

Java and C tips

Intro to Computer Systems, Fall 2022

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Upcoming Due Dates

- ❖ HW10/11 (J compiler) is due Friday December 9th
 - HW10 & 11 make up a 2-part assignment that take a while to complete.
 - Recitation for this assignment has been VERY helpful
 - Can grant extensions on this, but there will be reduced office hours and Ed activity after a bit
 - **Took some students a long time in Fall 2021**

- ❖ Final Exam: Thursday December 15th
 - Cumulative
 - More info coming soon

J Compiler Common Mistakes

- ❖ DON'T FORGET TO ADD HEADER GUARDS

- ❖ `next_token`
 - When you read a comment, don't forget to read till the rest of the line

- ❖ ASM generation:
 - Some 16-bit LITERALS require both CONST and HICONST to load that value into a register
 - The prologue/epilogue is wrong, you can mostly copy this off of the slides though.
 - Generating unique labels/handling nested control structures

Lecture Outline

- ❖ **Java vs C**
 - **Java Datatypes**
 - Java Compilation
 - Java Garbage Collector
- ❖ C tips & Practice

None of this is on
the final exam or
HW10/HW11

Comparing Java and C

- ❖ Perquisite to this course: CIS 1100
 - You all have experience programming in Java
 - Java the first language for most of you
- ❖ "The Hardest programming language you learn is the second one that you learn."
 - May not fully be true, but it is common to struggle with the differences between the languages
 - Doesn't help that C and Java look VERY similar
- ❖ Hopefully this comparison gives you a better understanding of both Java and C

Disclaimer

- ❖ Java and C both can have multiple implementations.
 - Some things we discuss in this lecture may not be guaranteed, but instead may vary.
- ❖ C: Leaves some details that can vary from machine to machine and/or compiler to compiler
 - Example: what is the size and sign of the **char** datatype?
 - Example: What happens when we return the address of data in the current stack frame?
- ❖ Java: the language specification provides an abstraction
 - We can understand how the code should behave, but it may do things differently when actually compiled & run

Java Data Types: Primitives

- ❖ Primitive types are pretty much the same as C
 - `int`, `float`, `double`, etc.
 - Java doesn't have **unsigned** types to avoid issues with converting & comparing between signed/unsigned types
- ❖ `char`:
 - **char** in C is 1-byte which represents an ASCII character
 - **char** in Java is 2-bytes for 2-byte Unicode characters
- ❖ Primitive Size:
 - Java is designed to be portable, primitives are fixed in size
 - C primitive sizes can vary from machine to machine
 - Example: `int` **is 4-bytes** in Java, and **is usually 4-bytes** in C

Java Data Types: Pointers

- ❖ Pointers are a type of primitive in C. Can be used to access memory but we can also deal with the address directly (pointer arithmetic, get address of with `&`)
- ❖ Java has references, which are almost like "protected" or "hidden" pointers.
 - All Object variables are actually Object References
 - Much more constrained in how you use them to try and minimize possible memory usage errors
- ❖ Both have `NULL` or `null` to indicate an unused/empty pointer/references. (`NULL` typically represented as 0)

Java Data Types: Objects

- ❖ C doesn't have true Objects, but code can have "objects" or structs. This data can exist in many places in memory.
- ❖ Java has Object support. All objects in Java are stored on the heap. The "new" keyword allocates memory dynamically, like how `malloc` allocates space.

Java Data Types: Objects Example

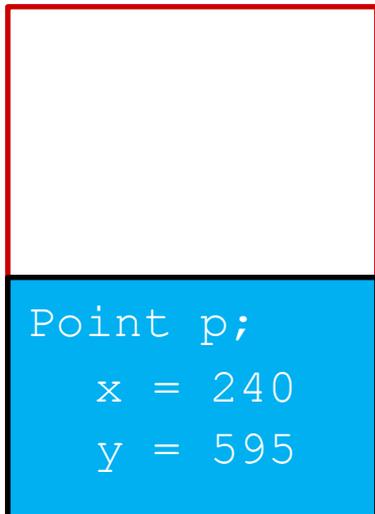
- ❖ Consider we have a struct **Point** in C and object **Point** in Java. Each contains two integers, an X and a Y.

