

The Heap, Processes

Computer Systems Programming, Fall 2025

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The Dish, Stanford California

Administrivia

- ❖ First Assignment (HW00 penn-vector)
 - Released already! Should have everything you need after this lecture
 - “Due” Friday next week 09/05
 - Mostly a C refresher

- ❖ Pre semester Survey
 - Anonymous
 - Short!
 - Out Yesterday , Due Friday the 5th

Administrivia

- ❖ Second Assignment (HW01 penn-shredder)
 - Releases after Tuesday's lecture
 - Due September 12th
 - Intro to system calls, processes, etc.
 - Short Q&A and demo in lecture on Tuesday 😊

- ❖ Recitation: Starts next week! Not today!

Lecture Outline

❖ C “Refresher”

- Dynamic Memory vs the Stack
- Structs

❖ Processes

- Overview
- `fork()`
- `exec()`

Demo: `get_input.c`

- ❖ Lets code together a small program that:
 - Reads at max 100 characters from stdin (user input)
 - Truncates the input to only the first word
 - Prints that word out
 - Not allowed to use `scanf`, `FILE*`, `printf`, etc

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- ❖ Two things are wrong with this function. What are they?
- ❖ How do we fix this function w/o changing the function signature?

```
#define MAX_INPUT_SIZE 100

char* read_stdin() {
    char str[MAX_INPUT_SIZE];

    ssize_t res = read(STDIN_FILENO, str, MAX_INPUT_SIZE);

    // error checking
    if (res <= 0) {
        return NULL;
    }

    return str;
}

// assuming this is how the function is called
char* result = read_stdin();
```

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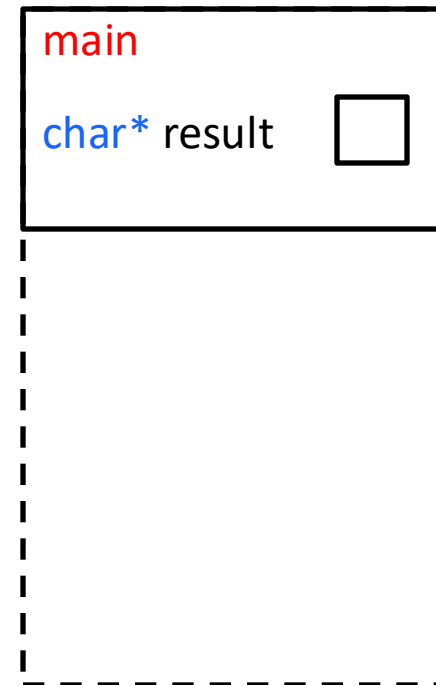
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The Stack



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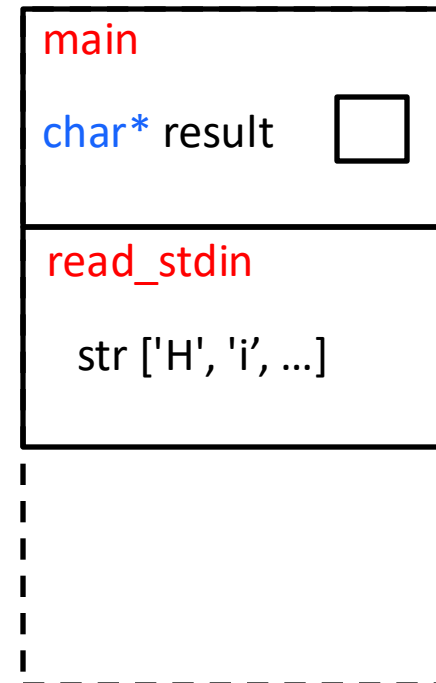
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The Stack



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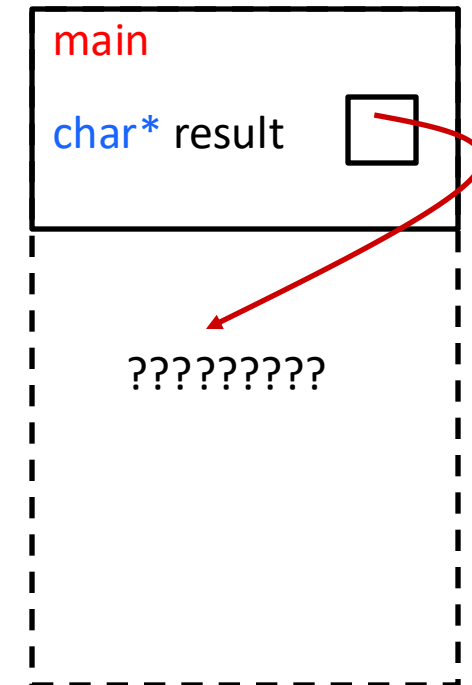
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        return NULL;
    }

    return str;
}

// assuming this is how the function is called
char* result = read_stdin();
```

The Stack



Memory Allocation

So far, we have seen two kinds of memory allocation:

```
int counter = 0;    // global var

int main() {
    counter++;
    printf("count = %d\n", counter);
    return 0;
}
```

- counter is **statically**-allocated
 - Allocated when program is loaded
 - Deallocated when program exits

```
int foo(int a) {
    int x = a + 1;    // local var
    return x;
}

int main() {
    int y = foo(10);  // local var
    printf("y = %d\n", y);
    return 0;
}
```

- a, x, y are **automatically**-allocated
 - Allocated when function is called
 - Deallocated when function returns



Aside: `sizeof`

- ❖ **`sizeof`** operator can be applied to a variable or a type and it evaluates to the size of that type in bytes
- ❖ Examples:
 - **`sizeof(int)`** – returns the size of an integer
 - **`sizeof(double)`** – returns the size of a double precision number
 - **`struct my_struct s;`**
 - **`sizeof(s)`** – returns the size of the struct `s`
 - **`my_type *ptr`**
 - **`sizeof (*ptr)`** – returns the size of the type pointed to by `ptr`
- ❖ Very useful for Dynamic Memory

What is Dynamic Memory Allocation?

- ❖ We want Dynamic Memory Allocation
 - Dynamic means “at run-time”
 - The compiler and the programmer don’t have enough information to make a final decision on how much to allocate
 - Your program explicitly requests more memory at run time
 - The language allocates it at runtime, maybe with help of the OS
- ❖ Dynamically allocated memory persists until either:
 - A garbage collector collects it (automatic memory management)
 - Your code explicitly deallocates it (manual memory management)
- ❖ C requires you to manually manage memory
 - More control, and more headaches

Heap API

- ❖ Dynamic memory is managed in a location in memory called the "Heap"
 - The heap is managed by user-level runtime library (libc)
 - Interface functions found in `<stdlib.h>`
- ❖ Most used functions:
 - `void *malloc(size_t size);`
 - Allocates memory of specified size
 - `void free(void *ptr);`
 - Deallocates memory
- ❖ Note: `void*` is “generic pointer”. It holds an address, but doesn’t specify what it is pointing at.
- ❖ Note 2: `size_t` is the integer type of `sizeof()`

malloc()

```
void *malloc(size_t size);
```

malloc allocates a block of memory of the requested size

- Returns a pointer to the first byte of that memory
 - And **returns NULL** if the memory allocation failed!
- You should assume that the memory initially contains garbage
- You'll typically use **sizeof** to calculate the size you need

