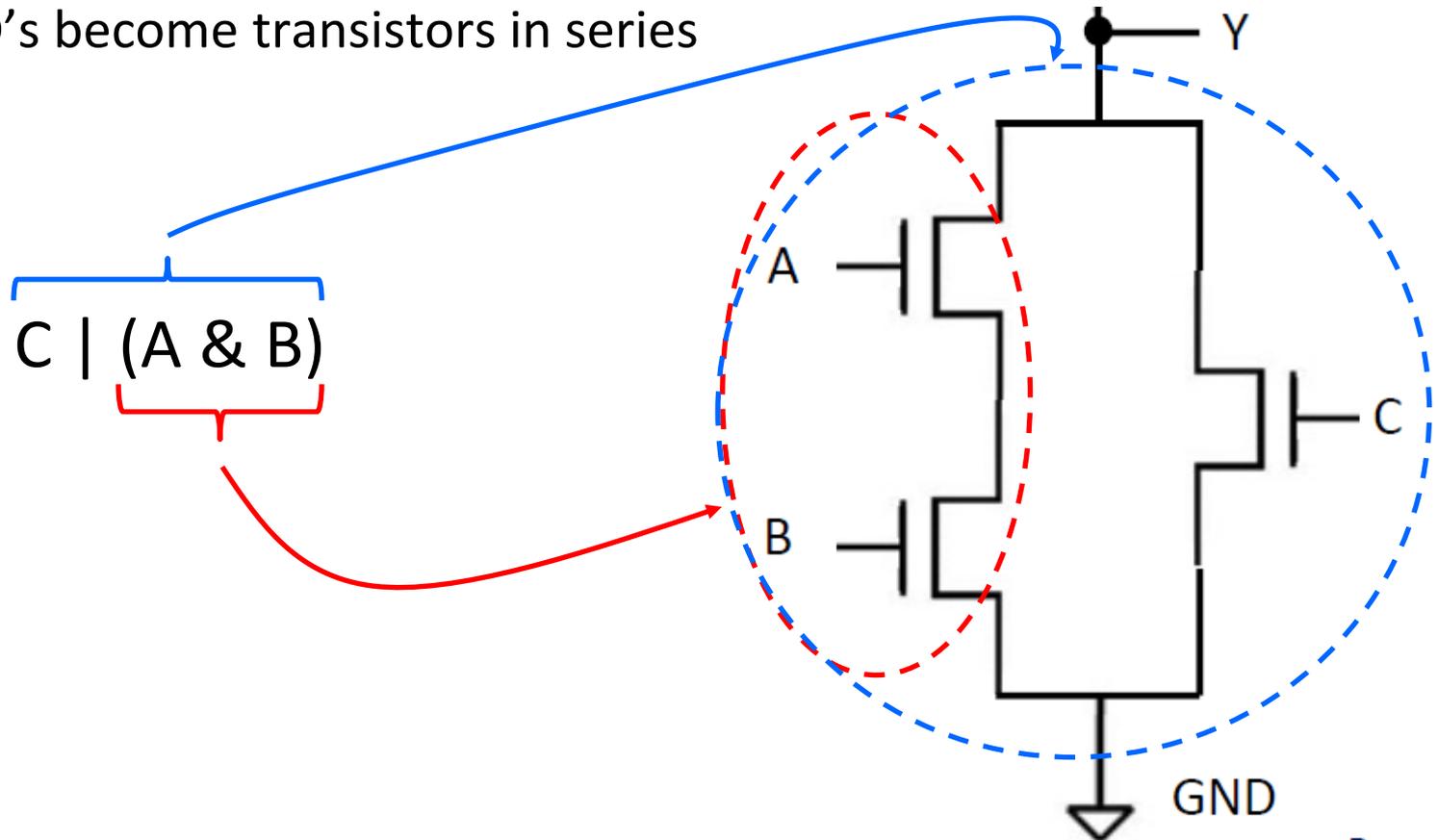
- ❖ Let's Practice, create the CMOS circuit that takes in 3 inputs (A, B, C) and has a high output (1) if C is 0 and at least one of A or B is 0.
- ❖ First, what is the expression?
 - $\sim C \& (\sim A \mid \sim B)$

Complex Practice

- ❖ Second, what is the expression for the PDN?
 - $\sim C \& (\sim A \mid \sim B)$ // original overall expression
 - $\sim(\sim C \& (\sim A \mid \sim B))$ // Negate for PDN expression
 - $\sim\sim C \mid \sim(\sim A \mid \sim B)$ // by De Morgan's
 - $C \mid \sim(\sim A \mid \sim B)$ // By identity property
 - $C \mid (\sim\sim A \& \sim\sim B)$ // by De Morgan's
 - $C \mid (A \& B)$ // by identity property

