

Lecture 1

CIS 500: SOFTWARE FOUNDATIONS

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Fall 2016

How do we build software?

^
that works
^
(and be confident
that it does)

Critical Software

```
graph TD; CS[Critical Software] --> IP[Individual programs]; CS --> PL[Programming languages];
```

Individual programs

- Operating systems
- Network stacks
- Crypto
- Medical devices
- Flight control systems
- Power plants
- Home security
- ...

Programming languages

- Static type systems
- Data abstraction and modularity
- Security controls
- Compiler correctness

Logic

+ Reasoning about
individual programs

+ Reasoning about
whole programming
languages

SOFTWARE FOUNDATIONS



LOGICAL FOUNDATIONS



Q: How do we know something is true?

A: We test it out

Q: But that isn't truth; testing can only give us evidence. How do we know something is **true**?

A: We prove it

Q: How do we know that we have a proof?

A: We need to define what it means to be a proof.

A proof is a logical sequence of arguments, starting from some initial assumptions

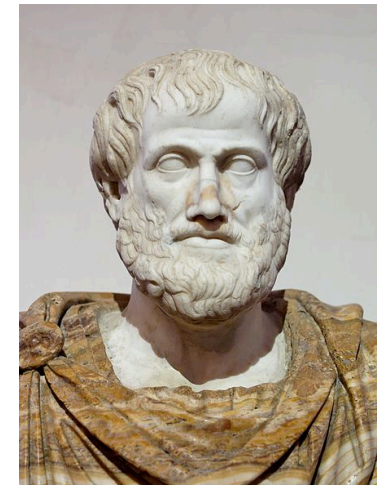
Q: How do we know that we have a valid sequence of arguments? Can any list be a proof?

All humans are mortal

All Greeks are human

Therefore I am a Greek!

A: No, no, no! We need to think about how we *think*....



Aristotle
384 – 322 BC



Euclid
~300 BC

First we need a language...

- **Gottlob Frege**: a German mathematician who started in geometry but became interested in logic and foundations of arithmetic.
- 1879 Published "*Begriffsschrift, eine der arithmetischen nachgebildete Formelsprache des reinen Denkens*" (Concept-Script: A Formal Language for Pure Thought Modeled on that of Arithmetic)
 - First rigorous treatment of functions and quantified variables
 - $\vdash A, \neg A, \forall x.F(x)$
 - First notation able to express arbitrarily complicated logical statements



Gottlob Frege
1848-1925



Formalization of Arithmetic

- 1884: *Die Grundlagen der Arithmetik* (The Foundations of Arithmetic)
- 1893: *Grundgesetze der Arithmetik* (Basic Laws of Arithmetic, Vol. 1)
- 1903: *Grundgesetze der Arithmetik* (Basic Laws of Arithmetic, Vol. 2)
- Frege's Goals:
 - isolate logical principles of inference
 - derive laws of arithmetic from first principles
 - set mathematics on a solid foundation of logic
- **David Hilbert**: a German recognized as one of the most influential mathematicians ever.
 - algebra, axiomatization of geometry, physics,...
 - 1900: published his "23 Problems"
 - Problem #2: Prove that the axioms of arithmetic are consistent



Hilbert
1943

The plot thickens...

Just as Volume 2 was going to print in 1903,
Frege received a letter...

Bertrand Russell

- *Russell's paradox:*

1. Set comprehension notation:

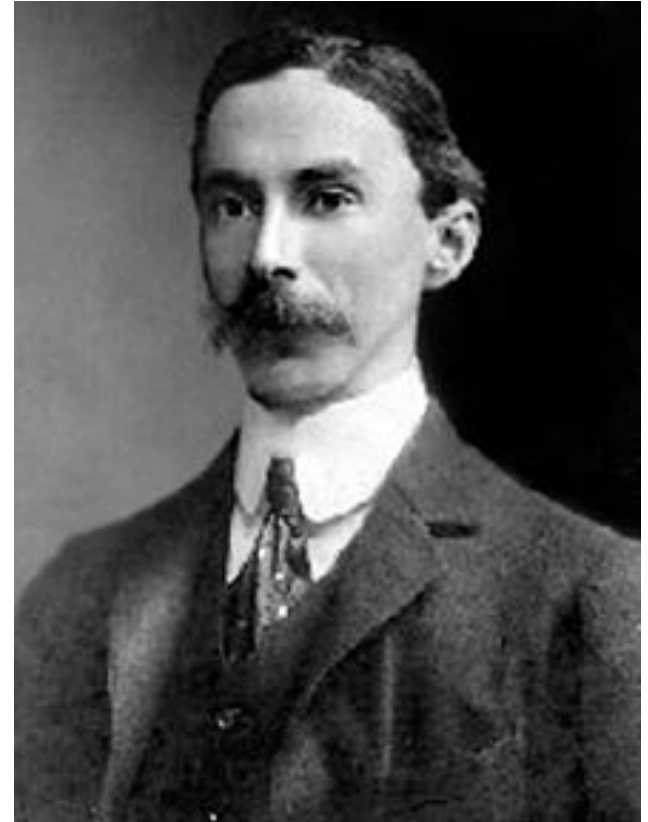
$\{ x \mid P(x) \}$ "The set of x such that $P(x)$ "

2. Let X be the set $\{ Y \mid Y \notin X \}$.

3. Ask the logical question:

Does $X \in X$ hold?