```
// allocate a 10-float array
float* arr = malloc(10*sizeof(float));
if (arr == NULL) {
    return errcode;
}
... // do stuff with arr
```

ALWAYS CHECK FOR NULL

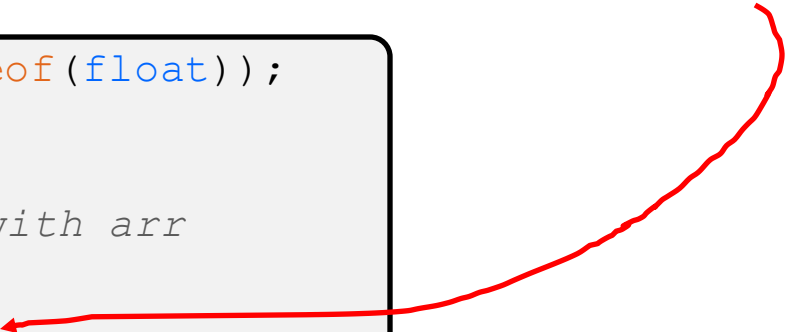
free()

```
free(pointer);
```

Deallocates the memory pointed-to by the pointer

- Pointer *must* point to the first byte of heap-allocated memory
 - (i.e. something previously returned by **malloc**)
- Freed memory becomes eligible for future allocation
- The bits in the pointer are *not changed* by calling free
 - Defensive programming: can set pointer to `NULL` after freeing it

```
float* arr = malloc(10*sizeof(float));  
if (arr == NULL)  
    return errcode;  
...           // do stuff with arr  
free(arr);  
arr = NULL;   // Preferred
```



```
free(NULL);
```

p.s. This is a No-Op.

The Heap

- ❖ The Heap is a large pool of available memory to use for Dynamic allocation
- ❖ This pool of memory is kept track of with a small data structure indicating which portions have been allocated, and which portions are currently available.
- ❖ **malloc:**
 - searches for a large enough unused block of memory
 - marks the memory as allocated.
 - Returns a pointer to the beginning of that memory
- ❖ **free:**
 - Takes in a pointer to a previously allocated address
 - Marks the memory as free to use.

Dynamic Memory Example

```
#include <stdlib.h>

int main() {
    char* ptr = malloc(4*sizeof(char));
    if (ptr == NULL)
        return EXIT_FAILURE;
    ...           // do stuff with ptr
    free(ptr);
}
```

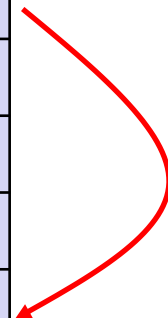
addr	var	value
0x2001	ptr	--
...	...	--
0x4000	HEAP START	USED
0x4001		USED
0x4002		Free
0x4003		Free
0x4004		Free
0x4005		Free
0x4006		Free
0x4007		Free
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0x4009		USED

Dynamic Memory Example

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    free(ptr);
}
```

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...	...	--
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0x4003	e	USED
0x4004	y	USED
0x4005	\0	USED
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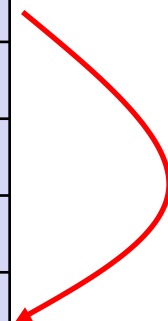


Dynamic Memory Example

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Demo (continued): `get_input.c`

- ❖ Lets code together a small program that:
 - Reads at max 100 characters from stdin (user input)
 - Truncates the input to only the first word
 - Prints that word out
 - Not allowed to use `scanf`, `FILE*`, `printf`, etc

- ❖ Let's fix the Stack Array issue!
- ❖ What was the other issue? (other than not using `malloc`)

Partially Fixed read_stdin()

```
#define MAX_INPUT_SIZE 100

char* read_stdin() {
    char str = (char*) malloc(sizeof(char) * MAX_INPUT_SIZE);
    if (str == NULL) {
        return NULL;
    }

    ssize_t res = read(STDIN_FILENO, str, MAX_INPUT_SIZE);

    // error checking
    if (res <= 0) {
        return NULL;
    }

    return str;
}
```

Fully Fixed read_stdin()

```
#define MAX_INPUT_SIZE 100

char* read_stdin() {
    char str = (char*) malloc(sizeof(char) * MAX_INPUT_SIZE);
    if (str == NULL) {
        return NULL;
    }

    ssize_t res = read(STDIN_FILENO, str, MAX_INPUT_SIZE);

    // error checking
    if (res <= 0) {
        return NULL;
    }
    str[res] = '\0';

    return str;
}
```

Reminder: read is a very *exact function* in that it only writes exactly what there is to read.

The Null-Terminator must be added manually.

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❖ Which function works as intended?

```
typedef struct point_st {  
    float x;  
    float y;  
} Point;
```

```
Point* make_point() {  
    Point p = (Point) {  
        .x = 2.0f;  
        .y = 1.0f;  
    };  
    Point* ptr = &p;  
    return ptr;  
}
```

A

```
typedef struct point_st {  
    float x;  
    float y;  
} Point;
```

```
Point make_point() {  
    Point p = (Point) {  
        .x = 2.0f;  
        .y = 1.0f;  
    };  
    return p;  
}
```

B

Dynamic Memory Pitfalls

- ❖ Buffer Overflows
 - E.g. ask for 10 bytes, but write 11 bytes
 - Could overwrite information needed to manage the heap
 - Common when forgetting the null-terminator on malloc'd strings
- ❖ Not checking for **NULL**
 - Malloc returns NULL if out of memory
 - Should check this after every call to malloc
- ❖ Giving **free()** a pointer to the middle of an allocated region
 - Free won't recognize the block of memory and probably crash
- ❖ Giving **free()** a pointer that has already been freed
 - Will interfere with the management of the heap and likely crash
- ❖ **malloc** does NOT initialize memory
 - There are other functions like **calloc** that will zero out memory

Memory Leaks

- ❖ The most common Memory Pitfall
- ❖ What happens if we malloc something, but don't free it?
 - That block of memory cannot be reallocated, even if we don't use it anymore, until it is **freed**
 - If this happens enough, we run out of heap space and program may slow down and eventually crash
- ❖ Garbage Collection
 - C doesn't have this. You're on your own.

Discuss: What is wrong with this code? (Multiple bugs)

```
int main() {
    char* literal = "Hello!";
    char* duplicate = dup_str(literal);
    char* ptr = duplicate;

    while (*ptr != '\0') {
        printf("%s\n", ptr);
        // printf line is fine
        ptr += 1;
    }

    free(duplicate);
    free(ptr);
    free(literal);
}
```

```
char* dup_str(char* to_copy) {
    size_t len = strlen(to_copy);
    char* res = malloc(sizeof(char) * len);
    for (size_t i = 0; i < len; i++) {
        res[i] = to_copy[i];
    }
    return res;
}
```

strlen()

returns the number of characters before the null-terminator

static function variables

❖ Functions can declare a variable as static

```
#include <stdio.h> // for printf
#include <stdlib.h> // for EXIT_SUCCESS
```

```
int next_num();
```

This is how some functions
can “remember” things.

```
int main(int argc, char** argv) {
    printf("%d\n", next_num()); // prints 1
    printf("%d\n", next_num()); // then 2
    printf("%d\n", next_num()); // then 3
    return EXIT_SUCCESS;
}
```

```
int next_num() {
    // marking this variable as static means that
    // the value is preserved between calls to the function
    // this allows the function to “remember” things
    static int counter = 0;
    counter++;
    return counter;
}
```

Can be thought of as a
global variable that is
“private” to a function

Lecture Outline

❖ C “Refresher”