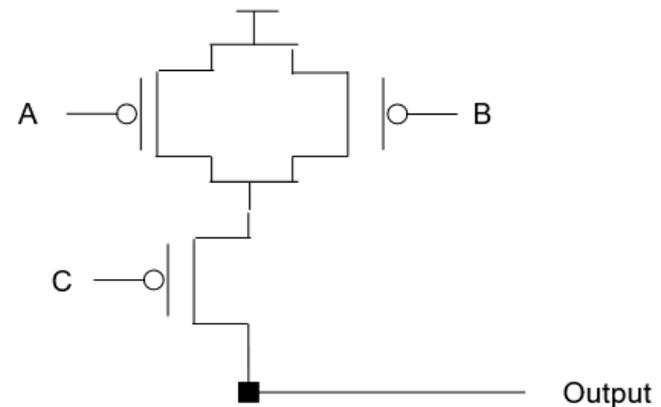
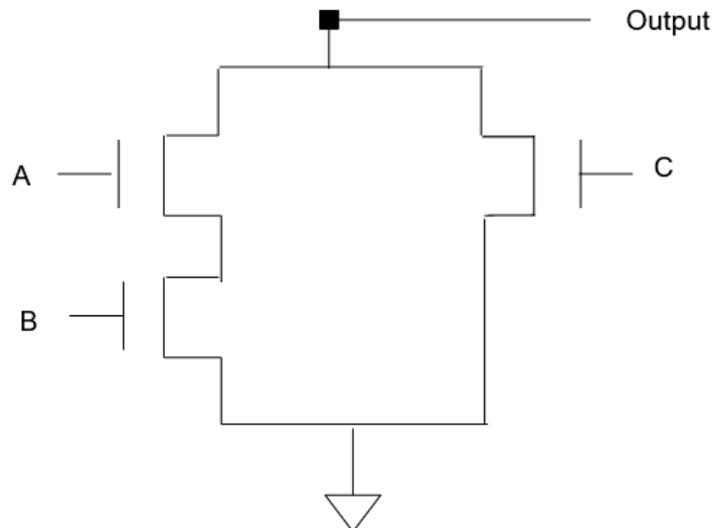
Complex Practice

- ❖ Third, convert the PDN expression into a circuit
 - OR's become transistors in parallel
 - AND's become transistors in series



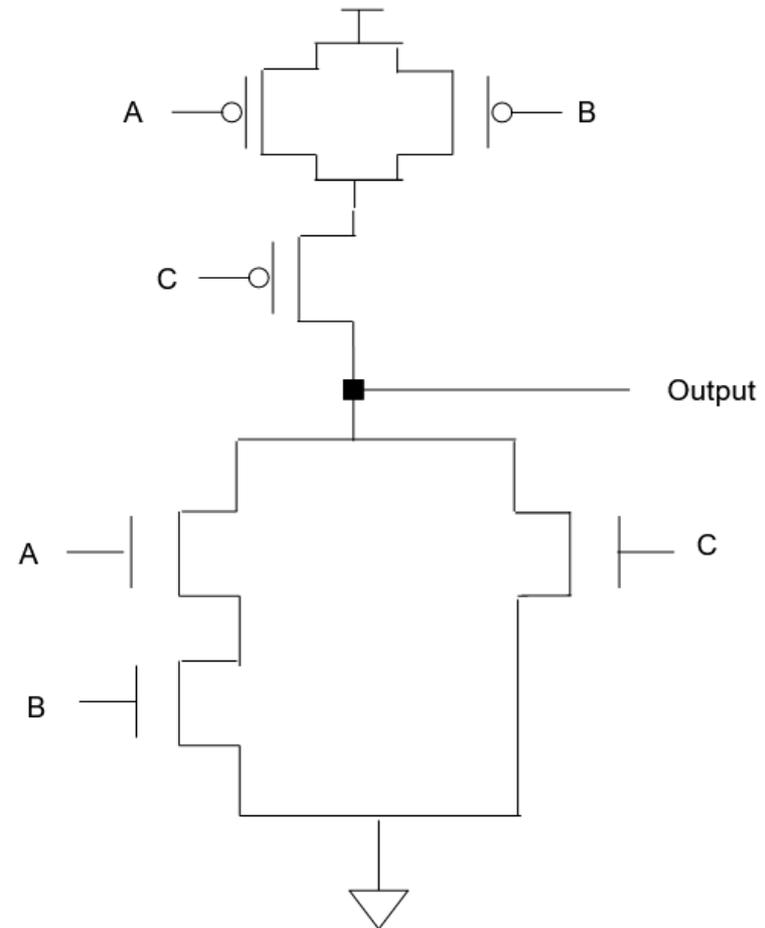
Complex Practice

- ❖ Fourth get the PUN from the PDN (I'll use the short-cut)
 - Series relations in PDN become Parallel relations in PUN
 - Parallel relations in PDN become Series relations in PUN
 - nMOS transistors become pMOS transistors



