

**CIS 500**

Software Foundations

Fall 2003

**22 October**

# References

# Mutability

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- ◆ In most programming languages, **variables** are mutable — i.e., a variable provides both
  - ◆ a name that refers to a previously calculated value, and
  - ◆ the possibility of **overwriting** this value with another (which will be referred to by the same name)
- ◆ In some languages (e.g., OCaml), these two features are kept separate
  - ◆ variables are only for naming — the binding between a variable and its value is immutable
  - ◆ introduce a new class of **mutable values** (called **reference cells** or **references**)
  - ◆ at any given moment, a reference holds a value (and can be **dereferenced** to obtain this value)
  - ◆ a new value may be **assigned** to a reference

We choose OCaml's style, which is easier to work with formally.

So a variable of type `T` in most languages (except OCaml) will correspond to a `Ref T` (actually, a `Ref(Option T)`) here.

# Basic Examples

---

```
r = ref 5
```

```
!r
```

```
r := 7
```

```
(r:=succ(!r); !r)
```

```
(r:=succ(!r); r:=succ(!r); r:=succ(!r); r:=succ(!r); !r)
```

# Basic Examples

---

```
r = ref 5
```

```
!r
```

```
r := 7
```

```
(r:=succ(!r); !r)
```

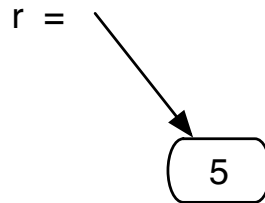
```
(r:=succ(!r); r:=succ(!r); r:=succ(!r); r:=succ(!r); !r)
```

**i.e.,**

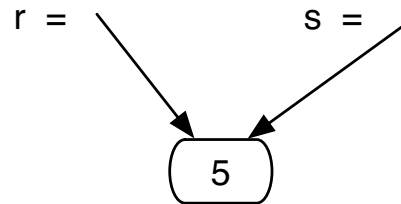
```
((((r:=succ(!r); r:=succ(!r)); r:=succ(!r)); r:=succ(!r)); !r)
```

# Aliasing

A value of type `Ref T` is a **pointer** to a cell holding a value of type `T`.



If this value is “copied” by assigning it to another variable, the cell pointed to is not copied.



So we can change `r` by assigning to `s`:

```
(s:=6; !r)
```

## Aliasing all around us

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Reference cells are not the only language feature that introduces the possibility of aliasing.

- ◆ arrays
- ◆ communication channels
- ◆ I/O devices (disks, etc.)



## The difficulties of aliasing

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The possibility of aliasing invalidates all sorts of useful forms of reasoning about programs, both by programmers...

The function

```
 $\lambda r:\text{Ref Nat. } \lambda s:\text{Ref Nat. } (r:=2; s:=3; !r)$ 
```

always returns 2 **unless**  $r$  and  $s$  are aliases for the same cell.

...and by compilers:

Code motion out of loops, common subexpression elimination, allocation of variables to registers, and detection of uninitialized variables all depend upon the compiler knowing which objects a load or a store operation could reference.

High-performance compilers spend significant energy on **alias analysis** to try to establish when different variables cannot possibly refer to the same storage.

## The benefits of aliasing

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The problems of aliasing have led some language designers simply to disallow it (e.g., Haskell).

But there are good reasons why most languages do provide constructs involving aliasing:

- ◆ efficiency (e.g., arrays)
- ◆ “action at a distance” (e.g., symbol tables)
- ◆ shared resources (e.g., locks) in concurrent systems
- ◆ etc.

## Example

---

```
c = ref 0
```

```
incc = λx:Unit. (c := succ (!c); !c)
```

```
decc = λx:Unit. (c := pred (!c); !c)
```

```
incc unit
```

```
decc unit
```

```
o = {i = incc, d = decc}
```

```
let newcounter =  
  λ_:Unit.  
    let c = ref 0 in  
    let incc = λx:Unit. (c := succ (!c); !c) in  
    let decc = λx:Unit. (c := pred (!c); !c) in  
    let o = {i = incc, d = decc} in  
    o
```

# Syntax

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$t ::=$

`unit`

`x`

`$\lambda x:T.t$`

`t t`

`ref t`

`!t`

`t:=t`

terms

unit constant

variable

abstraction

application

reference creation

dereference

assignment

... plus other familiar types, in examples.

# Typing Rules

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$$\frac{\Gamma \vdash t_1 : T_1}{\Gamma \vdash \text{ref } t_1 : \text{Ref } T_1} \quad (\text{T-REF})$$

$$\frac{\Gamma \vdash t_1 : \text{Ref } T_1}{\Gamma \vdash !t_1 : T_1} \quad (\text{T-DEREF})$$

$$\frac{\Gamma \vdash t_1 : \text{Ref } T_1 \quad \Gamma \vdash t_2 : T_1}{\Gamma \vdash t_1 := t_2 : \text{Unit}} \quad (\text{T-ASSIGN})$$

## Final example

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```
NatArray = Ref (Nat → Nat);
```

```
newarray = λ_:Unit. ref (λn:Nat.0);  
          : Unit → NatArray
```

```
lookup = λa:NatArray. λn:Nat. (!a) n;  
        : NatArray → Nat → Nat
```

```
update = λa:NatArray. λm:Nat. λv:Nat.  
         let oldf = !a in  
         a := (λn:Nat. if equal m n then v else oldf n);  
         : NatArray → Nat → Nat → Unit
```