CIS 500 Software Foundations Fall 2006

November 6

Some Hints

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- The exam will potentially cover everything in the course so far, but will focus on material we've seen since the first midterm.
- ► There *will* be a question that is also a one-star exercise from the book.
- There will be a question similar to problem 6 from midterm 1 ("Which properties remain true if we change one fo the type systems we've studied in the following way...?")
- ▶ There will be (at least) one question based on one of the proofs in chapter 15.
- ► For PhD students, there *will* be a question involving subtyping and references.

Review

```
\blacktriangleright \lambda x:Bool \rightarrow Bool. x (x (true)))
```

- $\blacktriangleright \lambda x:Bool \rightarrow Bool. x (x (true)))$
- ► (λ x:Bool. λ y:Bool \rightarrow Bool. true) false (λ z:Bool \rightarrow Bool. true)

- $\blacktriangleright \lambda x:Bool \rightarrow Bool. x (x (true)))$
- ▶ $(\lambda x:Bool. \lambda y:Bool \rightarrow Bool. true)$ false $(\lambda z:Bool \rightarrow Bool. true)$
- (λx :Bool. λy :Bool. error) false false false false

- $\blacktriangleright \lambda x:Bool \rightarrow Bool. x (x (true)))$
- $(\lambda x:Bool. \lambda y:Bool \rightarrow Bool. true)$
- false (λ z:BooloBool. true)
- (λ x:Bool. λ y:Bool. error) false false false false false
- (λx :Bool. λy :Bool. true) false false false false false

```
What are the types of these expressions?
```

- $ightharpoonup \lambda x: Bool
 ightharpoonup Bool. x (x (true)))$
- $(\lambda x: Bool. \lambda y: Bool \rightarrow Bool. true)$
 - (λ x:Bool. λ y:Bool. error) false false false false false
 - ► (\(\lambda\x:\)Bool. \(\lambda\x:\)Bool. \(\text{true}\)
 false false false false false
 - ▶ try
 (if (\(\lambda x\):Bool. x) error
 then (error false)
 else error)
 - with $\lambda y : Bool \rightarrow Bool. y$

For reference: Typing rules for exceptions

Give the result of evaluation and the final store after each of these expressions is evaluated to a normal form starting in the empty store.

```
let x = ref 0 in
let y = ref 1 in
y:=3
!x
```

Give the result of evaluation and the final store after each of these expressions is evaluated to a normal form starting in the empty store.

```
let y = ref 1 in
  y:=3
  ! x
\triangleright let x = ref 0 in
```

 \triangleright let x = ref 0 in

- let y = ref 1 in
- let x = y in !x

Give the result of evaluation and the final store after each of these expressions is evaluated to a normal form starting in the empty store.

```
\triangleright let x = ref 0 in
  let y = ref 1 in
  y:=3
  ! X
\triangleright let x = ref 0 in
  let y = ref 1 in
  let x = y in !x
\triangleright let x = ref 0 in
  let y = x in
  let x = y in !x
```

For reference: Evaluation rules for references

$$\frac{\textit{I} \notin \textit{dom}(\mu)}{\text{ref } v_1 \mid \mu \longrightarrow \textit{I} \mid (\mu, \textit{I} \mapsto v_1)} \qquad \text{(E-RefV)}$$

$$\frac{\mu(\textit{I}) = v}{!\textit{I} \mid \mu \longrightarrow v \mid \mu} \qquad \text{(E-DerefLoc)}$$

$$\textit{I} := v_2 \mid \mu \longrightarrow \text{unit} \mid [\textit{I} \mapsto v_2]\mu \qquad \text{(E-Assign)}$$

(Plus several congruence rules.)

 $\triangleright \lambda x: Ref Nat. !x+1$

```
\triangleright \lambda x: Ref Nat. !x+1
```

```
▶ λx:Ref Nat. x
```

```
\triangleright \lambda x:Ref Nat. !x+1
```

```
\triangleright \lambda x:Ref Nat. x
```

```
▶ \lambda x:Ref Nat. (x:=3; 1_1:=42; !1_1)
```

```
\triangleright \lambda x:Ref Nat. !x+1
```

 $\triangleright \lambda x$:Ref Nat. x

```
▶ \lambda x:Ref Nat. (x:=3; 1_1:=42; !1_1)
```

 \blacktriangleright λ f:Unit \rightarrow Unit. (l₁:=3; f unit; !l₁)

Preservation and progress for chapter 13

- ▶ The preservation and progress proofs for λ with references are just sketched in TAPL.
- Working out the details for yourself is an excellent exercise
- ▶ A question based on this proof may appear on the final exam, but will not appear on the coming midterm

For each of the following pairs of terms, say whether the one on the left is a subtype of the one on the right, a supertype, equivalent, or incomparable.