Complex Practice

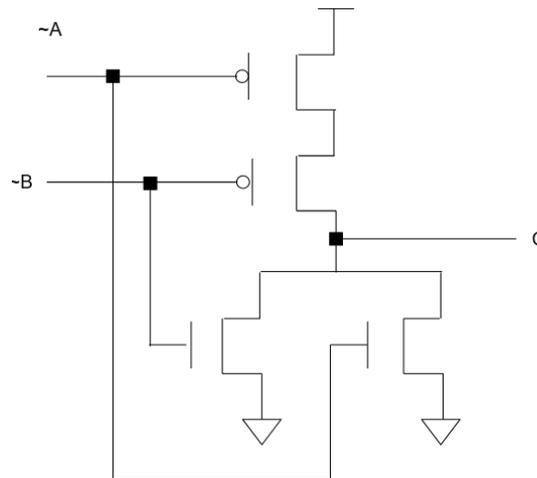
- ❖ Last step:
combine PDN and PUN
- ❖ Suggested:
Check your work



Closing Details on CMOS

- ❖ Sometimes, there will be 3 or more inputs (A, B, C)
- ❖ Sometimes you will also have the complements of the inputs ($\sim A$, $\sim B$, $\sim C$) that you can use as inputs in a circuit

- Example AND circuit:



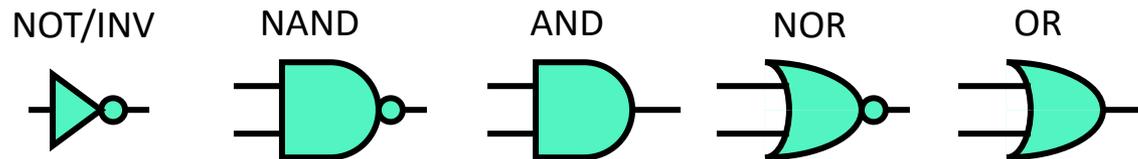
- ❖ Highly suggest you simplify before creating circuits
 - More practice available in recitation

Lecture Outline

- ❖ CMOS Circuit Design
- ❖ **Gates & PLAs**
- ❖ Transistors in the Real World

Logic Gates

❖ Basic Logic Gates

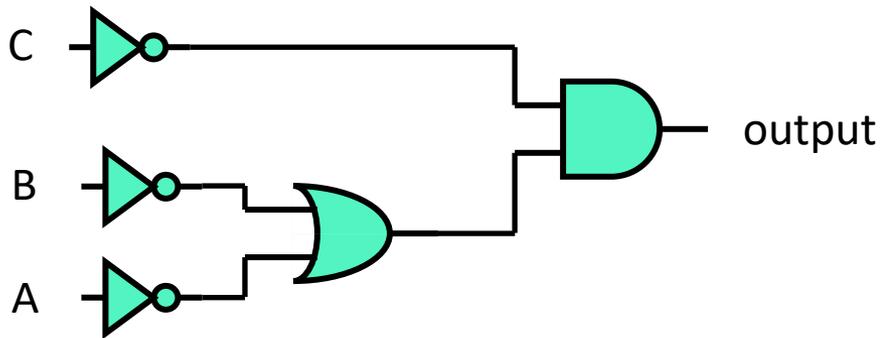


❖ Moving an Abstraction layer “up”

- + Abstracting away messy details of CMOS Circuits!
 - + Technology independent at this level, doesn't have to be CMOS
 - - Designing only at this level can be less efficient than CMOS
- ❖ Gates of more complex operations exist (e.g. an XOR gate)
- ❖ Gates can be combined to make more complex functions