C

```
int main() {
    Point p;
    p.x = 240;
    p.y = 595;
}
```

Java

```
public static void main(String args[]) {
    Point p;
    p.x = 240;
    p.y = 595;
}
```



Stack frame for
main()

NULL POINTER EXCEPTION

Point p is an uninitialized references
(an uninitialized pointer)

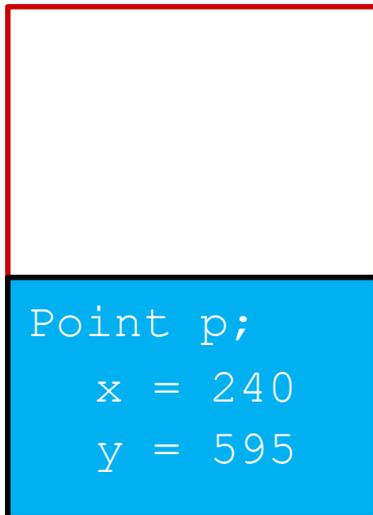
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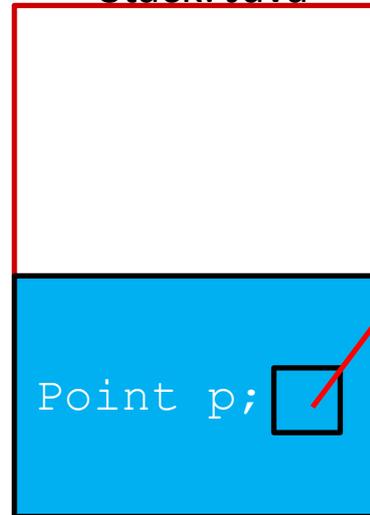
Stack: C


 Stack frame for
main()

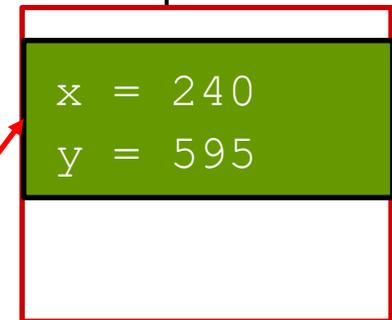
Java

```
public static void main(String args[]) {
    Point p = new Point();
    p.x = 240;
    p.y = 595;
}
```

Stack: Java


 Stack frame for
main()

Heap: Java



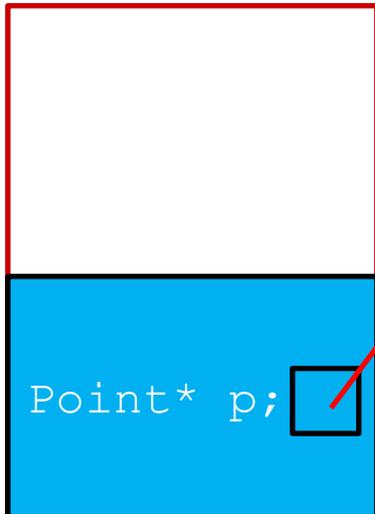
C Objects Heap Example

- ❖ C can also have “references” to things on the heap, but it is more explicit in the code

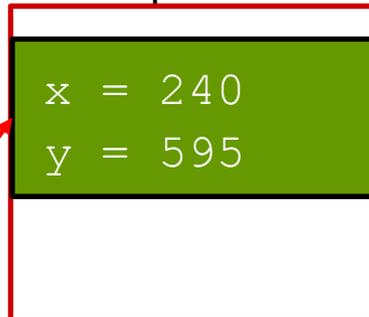
C

```
int main() {
    Point* p = malloc(sizeof(Point));
    p->x = 240;
    p->y = 595;
}
```

Stack: C



Heap: C



Stack frame for
main()

Java Data Types: Arrays

❖ C Arrays:

- elements are garbage by default
- Length not stored
- Does not check bounds when accessing array

❖ Java Arrays:

- elements are initialized to 0 or `null`
- Length stored as an immutable field at start of the array
- Every access to the array does a bounds check, throwing an exception if the index is illegal

Java Data Types: Arrays in Memory

❖ Example Code:

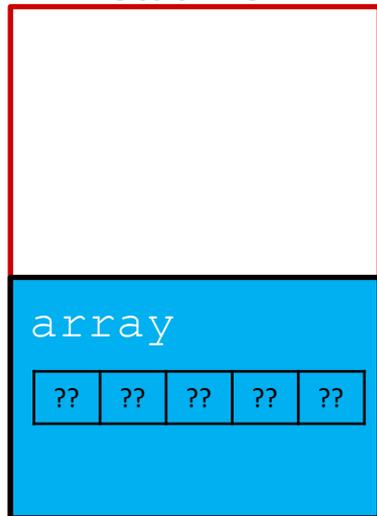
C

```
int main() {  
    int array[5];  
}
```

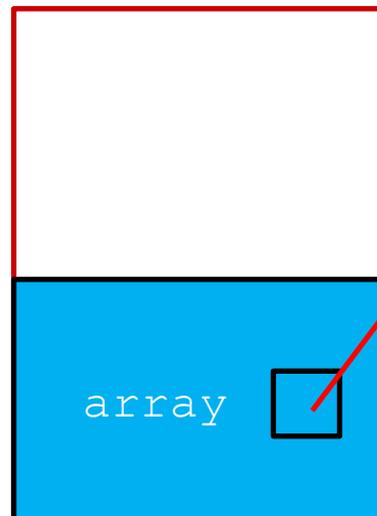
Java

```
public static void main(String args[]) {  
    int[] array = new int[5];  
}
```