4. **Paradox!** If $X \in X$ then $X \notin X$.
If $X \notin X$ then $X \in X$.



Bertrand Russell
1872 - 1970

- Frege's language could derive Russell's paradox \Rightarrow it was *inconsistent*.
- Frege's logical system could derive anything. Oops(!!)

Addendum to Frege's 1903 Book

*“Hardly anything more unfortunate can befall a scientific writer than to have one of the **foundations** of his edifice shaken after the work is finished. This was the position I was placed in by a letter of Mr. Bertrand Russell, just when the printing of this volume was nearing its completion.”*

– Frege, 1903

Aftermath of Frege and Russell

- Frege came up with a fix, but it made his logic trivial...
- **1908**: Russell fixed the inconsistency of Frege's logic by developing a *theory of types*.
- **1910, 1912, 1913**, (revised **1927**):
Principia Mathematica (Whitehead & Russell)
 - Goal: axioms and rules from which *all* mathematical truths could be derived.
 - It was a bit unwieldy...

"From this proposition it will follow, when arithmetical addition has been defined, that $1+1=2$."

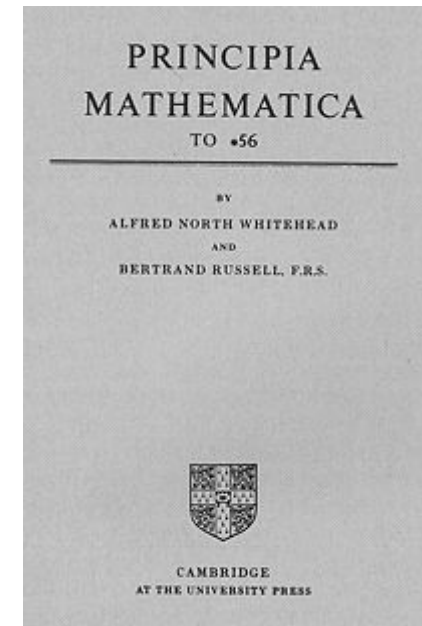
—Volume I, 1st edition, *page 379*



Whitehead



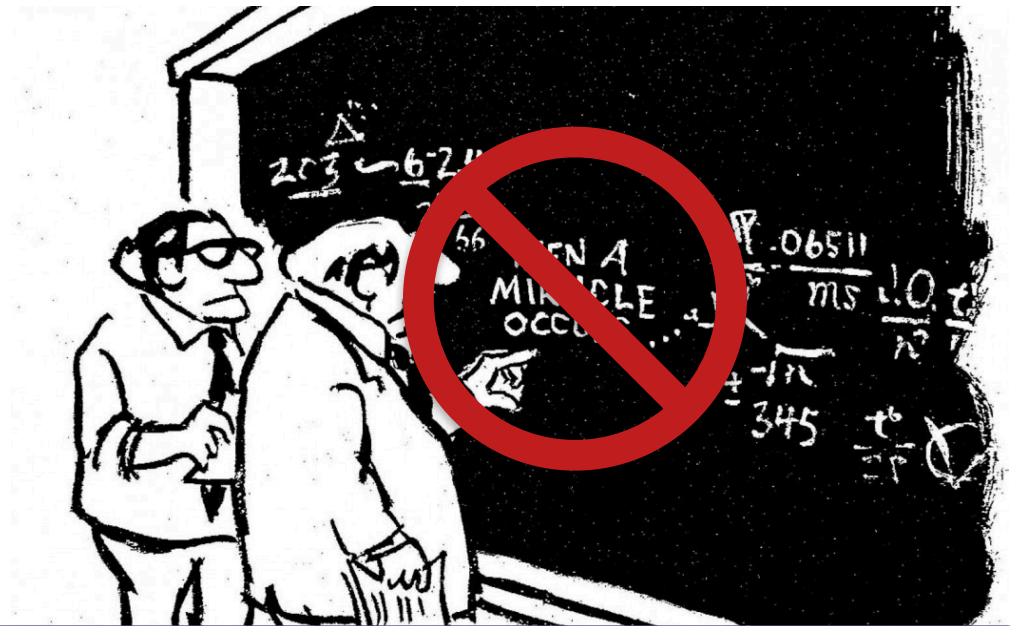
Russell



1920's: Hilbert's Program

A plan to secure the foundations of mathematics:

- Develop a *formal system* of all mathematics.
 - Mathematical statements should be written in a precise formal language
 - Mathematical proofs should proceed by well-specified rules
- Prove *completeness*
 - i.e. that all true mathematical statements can be proved
- Prove *consistency*
 - i.e. that no contradictory conclusions can be proved
- Prove *decidability*
 - i.e. there should be an algorithm for determining whether a given statement has a proof



Things were going well, following Russell & Whitehead, until...