Logic to Gates Practice

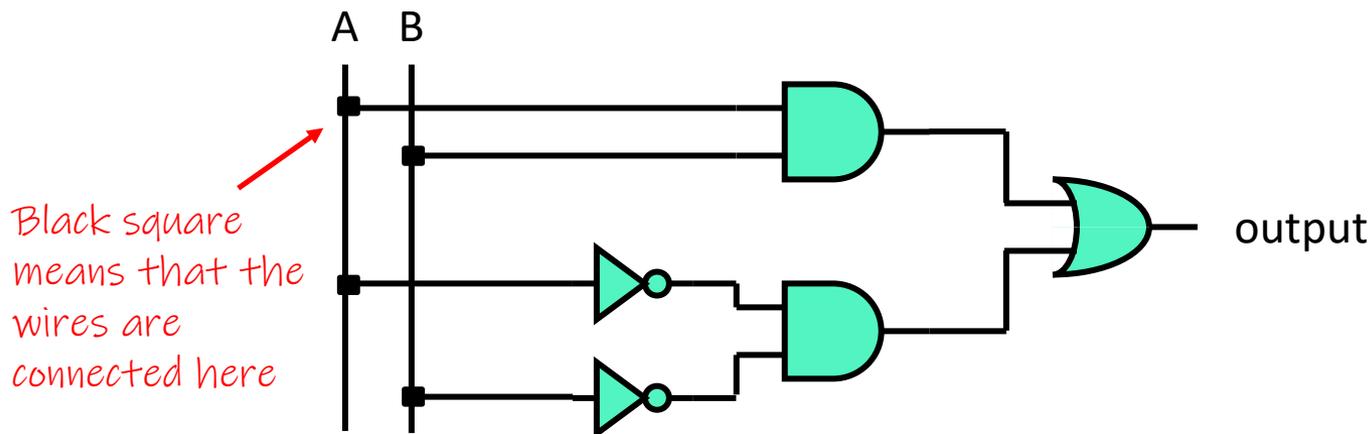
- ❖ Let's Practice, create a Gate-level circuit that takes in 3 inputs (A, B, C) and has a high output (1) if C is 0 and at least one of A or B is 0.
- ❖ First, what is the expression?
 - $\sim C \ \& \ (\sim A \ | \ \sim B)$
- ❖ Second, translate it to gates:



Logic to Gates Practice 2

- ❖ Create a gate-level circuit that takes two inputs (A and B) and output 1 if and only if A and B are the same.
 - Use only the basic logic gates I've shown you so far

- ❖ First, what is the expression?
 - $(A \& B) \mid (\sim A \& \sim B)$

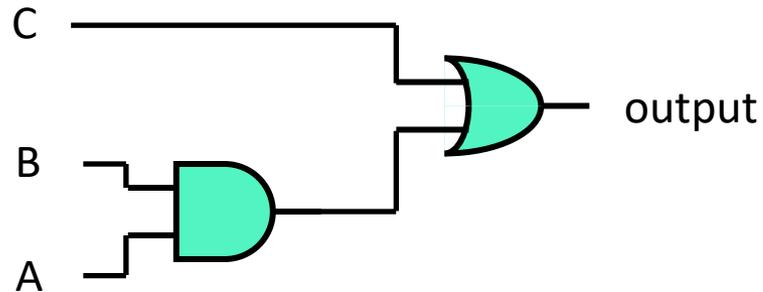


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❖ What is the equivalent expression of this gate circuit

- A. $(A \& B) | C$
- B. $C \& (B | A)$
- C. $(C | B) \& A$
- D. $(B \& C) | A$
- E. I'm not sure



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❖ What is the equivalent expression of this gate circuit

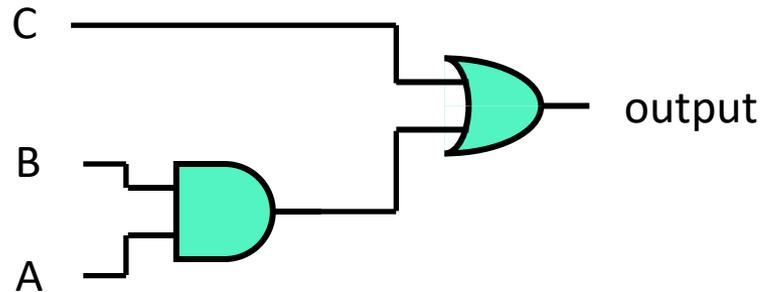
A. $(A \& B) | C$

B. $C \& (B | A)$

C. $(C | B) \& A$

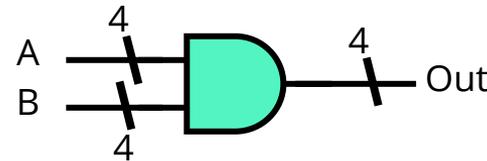
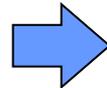
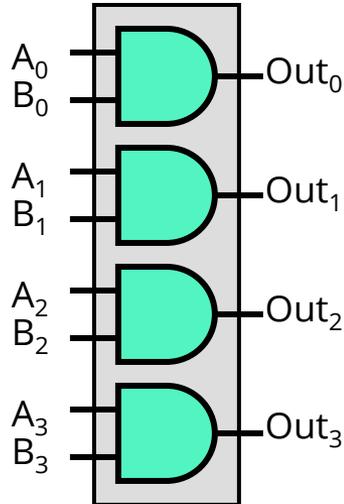
D. $(B \& C) | A$