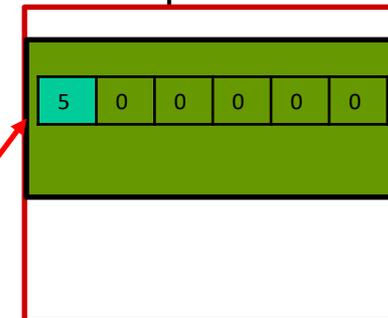
Stack: C

Stack frame for
`main()`

Stack: Java

Stack frame for
`main()`

Heap: Java



Java Data Types: Strings

❖ C strings:

- ASCII Characters
- Pretty much an array of characters
- Null terminated
- Can be modified

❖ Java strings:

- Unicode Characters
- An Object
- Bounded by length like arrays in Java (with a 4-byte int field)
- Are immutable

Java Data Types: Strings in Memory

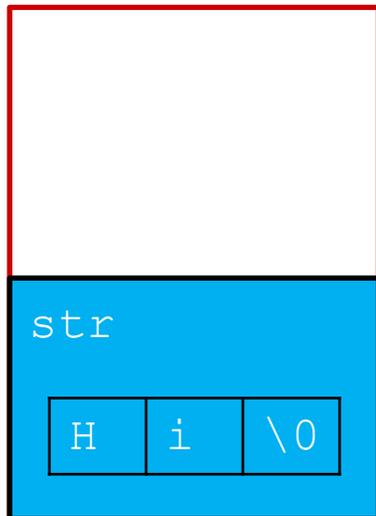
C

```
int main() {
    char str[3] = "Hi";
}
```

Java

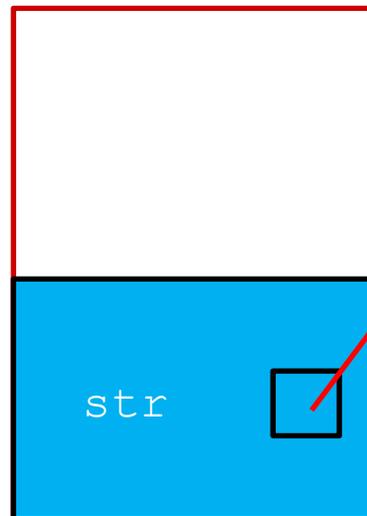
```
public static void main(String args[]) {
    String str = new String("Hi");
}
```

Stack: C



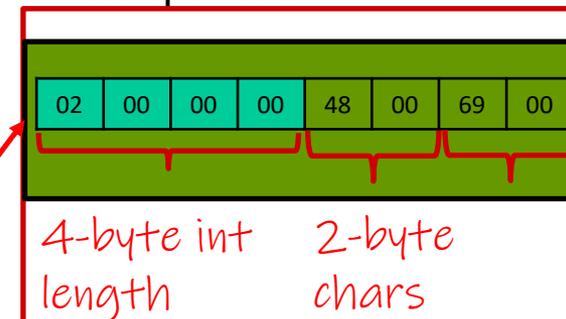
Stack frame for
`main()`

Stack: Java



Stack frame for
`main()`

Heap: Java



Byte level view

Lecture Outline

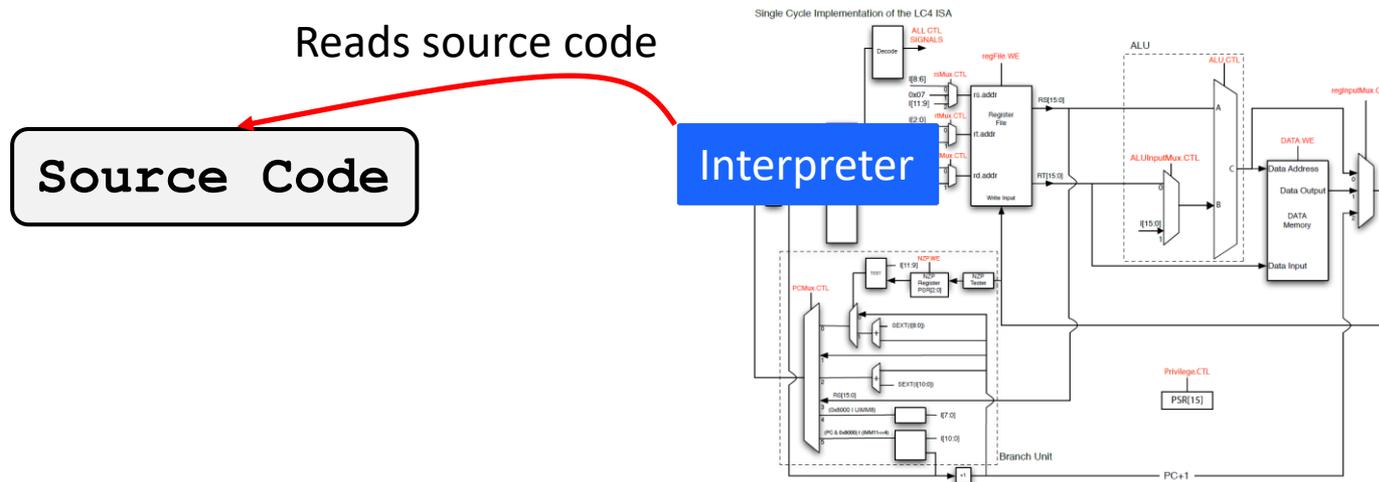
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 - Java Datatypes
 - **Java Compilation**
 - Java Garbage Collector
- ❖ C tips & Practice