- Dynamic Memory vs the Stack
- Structs

❖ **Processes**

- **Overview**
- **fork()**
- **exec()**

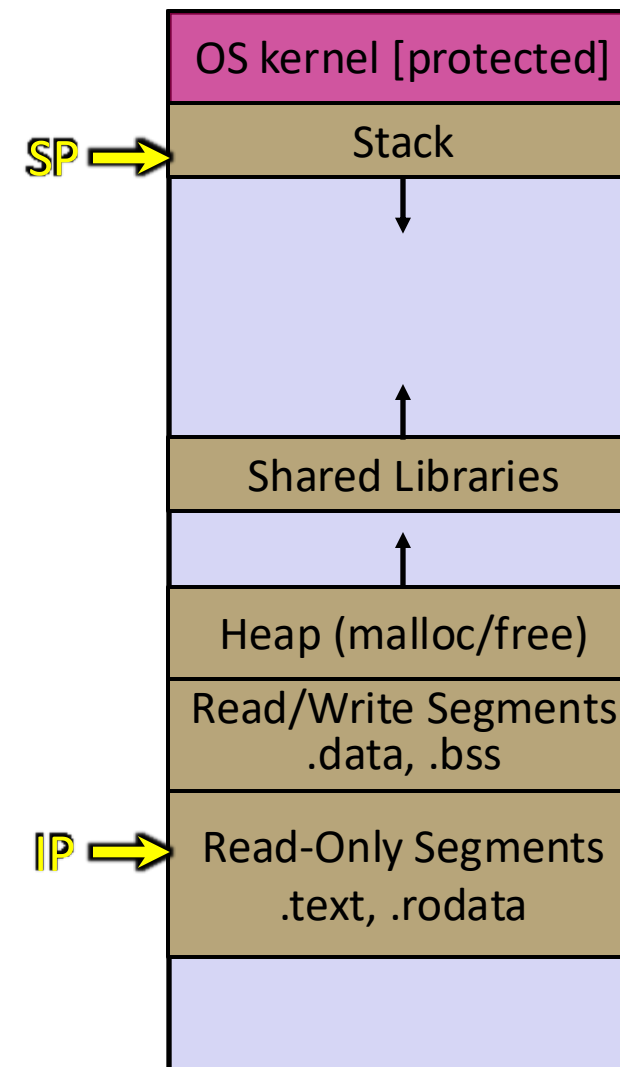
Definition: Process

Definition: An instance of a program that is being executed (or is ready for execution)

Consists of:

- Memory (code, heap, stack, etc)
- Registers used to manage execution (stack pointer, program counter, ...)
- Other resources

** This isn't quite true
more in a future lecture*

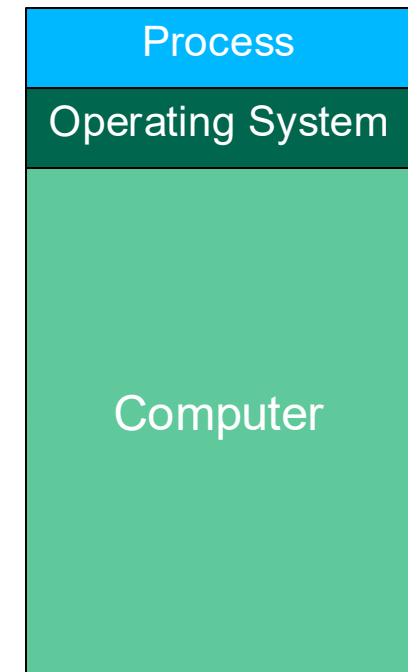


Computers as we know them now

- ❖ In CIS 2400, you learned about hardware, transistors, CMOS, gates, etc.
- ❖ Once we got to programming, our computer looks something like:

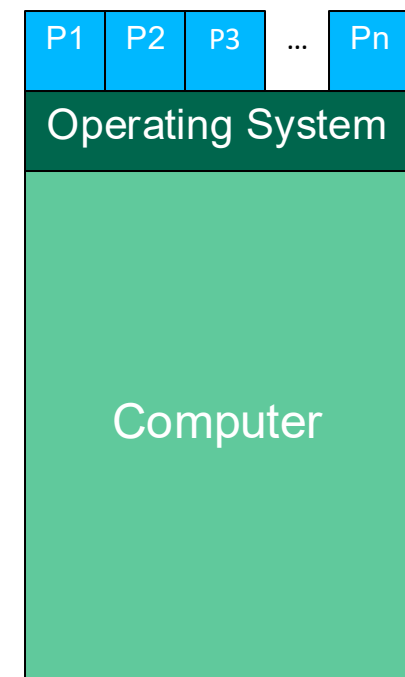
What is missing/wrong with this?

- ❖ This model is still useful, and can be used in many settings



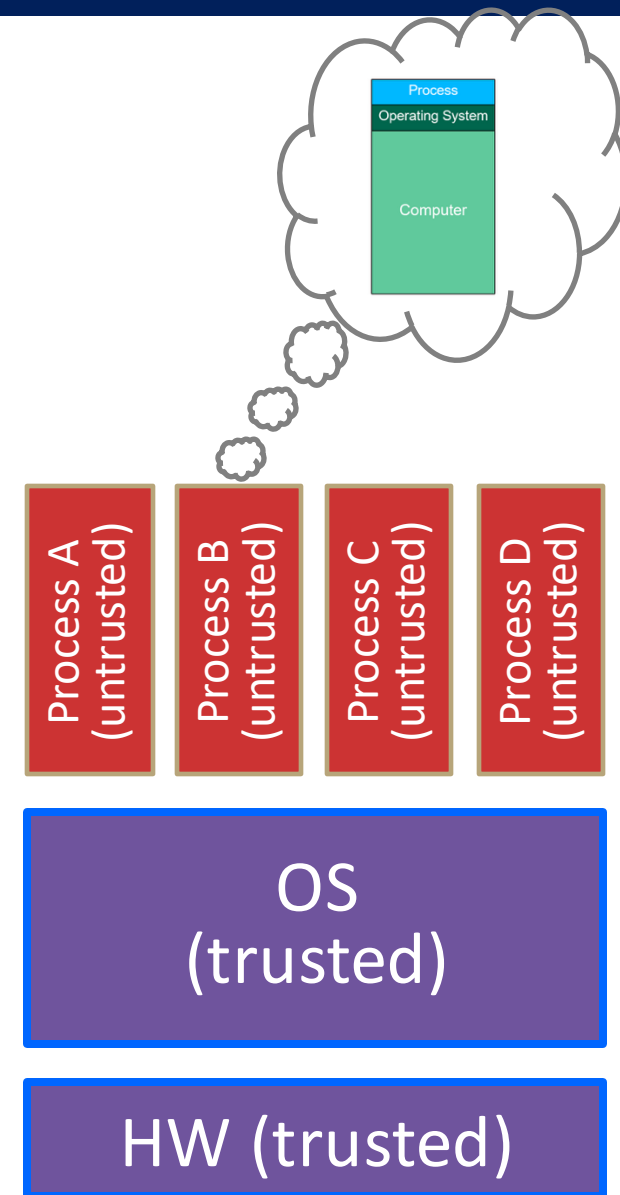
Multiple Processes

- ❖ Computers run multiple processes “at the same time”
- ❖ One or more processes for each of the programs on your computer
- ❖ Each process has its own...
 - Memory space
 - Registers
 - Resources