E. I'm not sure



Shortcuts: Multi-Bit Gates

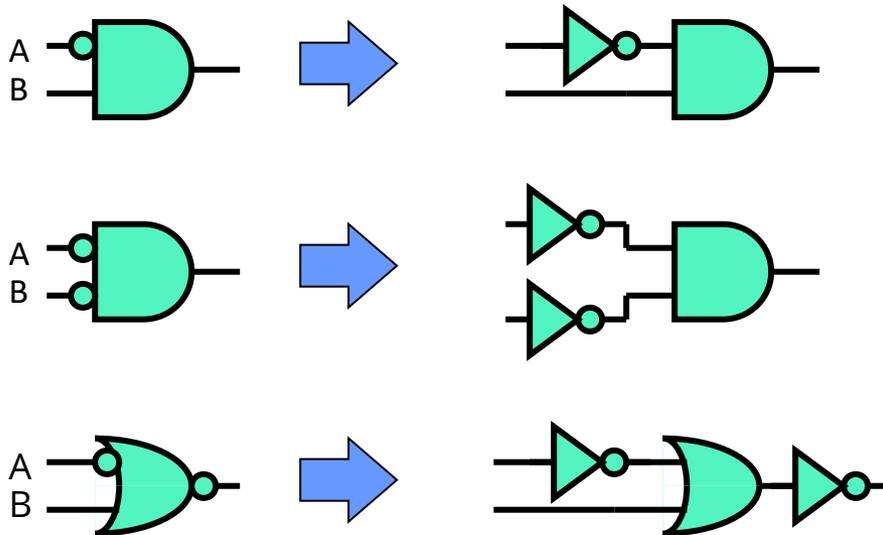
- ❖ Use a cross-hatch mark groups of wires:
 - Example, Say we had two 4-bit integers we wanted to AND
 - A_3 is the most-significant bit, A_0 is the least significant bit



This is still a 2-input NAND,
Except each input is 4-bits wide!

Shortcuts: Inverting Signals

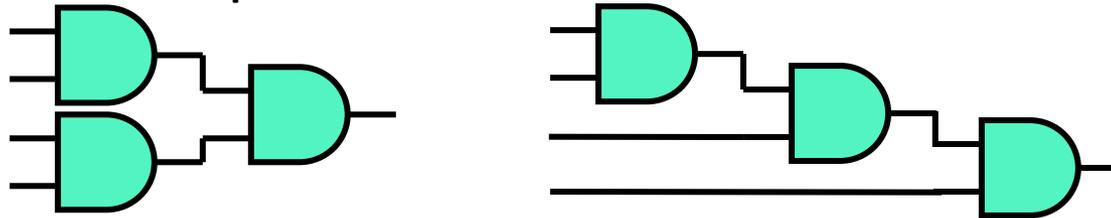
- ❖ Instead of always drawing out a NOT gate
 - A  represents a negation/inversion of a signal



Boolean algebra rules on gates

- ❖ Boolean rules still apply to gate-level circuits

- ❖ Associative Example:



- ❖ De Morgan's law Example:



- ❖ **NOTE:** There are actual differences between these circuits, but “logically” they are the same. More on this later in lecture.

PLA's

- ❖ What if we only had a truth table to create a gate circuit?
- ❖ PLA: **P**rogrammable **L**ogic **A**rray
 - A device where we can configure AND, OR and NOT gates to implement a function