None of this is on
the final exam or
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Interpreters

- ❖ There exist other ways for programming languages to run on a computer. A common method is using interpreters
 - Python, Lisp, Javascript, etc.

- ❖ The interpreter is a program that runs directly on the processor, reads your code, and interprets how to emulate the execution of your code.



Intermediate Formats

- ❖ Some languages are not read directly by the interpreter and instead are translated to some intermediary format
 - When we compile Java code, we are compiling from Java to Java bytecode
- ❖ Byte code provides an easier format for the interpreter to read our code
- ❖ Java bytecode can be used to implement other programming languages,
 - Kotlin, Scala, etc

Usual Java Compilation

- ❖ Java code is first compile to Java bytecode by a java compiler
- ❖ Java Byte code is then run on the Java Virtual Machine (JVM) which acts sort of like a Java bytecode interpreter
- ❖ There are other ways to compile and run Java and there are many optimizations that can be made to

JIT

- ❖ Just In Time (JIT) compilation
 - The interpreter/run-time environment will compile some bytecode into machine code while the program is running to try and execute the code faster.
- ❖ Translating to machine code has some overhead cost, especially if the code translation is complex or there are a lot of checks for optimization
- ❖ Some interpreters/environments will try to analyze the code to see which parts of bytecode is worth translating to machine code

Interpreters VS Compilers

- ❖ Interpreters make it easier to run on different architectures since the environment of the program is controlled by the interpreter
- ❖ Interpreters usually have deep connection to a debugger, making development of a debugger easier
- ❖ Allow for a garbage collector to implicitly work while the program is running
- ❖ Interpreters have more overhead cost than compiled languages and run slower
- ❖ Some languages aren't clearly interpreted or compiled

Lecture Outline

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 - **Java Garbage Collector**
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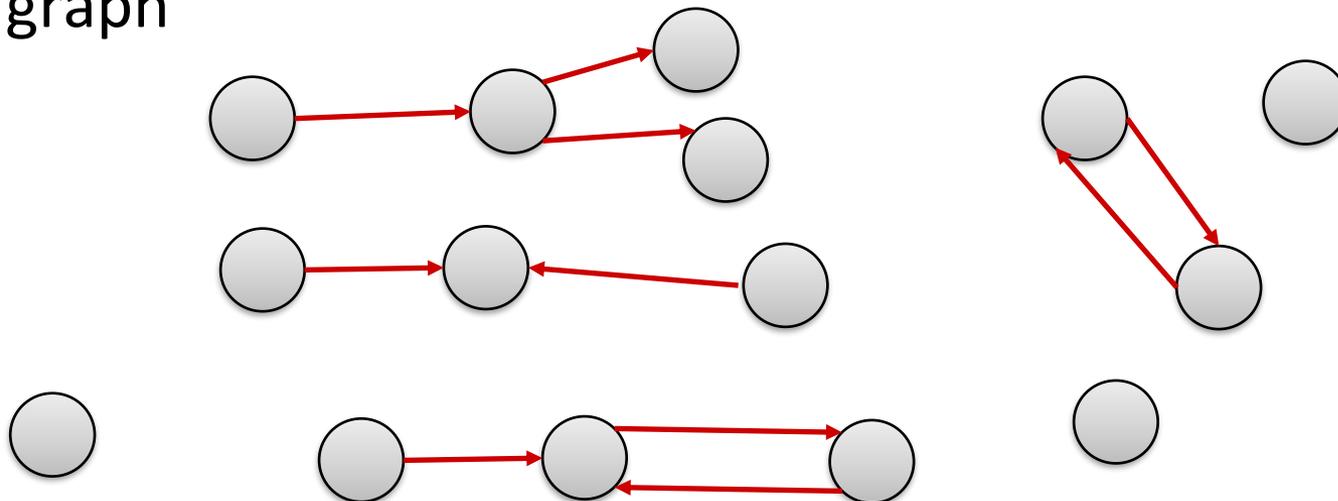
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Garbage Collection