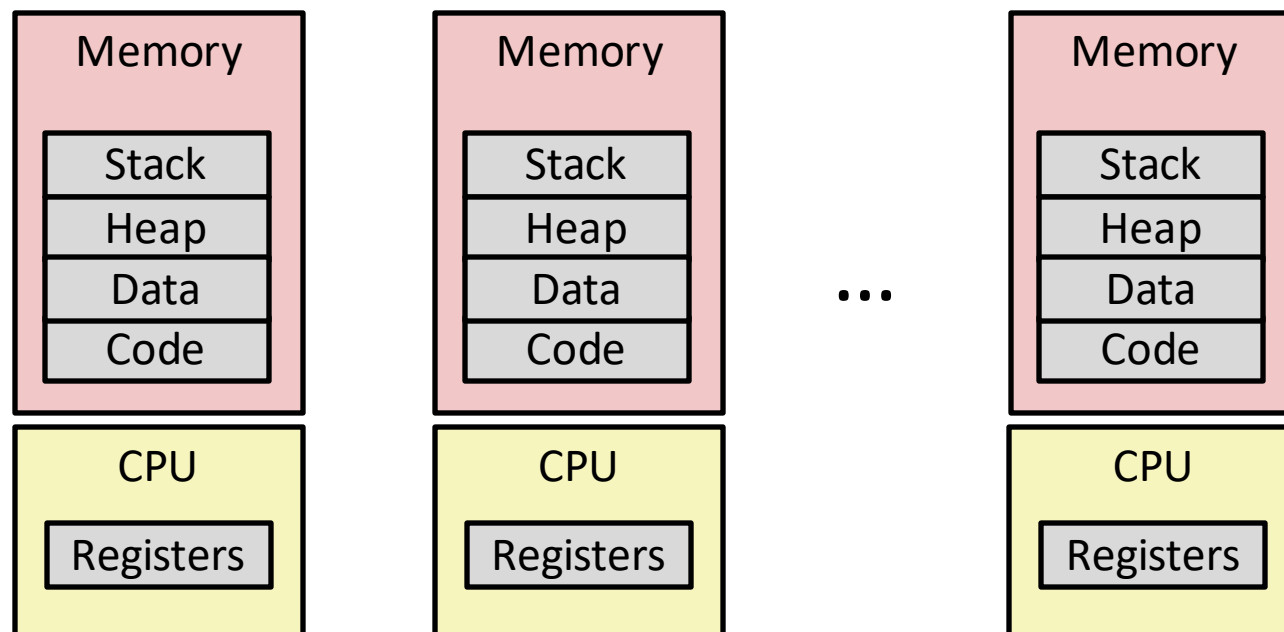


OS: Protection System

- ❖ OS isolates process from each other
 - Each process seems to have exclusive use of memory and the processor.
 - This is an **illusion**
 - More on Memory when we talk about virtual memory later in the course
 - OS permits controlled sharing between processes
 - E.g. through files, the network, etc.
- ❖ OS isolates itself from processes
 - Must prevent processes from accessing the hardware directly



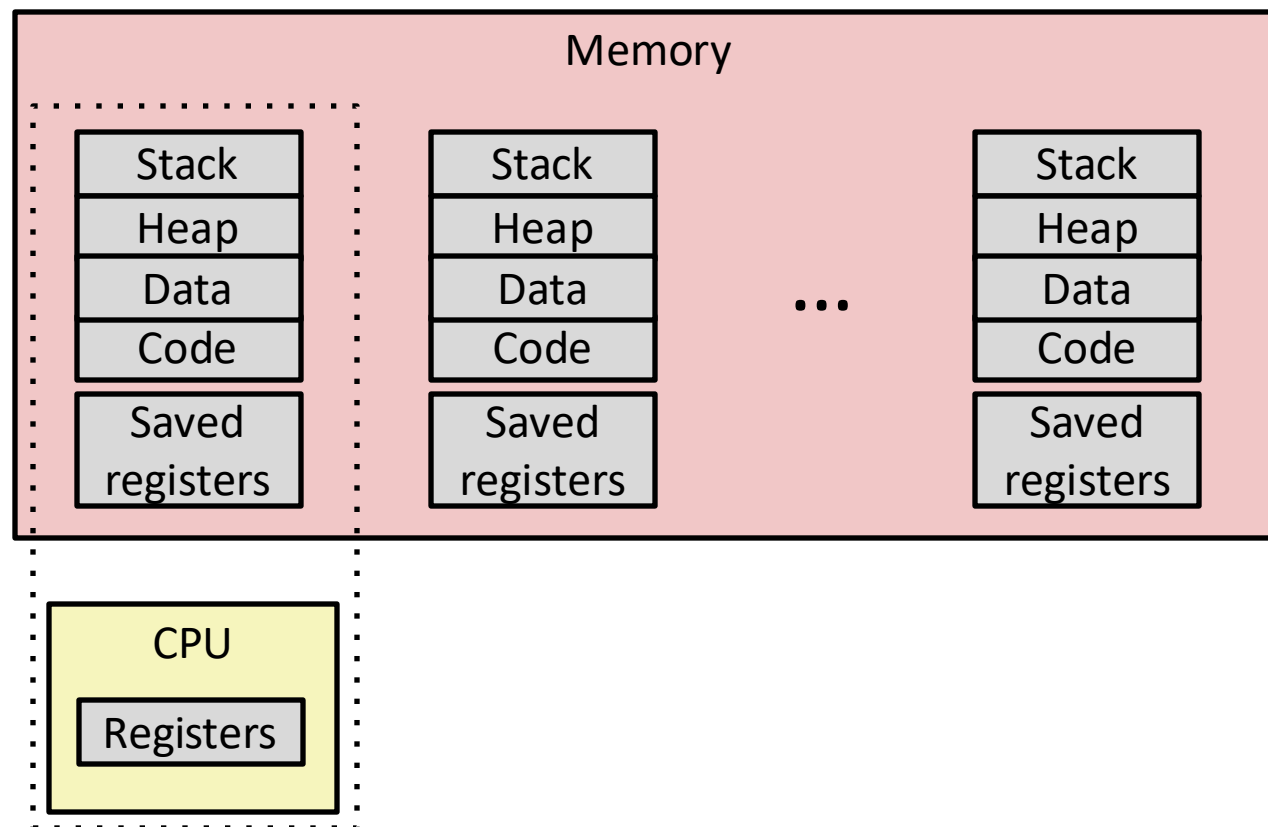
Multiprocessing: The Illusion



❖ Computer runs many processes simultaneously

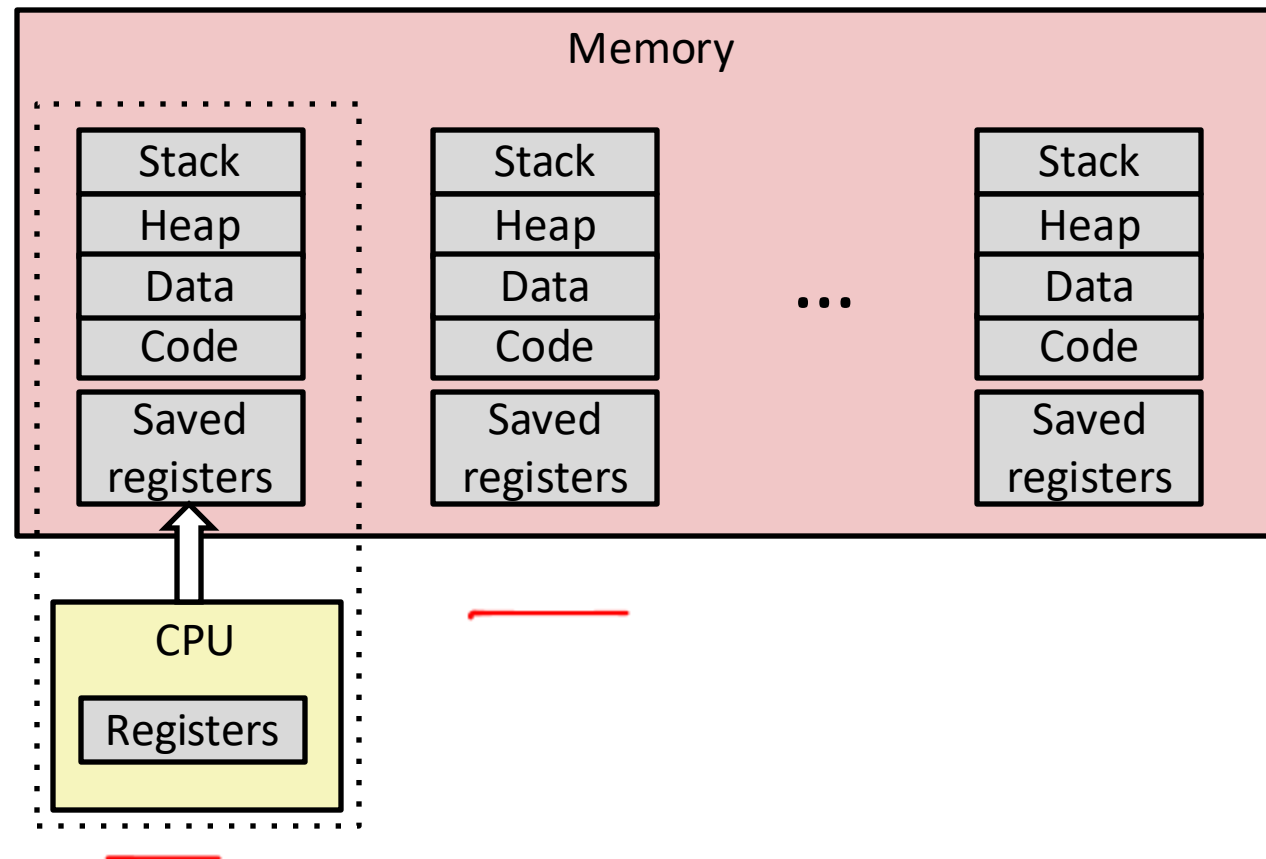
- Applications for one or more users
 - Web browsers, email clients, editors, ...
- Background tasks
 - Monitoring network & I/O devices