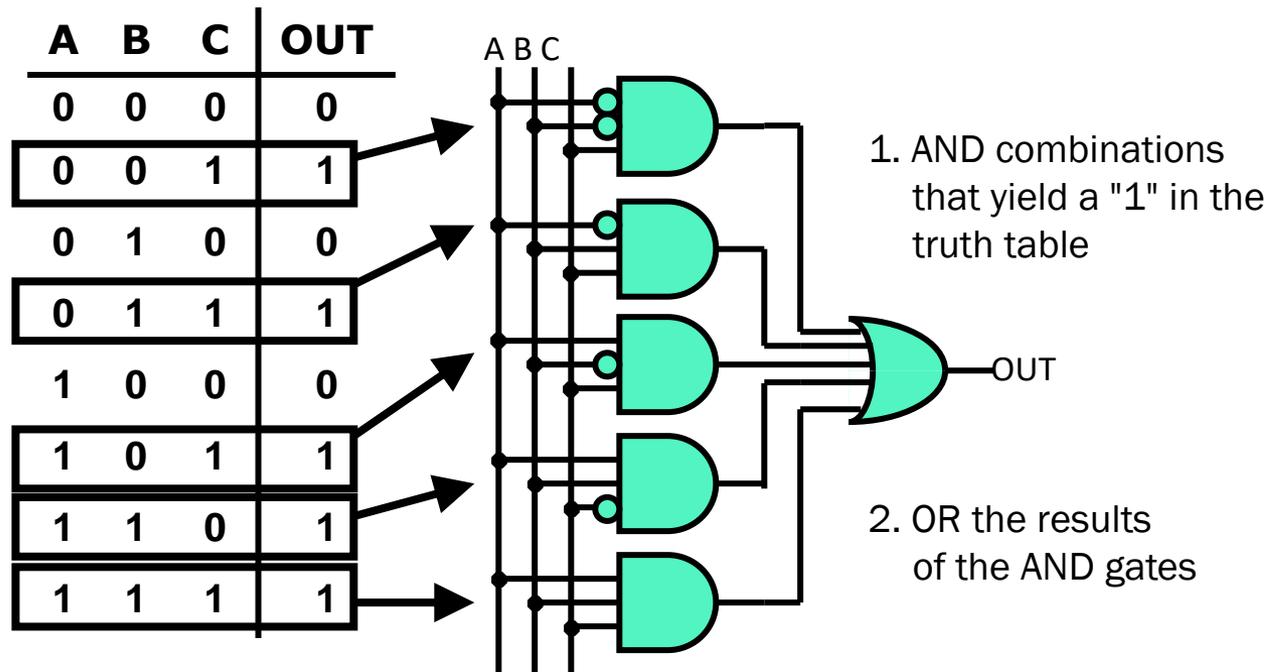
A	B	C	OUT
0	0	0	0
0	0	1	1
0	1	0	0
0	1	1	1
1	0	0	0
1	0	1	1
1	1	0	1
1	1	1	1



GATES

Implementing a PLA From a Truth Table

- ❖ NOT, AND, OR can implement any truth table function



IN a PLA, this structure is always followed

Notice, 5 rows that cause a "1" in the output...5 AND gates
 Notice, 1 output, only 1 OR gate
 Notice, negations always happen before the AND gates

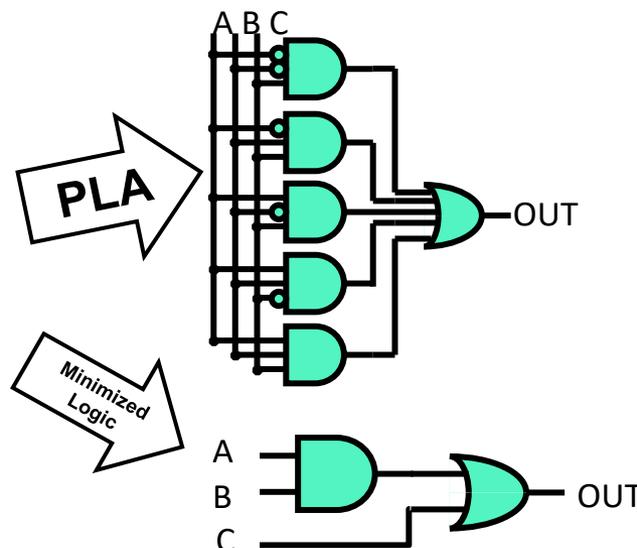
PLAs Pros & Cons

- ❖ A PLA can be used to implement ANY logical function
 - Provides you with an incredibly easy tool to use
 - If you can generate a truth table to model desired behavior
 - PLA gives you a way generate the gate level implementation
 - *However, PLAs don't give the most efficient solution*
 - *In terms of "run-time" and transistor cost*