Logic in the 1930s and 1940s

- **1931:** Kurt Gödel's first and second incompleteness theorems.
 - Demonstrated that any consistent formal theory capable of expressing arithmetic cannot be complete.
 - Write down: "This statement is not provable." as an arithmetic statement.
- **1936:** Genzen proves **consistency** of arithmetic.
- **1936:** Church introduces the **λ -calculus**.
- **1936:** Turing introduces Turing machines
 - Is there a decision procedure for arithmetic?
 - Answer: no, it's undecidable
 - The famous "halting problem"
 - only in 1938 did Turing get his Ph.D.
- **1940:** Church introduces the ***simple theory of types***



Kurt Gödel
1906 - 1978



Gerhard Gentzen
1909 - 1945



Alonzo Church
1903 - 1995



Alan Turing
1912 - 1954

Fast Forward...

- 1958 (Haskell Curry) and 1969 (William Howard) observe a remarkable correspondence:



Haskell Curry
1900 – 1982



William Howard
1926 –

types	~	propositions
programs	~	proofs
computation	~	simplification



N.G. de Bruijn
1918 - 2012

- 1967 – 1980's: N.G. de Bruijn runs Automath project
 - uses the Curry-Howard correspondence for computer-verified mathematics

- 1971: Jean-Yves Girard introduces System F
- 1972: Girard introduces F_ω
- 1972: Per Martin-Löf introduces intuitionistic type theory
- 1974: John Reynolds independently discovers System F

Basis for modern
type systems:
OCaml, Haskell,
Scala, Java, C#, ...

... to the Present

- 1984: Coquand and Huet first begin implementing a new theorem prover “Coq”
- 1985: Coquand introduces the **calculus of constructions**
 - combines features from intuitionistic type theory and $F\omega$
- 1989: Coquand and Paulin extend CoC to the **calculus of inductive constructions**
 - adds “inductive types” as a primitive
- 1992: Coq ported to Xavier Leroy’s OCaml
- 1990’s: up to Coq version 6.2
- 2000-2015: up to Coq version 8.4
- 2016: Coq version 8.5 ← CIS 500

- 2013: Coq receives ACM Software System Award



Thierry Coquand
1961 –



Gérard Huet
1947 –

Too many contributors to list here...

So much for foundations... what about the “software” part?

PROGRAMMING FOUNDATIONS

Building Reliable Software

- Suppose you work at (or run) a software company.
- Suppose, like Frege, you've sunk 30+ person-years into developing the "next big thing":
 - Boeing Dreamliner2 flight controller
 - Autonomous vehicle control software for Nissan
 - Gene therapy DNA tailoring algorithms
 - Super-efficient green-energy power grid controller
- Suppose, like Frege, your company has invested a lot of material resources that are also at stake.
- How do you avoid getting a letter like the one from Russell?

Or, worse yet, *not* getting the letter, with disastrous consequences down the road?

Approaches to Software Reliability

- Social
 - Code reviews
 - Extreme/Pair programming
- Methodological
 - Design patterns
 - Test-driven development
 - Version control
 - Bug tracking
- Technological
 - “lint” tools, static analysis
 - Fuzzers, random testing
- Mathematical
 - Sound type systems
 - “Formal” verification



Less “formal”: Techniques may miss problems in programs

This isn't a tradeoff... all of these methods should be used.

Even the most “formal” argument can still have holes:

- Did you prove the right thing?
- Do your assumptions match reality?
- Knuth: “Beware of bugs in the above code; I have only proved it correct, not tried it.”

More “formal”: eliminate *with certainty* as many problems as possible.

Five Interwoven Threads

1. basic tools from **logic** for making and justifying precise claims about programs
2. the use of **proof assistants** to construct rigorous, machine checkable, logical arguments
3. the idea of **functional programming**, both as a method of programming and as a bridge between programming and logic
4. techniques for formal **verification** of properties of specific programs
5. the use of **type systems** for establishing well-behavedness guarantees for all programs in a given language

Can it Scale?

Use of formal methods to verify full-scale software systems is a hot research topic!