Multiprocessing: The (Traditional) Reality



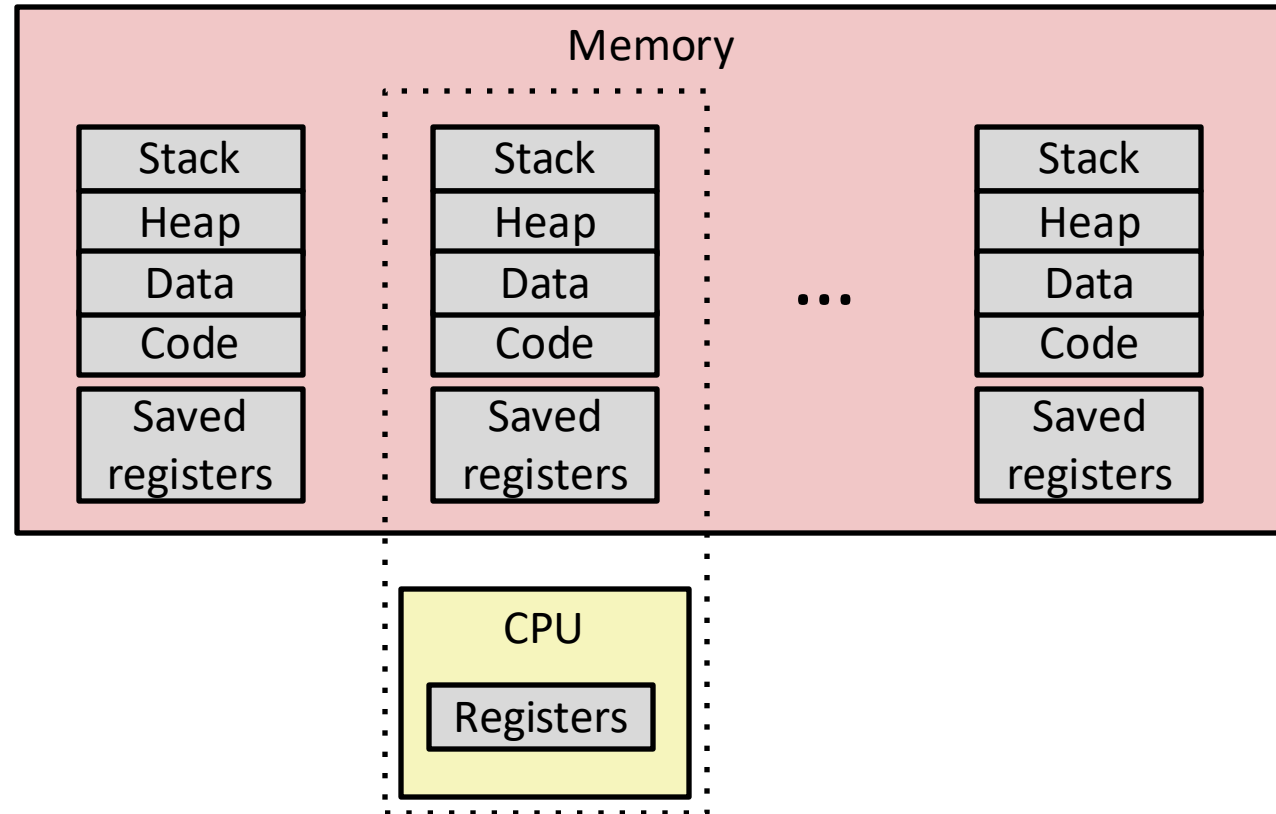
- ❖ Single processor executes multiple processes concurrently
 - Process executions interleaved (multitasking)
 - Address spaces managed by virtual memory system (later in course)
 - Register values for nonexecuting processes saved in memory