Logic Function
 $F = (A \text{ AND } B) \text{ OR } C$

Truth Table

A	B	C	OUT
0	0	0	0
0	0	1	1
0	1	0	0
0	1	1	1
1	0	0	0
1	0	1	1
1	1	0	1
1	1	1	1



Lecture Outline

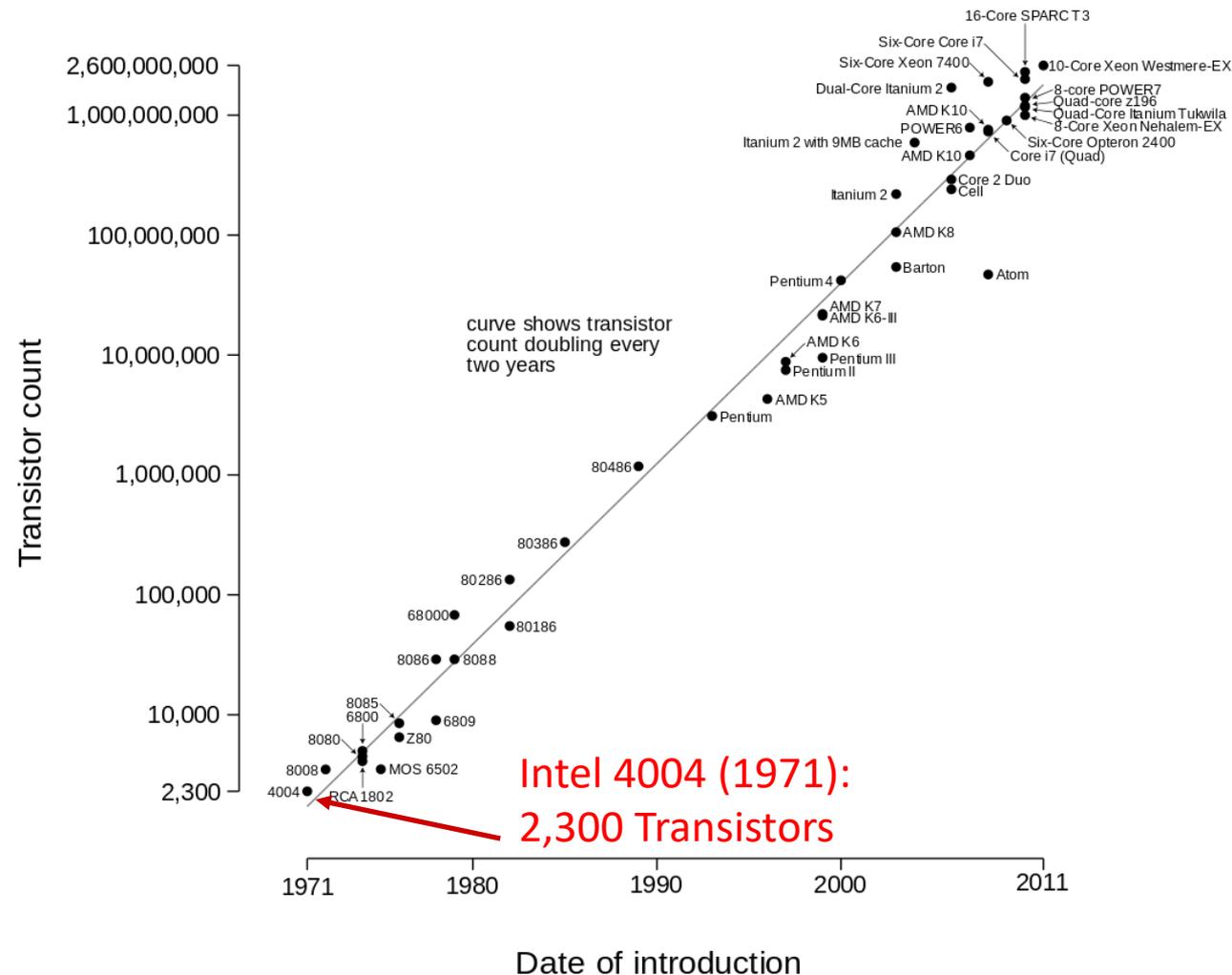
- ❖ CMOS Circuit Design
- ❖ Gates & PLAs
- ❖ **Transistors in the Real World**

Moore's Law

- ❖ Observation that number of transistors that can be put in an integrated circuit doubles every 18 months-ish

Intel 15-core Xeon IvyBridge-EX (2014):
4.3 billion
(not pictured)

Microprocessor Transistor Counts 1971-2011 & Moore's Law



Transistors are Small

- ❖ The Intel Skylake chip uses transistors that are 60 times smaller than the wavelength of light.
 - Skylake transistor size 14nm.
 - Wavelength of visible light 400-700nm

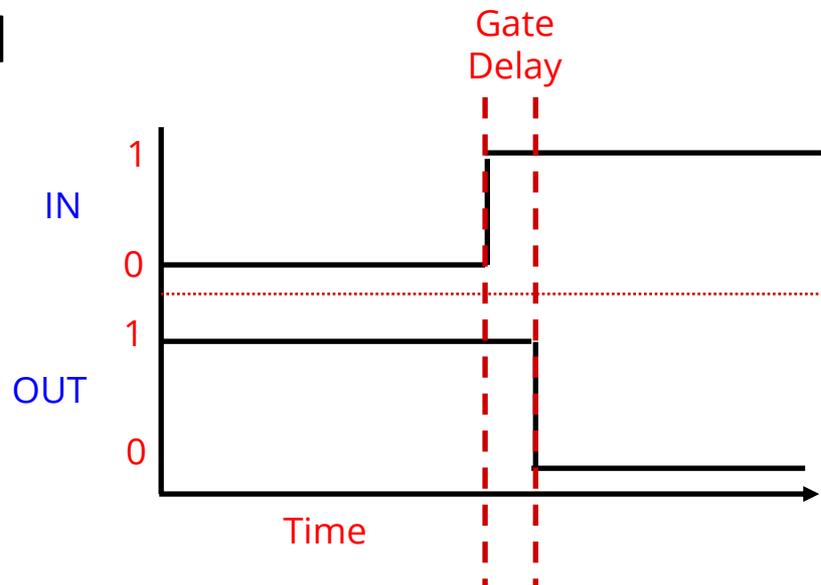
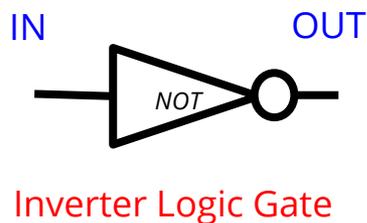
- ❖ There are now more transistors at work in the world (15 quintillion : 15,000,000,000,000,000,000) than there are leaves on all the trees in the world
 - From (The Perfectionists, Simon Winchester)