- ❖ Garbage Collection:
 - automatically deallocates memory on the heap.
- ❖ Commonly used in many programming Languages:
 - Java, C#, Go, Javascript, Ruby, Julia, ...
- ❖ Requires some overhead to check and see what memory can be deallocated and which is still being used
- ❖ Many implementations and optimizations on garbage collection

Trace Garbage Collection

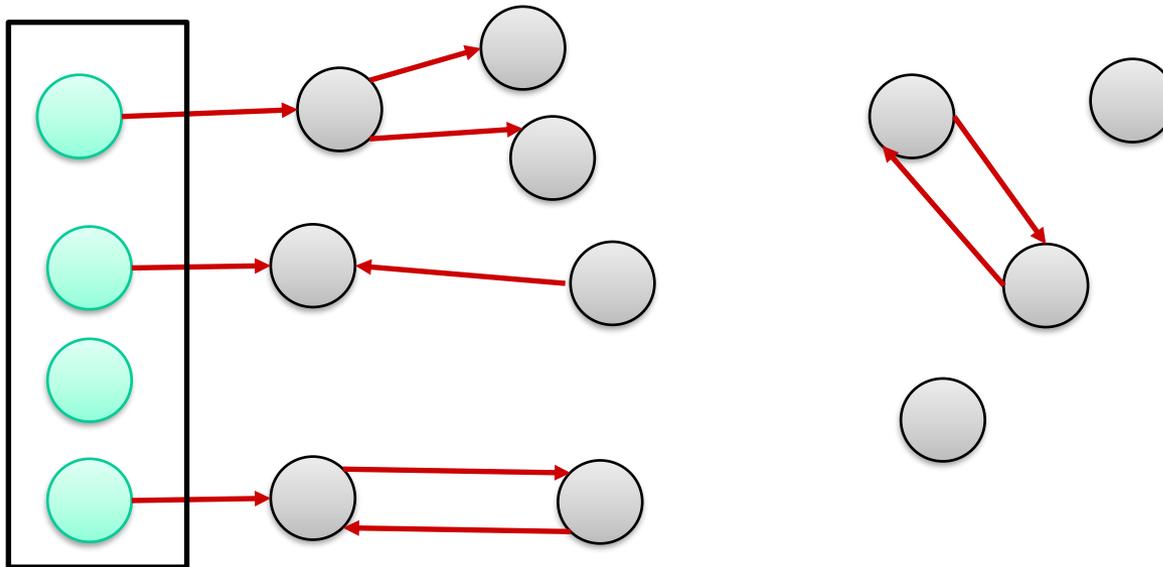
- ❖ To decide which memory can be deallocated, garbage collectors often trace memory to see which memory is still "reachable" by the user program.
- ❖ The garbage collector keeps track of all allocations and can draw memory references & allocations like a directed graph



Trace Garbage Collection

- ❖ We start with a set of allocations we know are reachable and call these Root Nodes (usually these are held as references in local variables still on the stack)

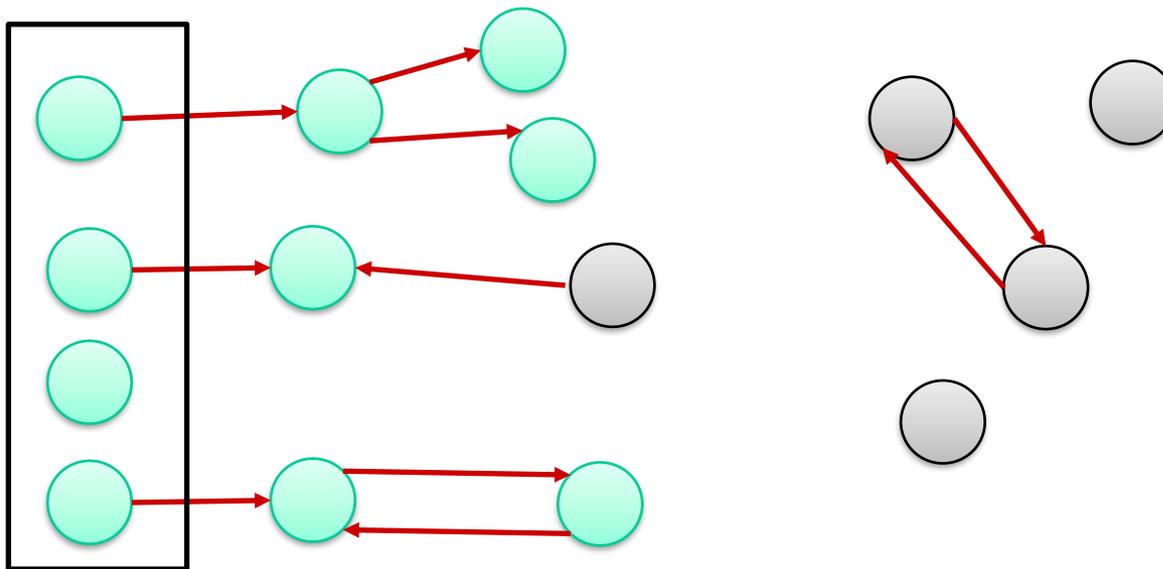
Root Nodes



Trace Garbage Collection

- ❖ We start with a set of allocations we know are reachable and call these Root Nodes (usually these are held as references in local variables still on the stack)
- ❖ We then trace through all references. Anything referenced *from* a reachable node is reachable.