Multiprocessing: The (Traditional) Reality



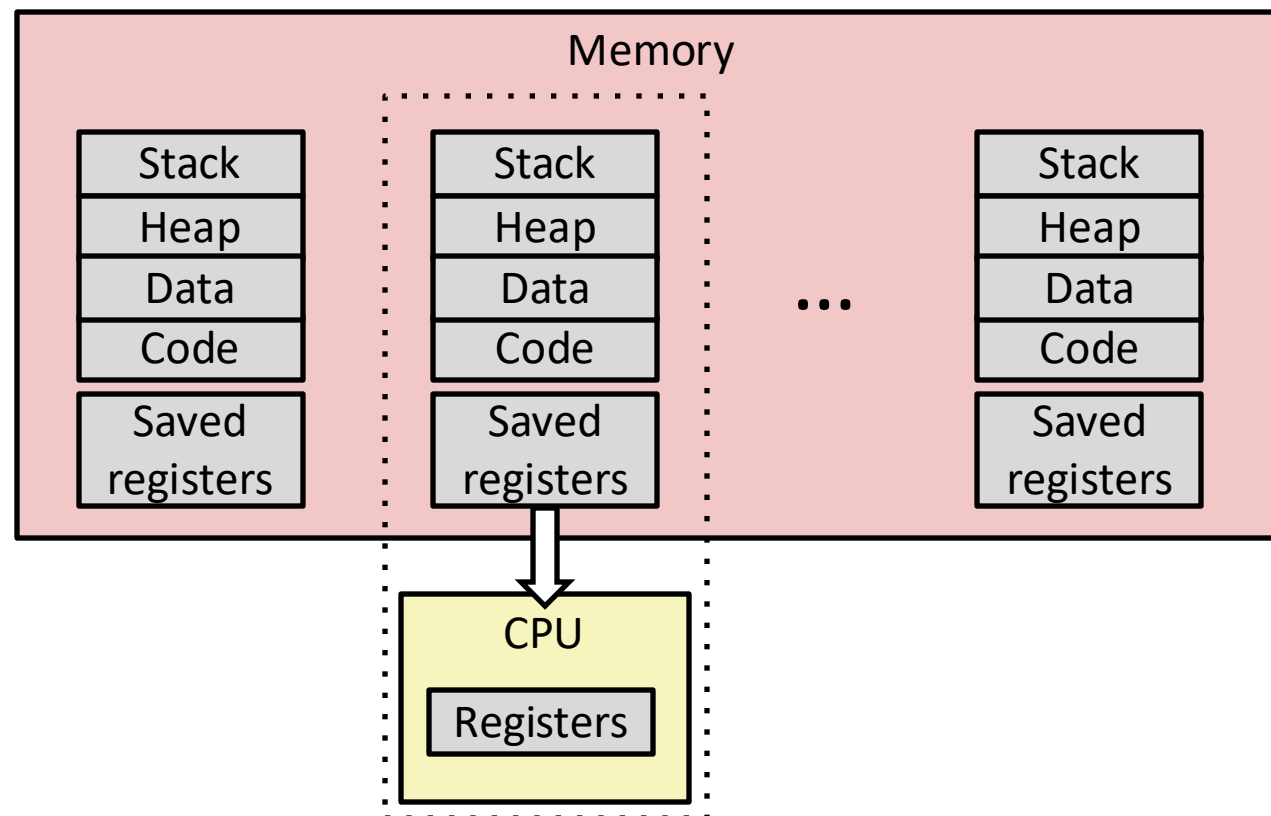
1. Save current registers in memory

Multiprocessing: The (Traditional) Reality



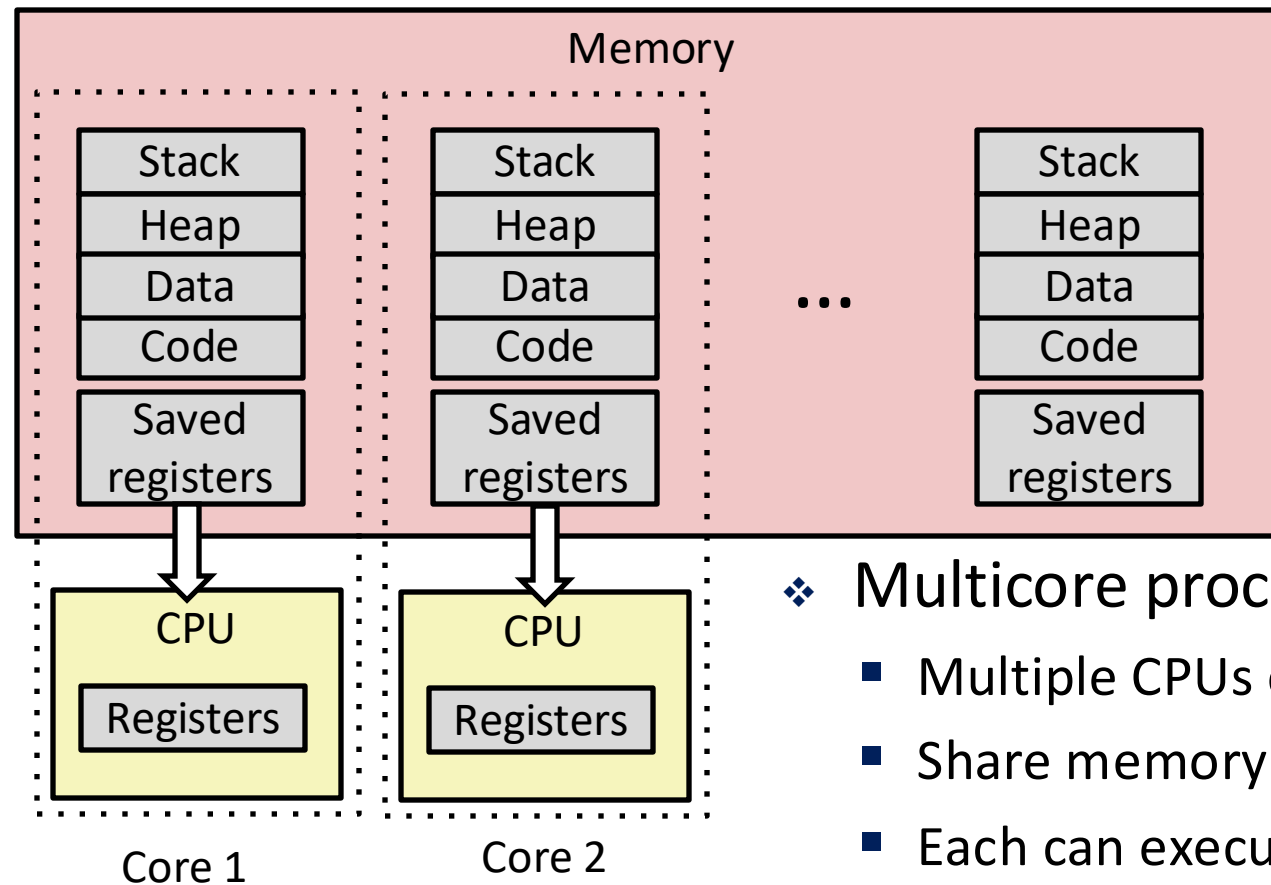
1. Save current registers in memory
2. **Schedule next process for execution**

Multiprocessing: The (Traditional) Reality



1. Save current registers in memory
2. **Schedule next process for execution**
3. Load saved registers and switch address space (**context switch**)

Multiprocessing: The (Traditional) Reality



- ❖ Multicore processors
 - Multiple CPUs on single chip
 - Share memory
 - Each can execute a separate process
 - Scheduling of processors onto cores done by kernel
 - This is called “Parallelism”

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- ❖ What I just went through was the big picture of processes. Many details left, some will be gone over in future lectures
- ❖ Any questions, comments or concerns so far?

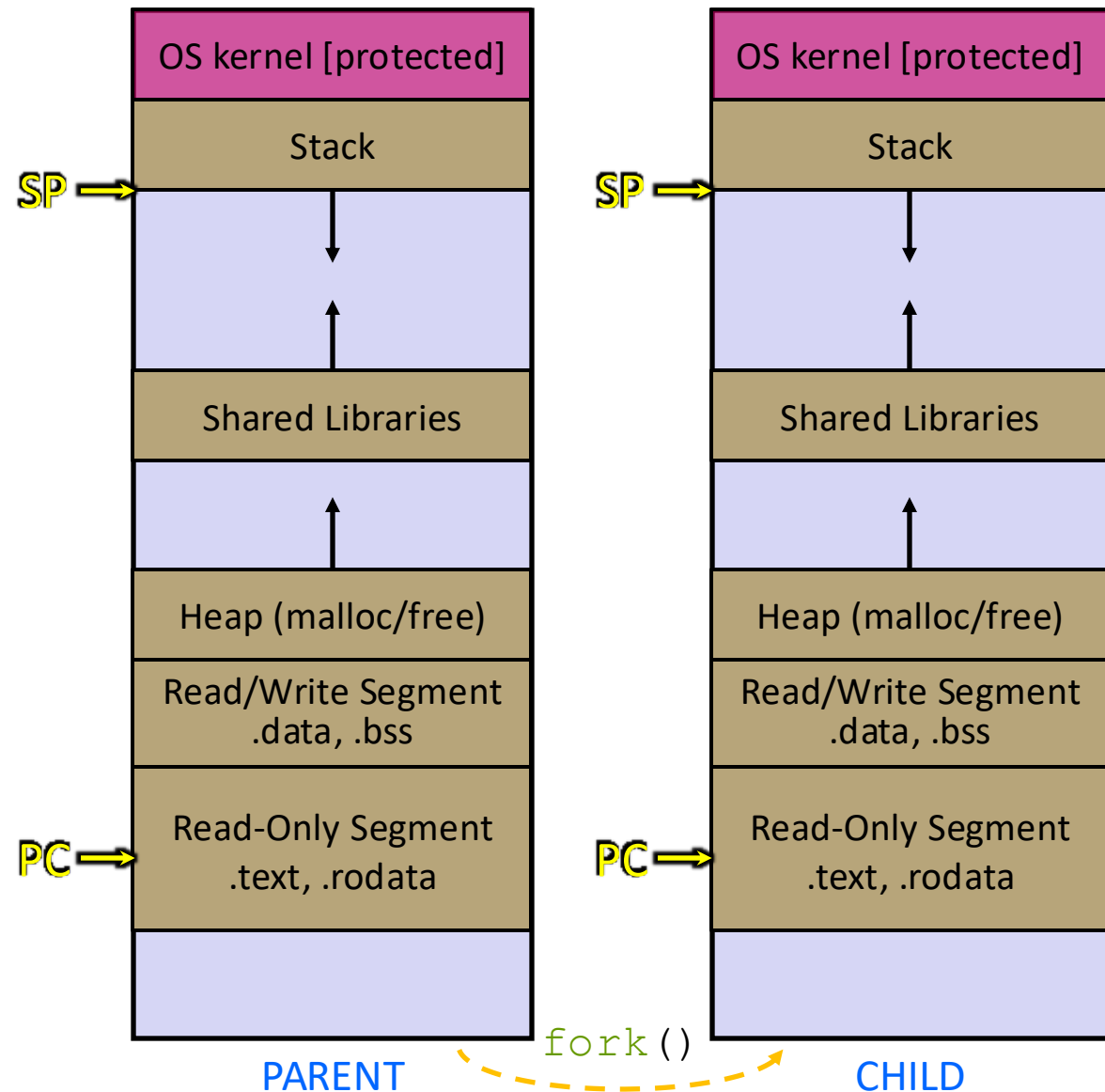
Creating New Processes

```
pid_t fork();
```