- **CompCert** – fully verified C compiler
Leroy, INRIA
- **Vellvm** – formalized LLVM IR
Zdancewic, Penn
- **Ynot** – verified DBMS, web services
Morrisett, Harvard
- **Verified Software Toolchain**
Appel, Princeton
- **Bedrock** – web programming, packet filters
Chlipala, MIT
- **CertiKOS** – certified OS kernel
Shao & Ford, Yale

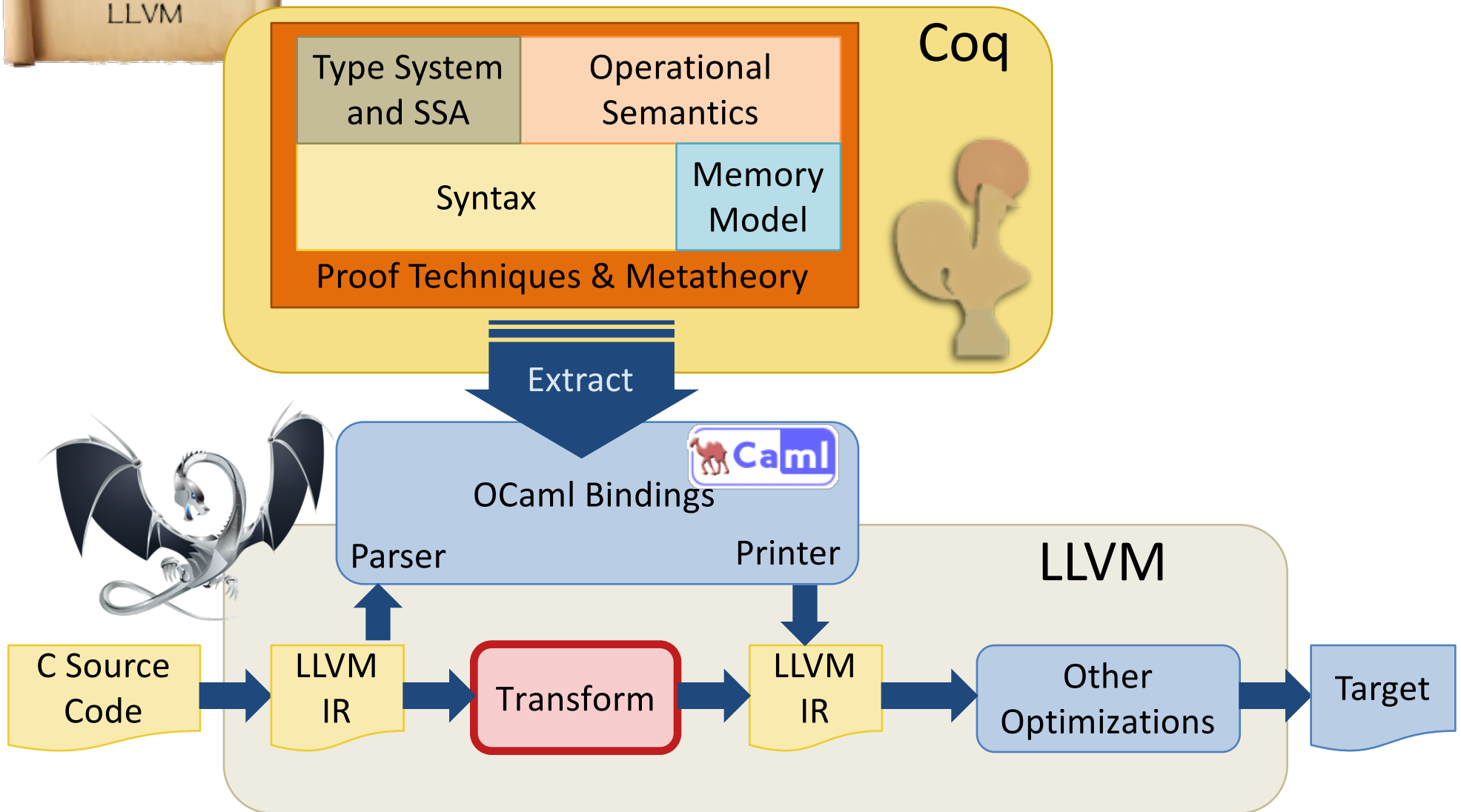


Bedrock



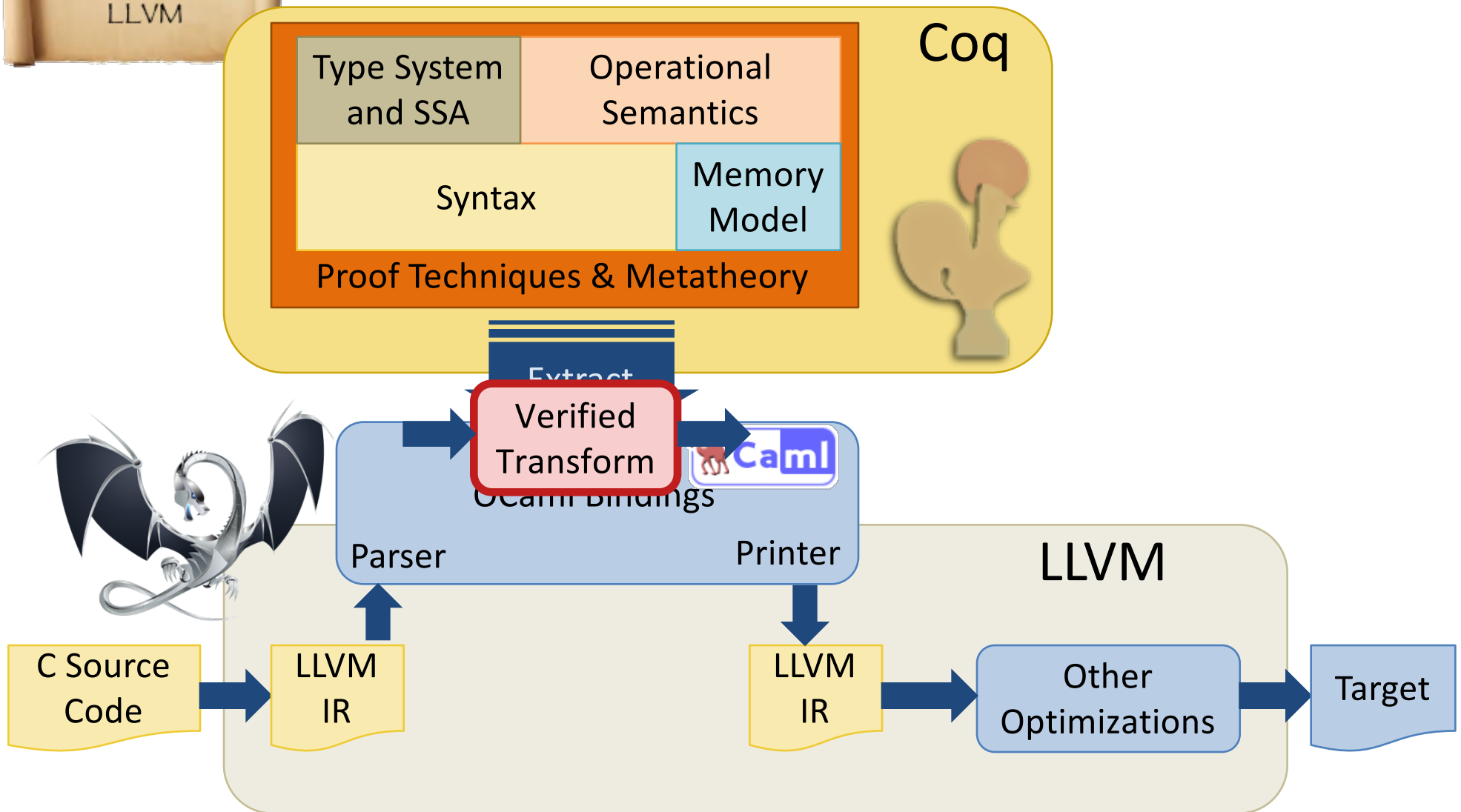
Vellvm Framework

Vellvm
verified
LLVM



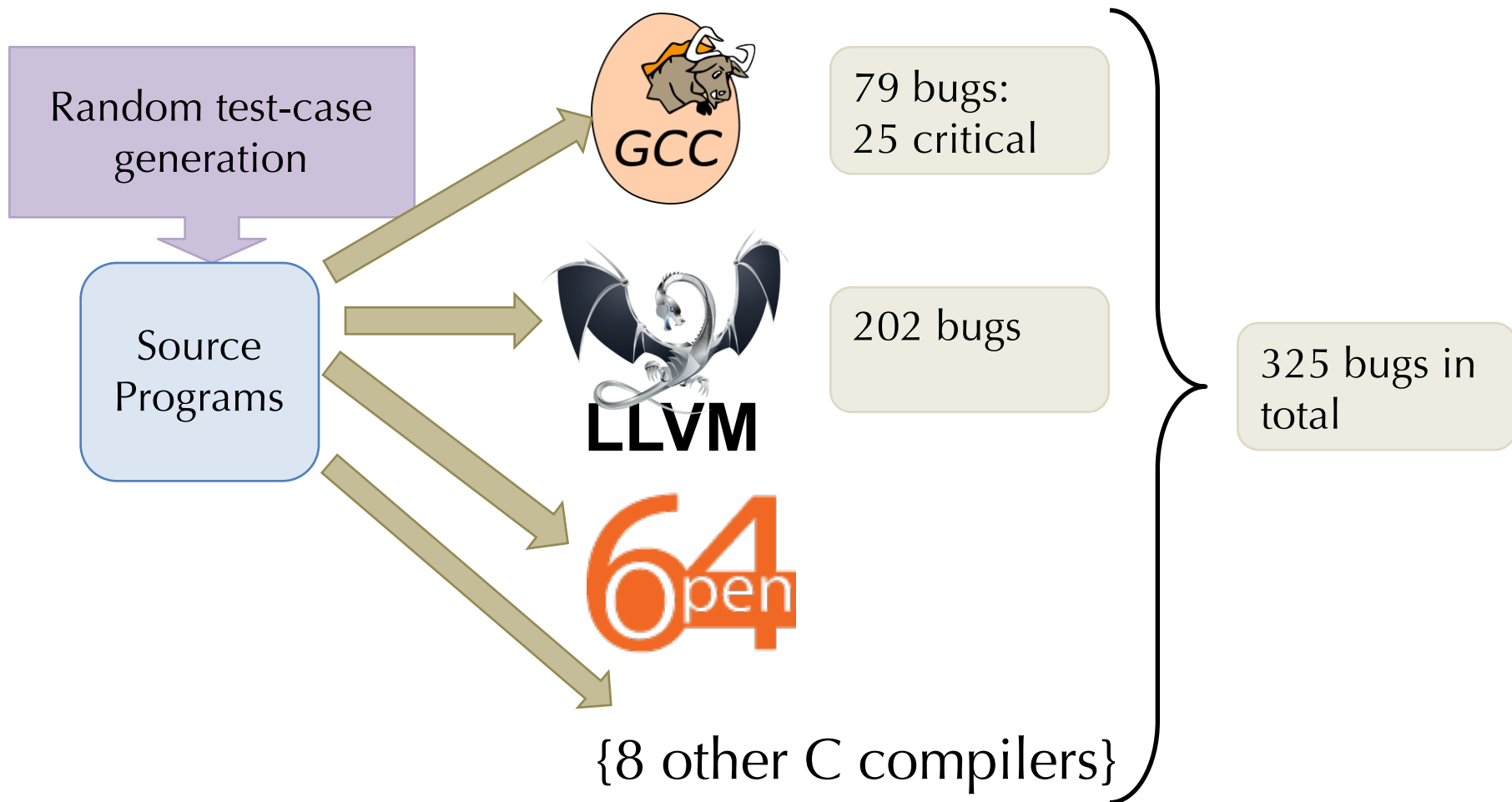
Vellvm Framework

Vellvm
verified
LLVM



Does it work?

Finding and Understanding Bugs in C Compilers [Yang et al. PLDI 2011]



Verified Compiler: CompCert [Leroy et al.]
<10 bugs found in *unverified* front-end component

Regehr's Group Concludes