Root Nodes



Lecture Outline

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- ❖ **C tips & Practice**

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C: Common Mistakes

- ❖ The most common mistakes I notice in office hours teaching usually deal with handling memory:
 - How parameters are passed
 - Using Output parameters
 - Exceeding the bounds of an array
 - Issues with deallocating memory

C is Call-By-Value

- ❖ C (and Java) pass arguments by *value*
 - Callee receives a **local copy** of the argument
 - Register or Stack
 - If the callee modifies a parameter, the caller's copy *isn't* modified

```
void swap(int a, int b) {  
    int tmp = a;  
    a = b;  
    b = tmp;  
}  
  
int main(int argc, char** argv) {  
    int a = 42, b = -7;  
    swap(a, b);  
    ...  
}
```

Broken Swap

Note: Arrow points to *next* instruction.

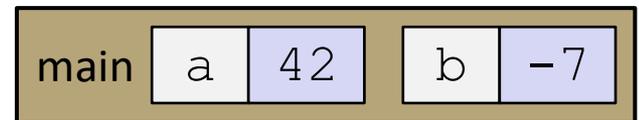
brokenswap.c

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void swap(int a, int b) {  
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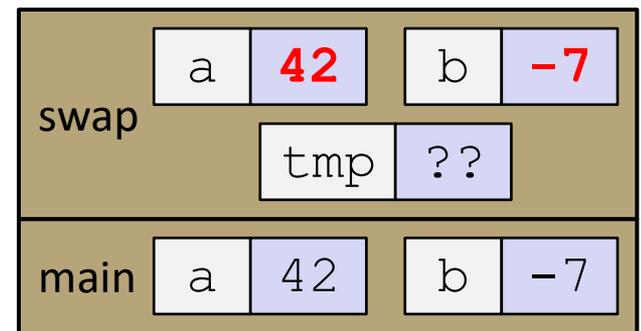
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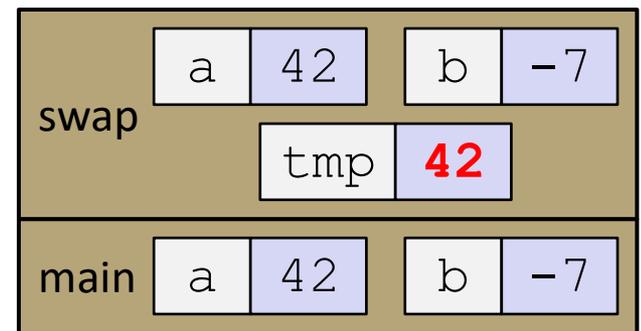
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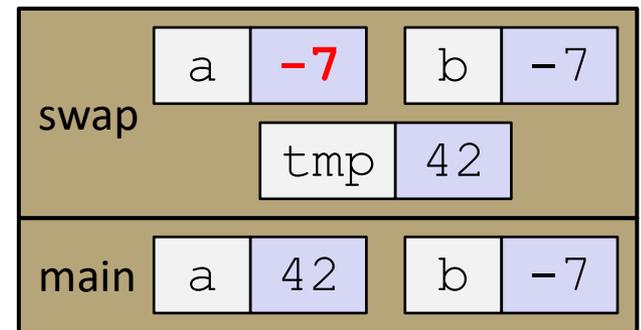
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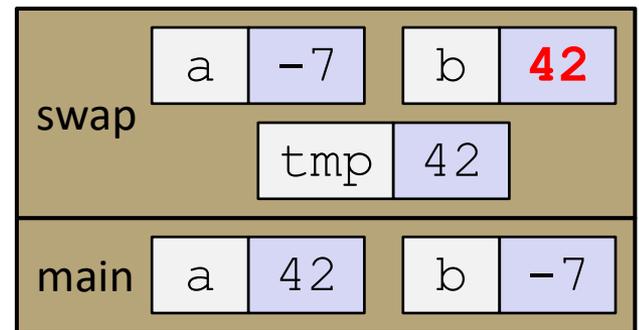
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    ...  
}
```



Faking Call-By-Reference in C

- ❖ Can use pointers to *approximate* call-by-reference
 - Callee still receives a **copy** of the pointer (*i.e.* call-by-value), but it can modify something in the caller's scope by dereferencing the pointer parameter