 \blacktriangleright ({} \rightarrow {}) \rightarrow Top and Top \rightarrow Top

- lacktriangle ({} \rightarrow {}) \rightarrow Top and Top \rightarrow Top
- $\qquad \qquad (\mathsf{Top} \rightarrow \mathsf{Top}) \rightarrow \{\} \rightarrow \{\} \text{ and } (\mathsf{Top} \rightarrow \{\}) \rightarrow \mathsf{Top}$

```
 (\{\} \rightarrow \{\}) \rightarrow \mathsf{Top} \text{ and } \mathsf{Top} \rightarrow \mathsf{Top}
```

- $\qquad \qquad (\mathsf{Top} \rightarrow \mathsf{Top}) \rightarrow \{\} \rightarrow \{\} \text{ and } (\mathsf{Top} \rightarrow \{\}) \rightarrow \mathsf{Top}$
- ► {a:Top, b:{c:Top,d:Top}} and {b:{d:Top,c:Top},a:Top}

- lacktriangle ({} \rightarrow {}) \rightarrow Top and Top \rightarrow Top
- $(Top \rightarrow Top) \rightarrow \{\} \rightarrow \{\} \text{ and } (Top \rightarrow \{\}) \rightarrow Top$
- ► {a:Top, b:{c:Top,d:Top}} and {b:{d:Top,c:Top},a:Top}
- ► <1:Top,m:{n:Top}> and <m:{n:Top,o:Top}>

- lacktriangle ({} \rightarrow {}) \rightarrow Top and Top \rightarrow Top
- $(Top \rightarrow Top) \rightarrow \{\} \rightarrow \{\} \text{ and } (Top \rightarrow \{\}) \rightarrow Top$
- ► {a:Top, b:{c:Top,d:Top}} and {b:{d:Top,c:Top},a:Top}
- ► <1:Top,m:{n:Top}> and <m:{n:Top,o:Top}>
- ► <>→Top and {}→Top

Draw a subtyping derivation for the following statement:

```
(\texttt{Top} \rightarrow \{\texttt{x}: \texttt{Nat}\}) \rightarrow \{\texttt{x}: \texttt{Nat}, \texttt{y}: \texttt{Nat}\} \quad <: \quad (\{\} \rightarrow \{\}) \rightarrow \{\texttt{y}: \texttt{Nat}\}
```

For reference: Subtyping rules

$$S <: S \qquad (S-Refl)$$

$$\frac{S <: U \quad U <: T}{S <: T} \qquad (S-TRANS)$$

$$\{1_i : T_i \stackrel{i \in 1...n+k}{} \} <: \{1_i : T_i \stackrel{i \in 1...n}{} \} \qquad (S-RCDWIDTH)$$

$$\frac{for \ each \ i \quad S_i <: T_i}{\{1_i : S_i \stackrel{i \in 1...n}{} \} <: \{1_i : T_i \stackrel{i \in 1...n}{} \}} \qquad (S-RCDDEPTH)$$

$$\frac{\{k_j : S_j \stackrel{j \in 1...n}{} \} \ is \ a \ permutation \ of \ \{1_i : T_i \stackrel{i \in 1...n}{} \}}{\{k_j : S_j \stackrel{j \in 1...n}{} \} <: \ \{1_i : T_i \stackrel{i \in 1...n}{} \}} \qquad (S-RCDPERM)}$$

$$\frac{T_1 <: \ S_1 \qquad S_2 <: \ T_2}{S_1 \rightarrow S_2 <: \ T_1 \rightarrow T_2} \qquad (S-ARROW)}{S <: \ Top \qquad (S-Top)}$$

Ascription as a derived form

- ▶ Someone asked to work exercise 11.4.1 part 2 today.
- ▶ But the solution is somewhat technical and would take too much time to discuss in detail.
- This exercise is not needed for the exam.

The Hints, again

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