The striking thing about our CompCert results is that the *middle-end bugs* we found in all other compilers are *absent*. As of early 2011, the under-development version of *CompCert is the only compiler we have tested for which Csmith cannot find wrong-code errors*. This is not for lack of trying: we have devoted about six CPU-years to the task. *The apparent unbreakability of CompCert supports a strong argument that developing compiler optimizations within a proof framework, where safety checks are explicit and machine-checked, has tangible benefits for compiler users.*




the science of deep specification

- National Science Foundation "Expedition" Project
 - \$10M over five years
 - Penn: Pierce / Weirich / Zdancewic
 - Princeton: Appel
 - Yale: Shao
 - MIT: Chlipala
- Many ways to get involved (especially after CIS 500!)
- See www.deepspec.org



CIS 500

- **Foundations**
 - Functional programming
 - Constructive logic
 - Logical foundations
 - Proof techniques for inductive definitions
- **Semantics**
 - Operational semantics
 - Modeling imperative “While” programs
 - Hoare logic for reasoning about program correctness
- **Type Systems**
 - Simply typed λ -calculus
 - Type safety
 - Subtyping
 - Dependently-typed programming
- **Coq interactive theorem prover**
 - turns doing proofs & logic into programming  fun!

COURSE MECHANICS

Administrivia

- Instructor: Benjamin Pierce
Office hours: Tuesdays, 4-5:30
Levine 562
- TAs:
 - Kenny Foner
Office hours: Fridays 10-noon
 - Antoine Voizard
Office hours: Monday afternoons
- Location: Berger Auditorium
- E-mail: cis500@seas.upenn.edu (goes to all course staff)
- Web site: <http://www.seas.upenn.edu/~cis500>
- Canvas: <https://upenn.instructure.com>
- PiazzaQ: <http://piazza.com/upenn/spring2016/cis500>