```
void swap(int* a, int* b) {  
    int tmp = *a;  
    *a = *b;  
    *b = tmp;  
}  
  
int main(int argc, char** argv) {  
    int a = 42, b = -7;  
    swap(&a, &b);  
    ...  
}
```

Fixed Swap

Note: Arrow points to *next* instruction.

swap.c

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void swap(int* a, int* b) {
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}

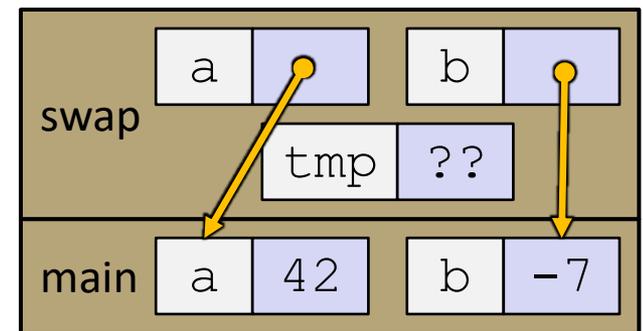
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Fixed Swap

swap.c

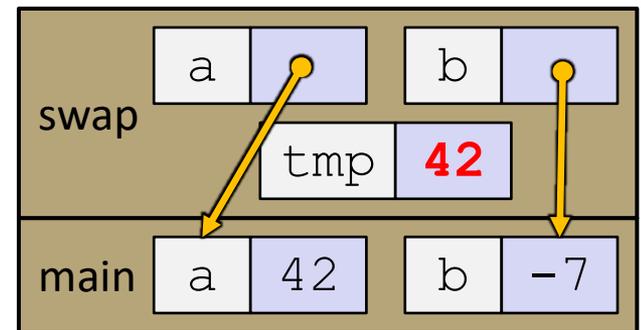
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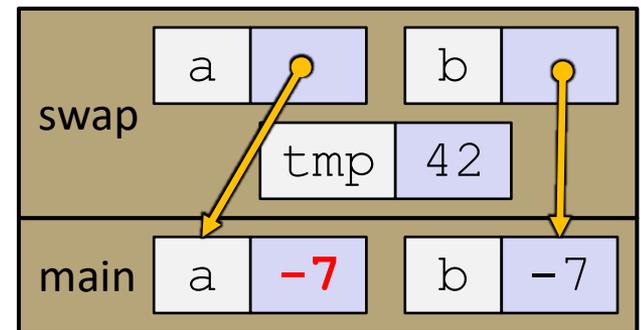
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swap.c

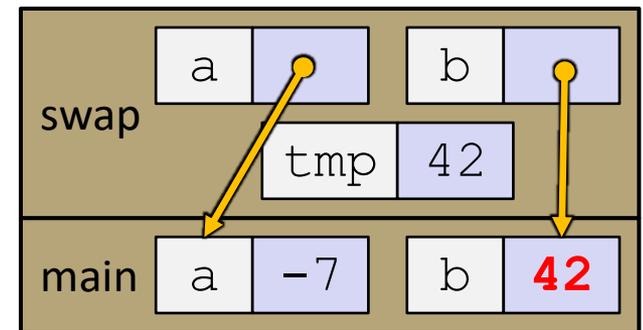
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Fixed Swap

swap.c

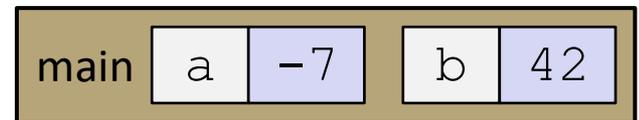
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    int a = 42, b = -7;  
    swap(&a, &b);  
    ...  
}
```



Practice Problem:

- ❖ What does this code print?

```
typedef struct point_st {
    int x, y;
} Point;

void increment_point(Point p) {
    p.x++;
    p.y++;
}

int main() {
    Point p = {1, 5};
    increment_point(p);
    printf("x: %d y: %d\n", p.x, p.y);
}
```

Practice Problem:

- ❖ What does this code print?

```
typedef struct point_st {
    int x, y;
} Point;

void increment_point(Point p) {
    p.x++;
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}

int main() {
    Point p = {1, 5};
    increment_point(p);
    printf("x: %d y: %d\n", p.x, p.y);
}
```

Structs are
passed by
value

This code prints “x: 1, y: 5”

Practice Problem: Fixed

- ❖ Fixed code that uses pointers to simulate pass-by-reference

```
typedef struct point_st {
    int x, y;
} Point;

void increment_point(Point* p) {
    p->x++;
    p->y++;
}

int main() {
    Point p = {1, 5};
    increment_point(&p);
    printf("x: %d y: %d\n", p.x, p.y);
}
```