- Creates a new process (the “child”) that is an *exact clone** of the current process (the “parent”)
 - *almost everything
- The new process has a separate virtual address space from the parent
- Returns a `pid_t` which is an integer type.
 - The parent receives the *real pid*
 - The child receives **0**.

fork () and Address Spaces

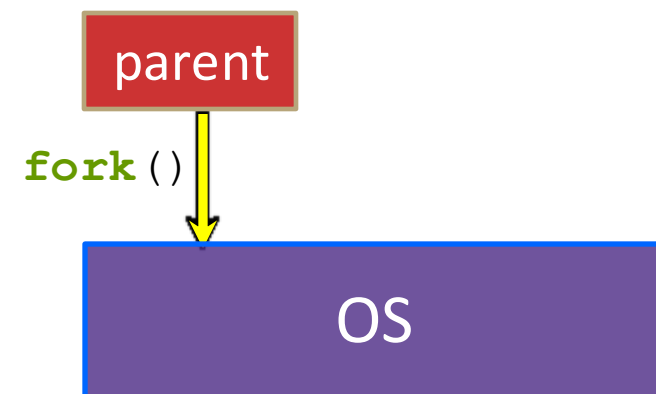
- ❖ Fork causes the OS to clone the address space
 - The *copies* of the memory segments are (nearly) identical
 - The new process has *copies* of the parent's data, stack-allocated variables, open file descriptors, etc.



fork()

❖ **fork()** has peculiar semantics

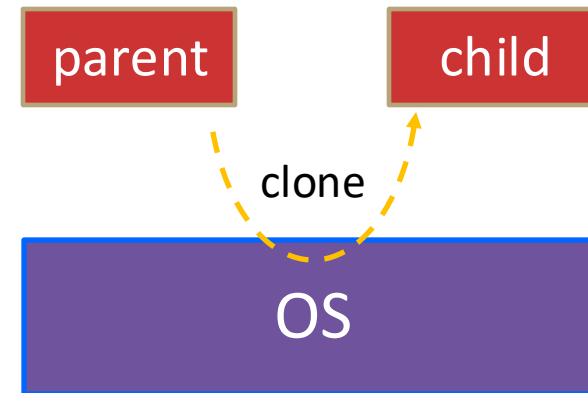
- The parent invokes **fork()**
- The OS clones the parent
- *Both* the parent and the child return from fork
 - Parent receives child's pid
 - Child receives a 0



fork()

❖ **fork()** has peculiar semantics

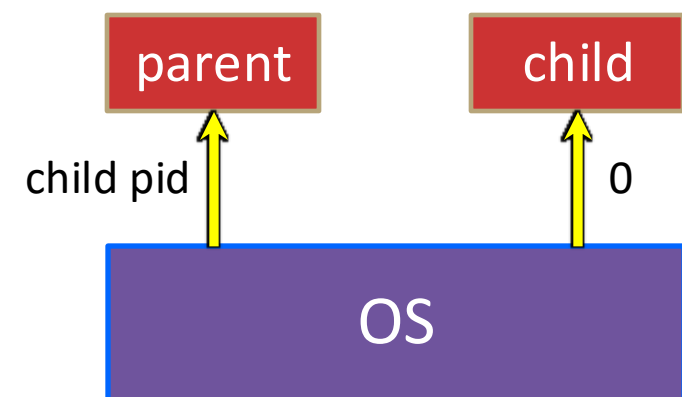
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fork()

❖ **fork()** has peculiar semantics

- The parent invokes **fork()**
- The OS clones the parent
- *Both* the parent and the child return from fork
 - Parent receives child's pid
 - Child receives a 0



"simple" `fork()` example

```
fork();  
printf("Hello!\n");
```

❖ What does this print?

"simple" `fork()` example

Parent Process (PID = X)



```
fork();  
printf("Hello!\n");
```

Child Process (PID = Y)



```
fork();  
printf("Hello!\n");
```

❖ What does this print?

❖ "Hello!\n" is printed twice

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```
fork();  
fork();  
printf("Hello!\n");
```

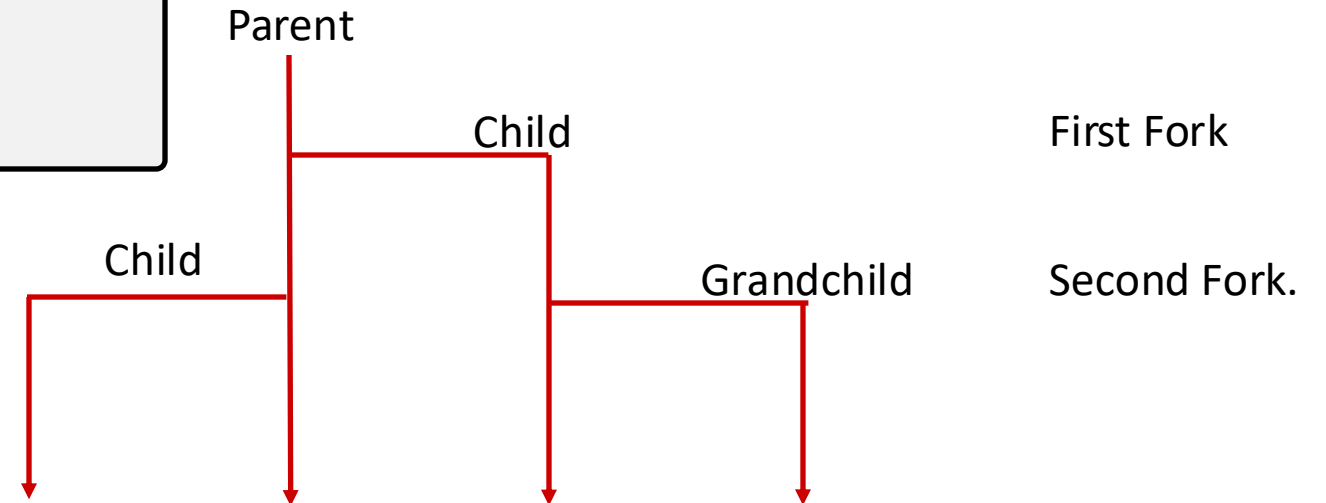
❖ What does this print?

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```
fork();  
fork();  
printf("Hello!\n");
```

❖ What does this print?

“hello!” four times!