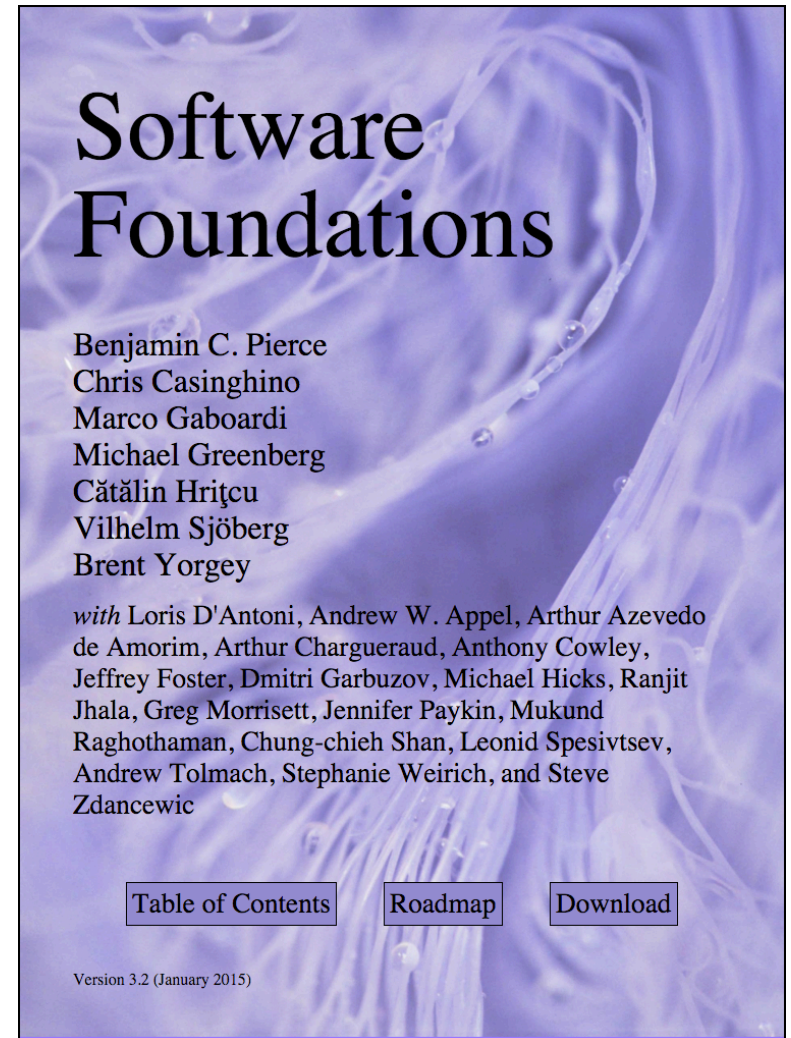
Resources

- Course textbook: *Software Foundations*
 - Electronic edition tailor-made for this class

Use the version available from the cis500 course web pages!!

(A new version of each chapter will generally go live just before class. :-)

- Additional resources:
 - *Types and Programming Languages* (Pierce, 2002 MIT Press)
 - *Interactive Theorem Proving and Program Development* (Bertot and Castéran, 2004 Springer)
 - *Certified Programming with Dependent Types* (Chlipala, electronic edition)



Course Policies

- Prerequisites:
 - Significant programming experience
 - Mathematical sophistication
 - Undergraduate functional programming or compilers class helpful

Grading:

- 24% Homework (~12 weekly assignments)
- 18% Midterm I (in class, probably Oct 4th)
- 18% Midterm 2 (in class, probably Nov. 8th)
- 36% Final (Tuesday, Dec. 20, noon-2PM)
- 4% Class participation

⇒ Lecture attendance is crucial!

“Regular” and “Advanced” tracks are graded separately

“Regular” vs. “Advanced” Tracks

- “Advanced” track:
 - More and harder exercises
 - More challenging exams
 - Covering a superset of the “regular” material
- Everybody starts in the advanced track by default.
- Students who wish to take CIS 500 for WPE I credit (Ph.D.) *must* complete the advanced track.
- Students may switch from advanced to regular track at any time.
 - Notify the course staff in writing (by e-mail).
 - The change is *permanent* after the first midterm.
- Students wishing to switch (back) to the advanced track:
 - Must do so *before* the first midterm exam.
 - Must make up all the advanced exercises (or accept the grade penalty).
- Only students taking the advanced track are eligible for an A+.

Class Participation

- Class attendance is mandatory.
- We will be using “clickers” for
 - in-class mini quizzes
 - Real-time “polls” during lectures
- TurningPoint clicker use will also be your attendance record.
- For next time: *buy a clicker.*
- *Any TurningPoint RF clicker will work; see note on course website.*



Homework Policies

- Homework is to be done *individually*
- Homework must be *submitted via Canvas*
- Homework that is late is subject to:
 - 25% penalty for 1 day late (up to 24 hours after deadline)
 - 50% penalty for 2 days late
 - 75% penalty for 3 days late
- Homework is due at *11:00am* on the due date
- Advanced track students must complete (or try to complete) all non-optional exercises.
 - Missing “advanced” exercises will count against your score.
- Regular track students must complete (or try to complete) all non-optional exercises except those marked “advanced”.
 - Missing “advanced” exercises will not count against your score.
 - (But you are welcome to try them!)