Practice Problem:

- ❖ What is wrong with this code?

```
#define LINE_LEN 250

int main() {
    FILE* f = fopen("Hi.txt", "r");
    if (f == NULL)
        return EXIT_FAILURE;

    char buf[10];

    while (fread(buf, sizeof(char), LINE_LEN, f)) {
        printf("%s", buf);
    }
    fclose(f);
    return EXIT_SUCCESS;
}
```

Practice Problem:

- ❖ What is wrong with this code?

```
#define LINE_LEN 250

int main() {
    FILE* f = fopen("Hi.txt", "r");
    if (f == NULL)
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    char buf[10];

    while (fread(buf, sizeof(char), LINE_LEN, f)) {
        printf("%s", buf);
    }
    fclose(f);
    return EXIT_SUCCESS;
}
```

buf only has space for 10 characters,
but fread tries to read 250!

This causes stack smashing,
program probably crashes

Practice Problem:

- ❖ What is printed by this code?

```
int main() {
    uint16_t i = 0;
    for (i = 0; i < 65536; i++) {
        printf("%d ", i);
    }
    return EXIT_SUCCESS;
}
```

Practice Problem:

- ❖ What is printed by this code?

```
int main() {
    uint16_t i = 0;
    for (i = 0; i < 65536; i++) {
        printf("%d ", i);
    }
    return EXIT_SUCCESS;
}
```

Code goes infinite!

`i` is of type `uint16_t` which only has
a max value of 65535!

Practice Problem:

- ❖ Similar Issue with unsigned types:

```
int main() {
    uint16_t i;
    for (i = 120; i >= 0; i--) {
        printf("%d ", i);
    }
    return EXIT_SUCCESS;
}
```

i never becomes negative, so the loop condition always evaluates to true

Buggy Program

- ❖ What is wrong with this program?
 - (ignoring style issues)

pair.h

```
#define FOO 240
typedef struct pair_st {
    int x, y;
} pair;
```

util.h

```
#include "pair.h"
#include <stdio.h>

void Pair_Allocate(pair *out) {
    out = (pair *) malloc(sizeof(pair))
    out->x = 0;
    out->y = 0;
}

void Pair_Print(pair *p) {
    printf("(x:%d, y:%d)", p.x, p.y);
}
```

main.c

```
#include "pair.h"
#include "util.h"

int main() {
    pair * p;
    Pair_Allocate(p);
    p->x = FOO;
    p->y = 595;
    Pair_Print(*p);
}
```

Buggy Program

- ❖ What is wrong with this program?
 - (ignoring style issues)

pair.h

```
#define FOO 240
typedef struct pair_st {
    int x, y;
} pair;
```

util.h

```
#include "pair.h"
#include <stdio.h>

void Pair_Allocate(pair *out) {
    out = (pair *) malloc(sizeof(pair))
    out->x = 0;
    out->y = 0;
}

void Pair_Print(pair *p) {
    printf("(x:%d, y:%d)", p.x, p.y);
}
```

Output parameter misuse

Needs to use -> syntax

main.c

```
#include "pair.h"
#include "util.h"

int main() {
    pair * p;
    Pair_Allocate(p);
    p->x = FOO;
    p->y = 595;
    Pair_Print(*p);
}
```

Memory leak

Shouldn't Dereference

Low Level Programming Review

- ❖ Complete the function "rand_string", which generates a random string of random length. Assume we have the following functions available to you:

- ```
int rand_len();
// returns a random int in the range of 1 - 256
```
- ```
char rand_char();  
// returns a random printable character  
// (no '\0' or other special characters)
```

- ❖ If you finish, write a small main function that calls `rand_string` and prints out the string

Low Level Programming Review

```
// returns a random string and its length. Returns -1 on error
int rand_string(char **output) {
    // generate random length
    int len = rand_len();

    // allocate space for the string (+1 for null terminator)
    char* result = (char *) malloc((len+1)*sizeof(char));
    // error checking
    if (result == NULL)
        return -1;

    // assign random characters
    for (int i = 0; i < len; i++)
        result[i] = rand_char();

    // add null terminator
    result[len] = '\0';

    // return results
    *output = result;
    return len;
}
```

Low Level Programming Review

```
int main() {
    char* str;
    // generate random string
    int len = rand_string(&str);

    if (len == -1)
        return EXIT_FAILURE;

    printf("%s\n", str);

    free(str);

    return EXIT_SUCCESS;
}
```

C Tips: Ownership

- ❖ In C and C++, there is no garbage collector and instead the programmer has to manage memory allocation/deallocation themselves.
- ❖ To help reason about code, try to think of who has the "Ownership" or "Responsibility" to free code and who is just "Borrowing" memory.
- ❖ In the previous example, the comment for `rand_string` should say something like "The caller is responsible for freeing the resulting string"