Transistors limitations

- ❖ After switching an input, it takes time for current to flow, and output to match changed output
- ❖ The faster you switch the circuit, the more current flows, the more heat is generated, the hotter your laptop gets.
 - This has proven to be an important barrier to speeding up CMOS circuitry

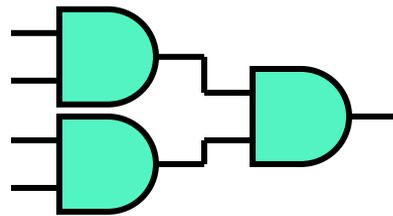
Gate Delays

- ❖ With any logic circuit there will be a short delay between when an input changes and when the output updates to it
 - This time is referred to as the gate delay
- ❖ For modern circuitry, gate delays are on the order of nanoseconds (10^{-9} seconds) or picoseconds (10^{-12} seconds)
- ❖ These delays ultimately limit the number of operations you can perform per second

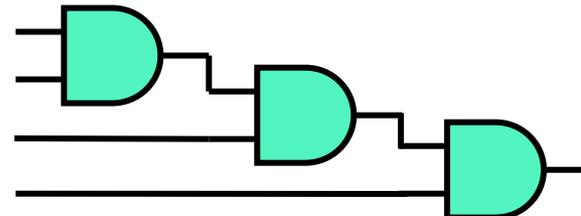


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- ❖ Consider the following logically equivalent circuits:



Circuit A



Circuit B

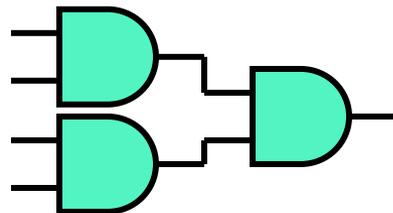
Which is the better implementation of 4-input AND?

- A. **Circuit A**
- B. **Circuit B**
- C. **They are equally good**
- D. **I'm not sure**

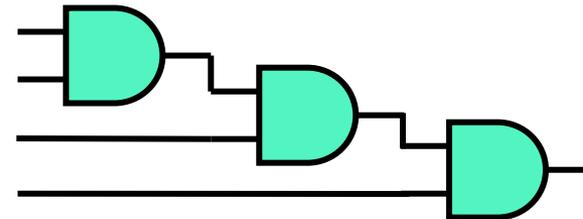
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- ❖ Consider the following logically equivalent circuits:



Circuit A



Circuit B

Which is the better implementation of 4-input AND?

A. Circuit A

Circuit A is better:

Why? It's faster, 2 "gate delays" instead of 3

B. Circuit B

C. They are equally good

- ❖ **Gate delays: longest path (in gates) through a circuit**

- Grossly over-simplified, ignores gate differences, wires
- Good enough for our purposes

D. I'm not sure

Abstraction

- ❖ Moving from CMOS to Gates is going to a higher abstraction level.

- ❖ Transistor Level:
 - More control over how many transistors there are and how they are oriented (orientation could affect how much delay there is)
 - More hardware details that must be worried about (PDN vs PUN, pMOS vs nMOS, etc)

- ❖ Gate level:
 - Abstract away the hardware details, focus only on logic.
 - Multiple gates -> longer delay

Multiple Approaches to things

- ❖ Lots of ways to implement things. Implementation choice depends on many factors:
 - Ease of Implementation (Readability, dev time, ...)
 - Kinds of gates available
 - Number of transistors available
 - Energy Consumption
 - Circuit delay
 - ...

Lecture Schedule

❖ This Lecture

- CMOS
- Gates
- PLAs
- Gate Delays

❖ Next Lecture:

- Use Gates & PLAs to start building more complex circuits and components of a computer