TODO for you

- Before next class:
 - Register for Piazza (if you are not already registered)
 - Try to log in to Canvas
 - Install Coq (version 8.5pl2, not 8.4 or 8.6)
 - Obtain a clicker
 - Start reading: Preface and Basics
- **HW1: Exercises in Basics.v**
 - Due: **Tuesday, September 6th at 11:00am**
 - Available from course web page
 - Complete all non-optional exercises
 - There are no “advanced” problems for this HW
 - Submit to Canvas



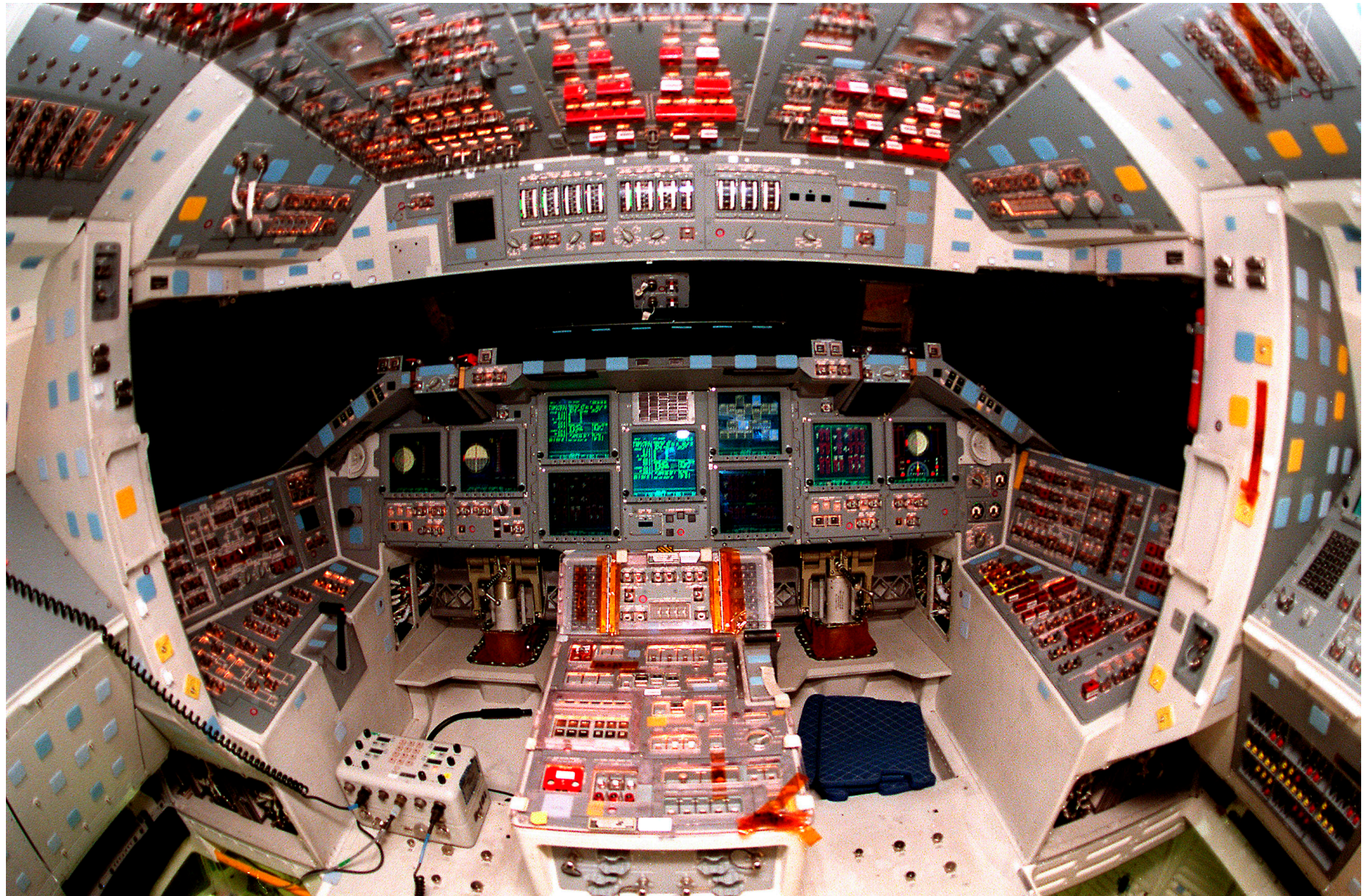
COQ

Coq in CIS 500

- We'll use Coq version 8.5
 - Available on CETS systems
 - Easy to install on your own machine
- See the web pages at: coq.inria.fr
- Two different user interfaces
 - CoqIDE – a standalone GUI / editor
 - ProofGeneral – an Emacs-based editing environment
- Course web pages have more information.



Coq's Full System



Subset Used in CIS 500



To start.



By the end of the semester.

Getting acquainted with Coq...

BASICS.V