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```
int x = 3;  
fork();  
x++;  
printf("%d\n", x);
```

❖ What does this print?

 Poll Everywhere[discusssss](#)

```
→ pid_t fork_ret = fork();  
  
if (fork_ret == 0) {  
    printf("Child\n");  
} else {  
    printf("Parent\n");  
}
```

❖ What does this print?

fork() example

Parent Process (PID = X)

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→ pid_t fork_ret = fork();  
  
if (fork_ret == 0) {  
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} else {  
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```

Child Process (PID = Y)

```
→ pid_t fork_ret = fork();  
  
if (fork_ret == 0) {  
    printf("Child\n");  
} else {  
    printf("Parent\n");  
}
```

fork()

fork() example

Parent Process (PID = X)

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Child Process (PID = Y)

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→ pid_t fork_ret = fork();  
  
if (fork_ret == 0) {  
    printf("Child\n");  
} else {  
    printf("Parent\n");  
}
```

fork_ret = Y

```
pid_t fork_ret = fork();  
  
if (fork_ret == 0) {  
    printf("Child\n");  
} else {  
→ printf("Parent\n");  
}
```

Prints "Parent"

fork_ret = 0

```
pid_t fork_ret = fork();  
  
if (fork_ret == 0) {  
→ printf("Child\n");  
} else {  
    printf("Parent\n");  
}
```

Prints "Child"

Which prints first?

Process States (incomplete)

FOR NOW, we can think of a process as being in one of three states:

- ❖ Running

- Process is currently executing

*More states in
future lectures*

- ❖ Ready

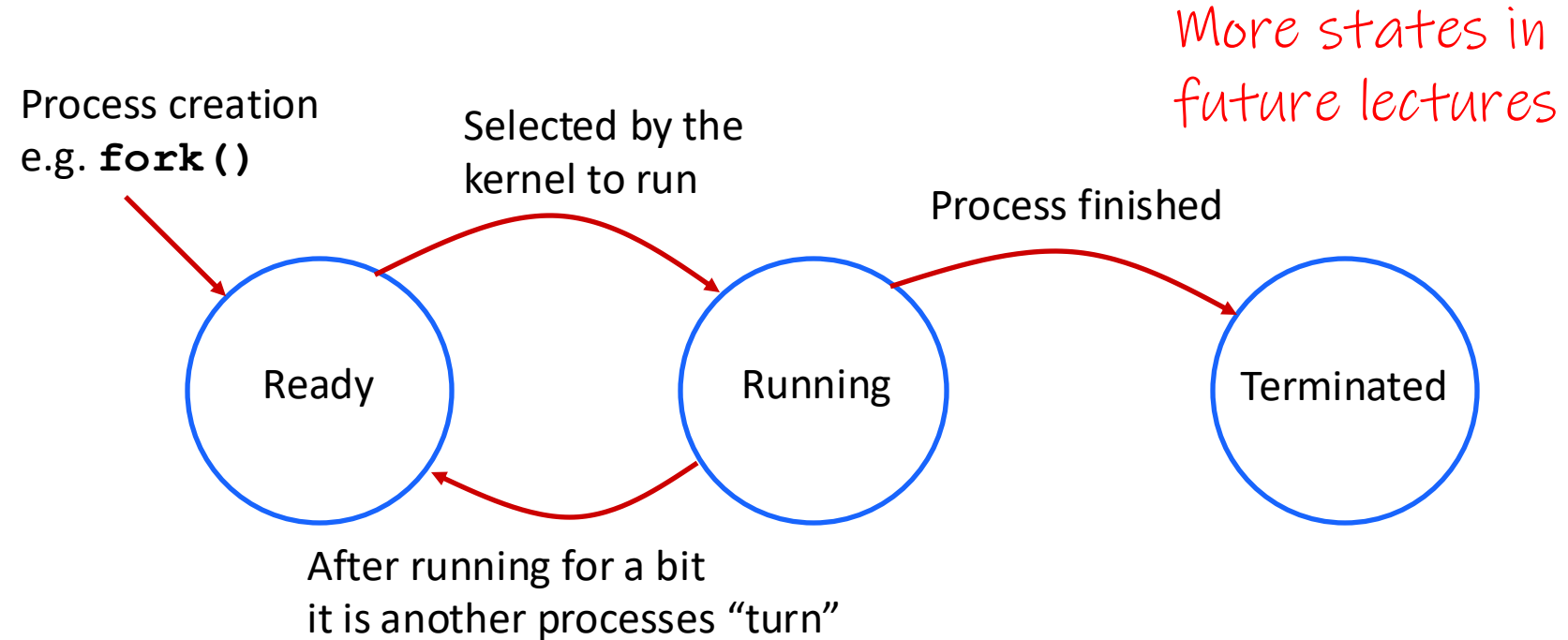
- Process is waiting to be executed and will eventually be *scheduled* (i.e., chosen to execute) by the kernel

*Scheduler to be covered
in a later lecture*

- ❖ Terminated

- Process is stopped permanently

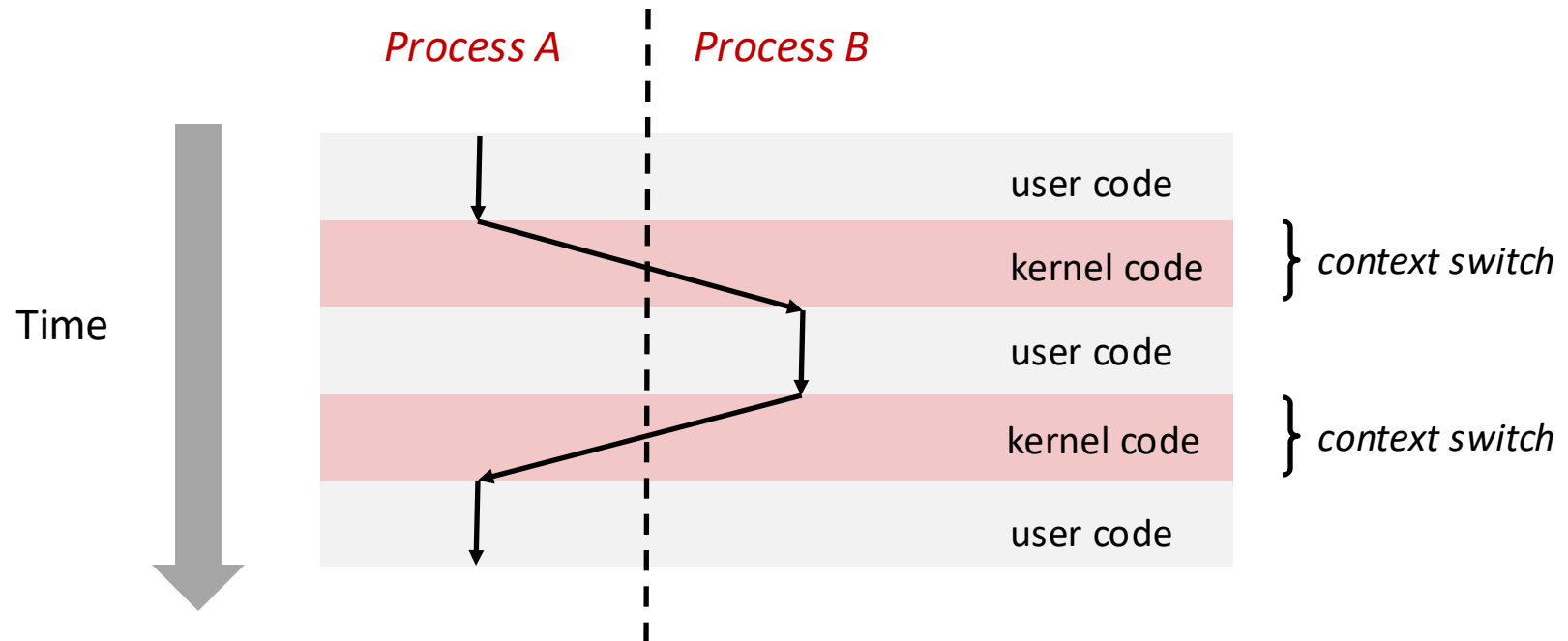
Process State Lifetime (incomplete)



Processes can be "interrupted" to stop running. Through something like a hardware timer interrupt

Context Switching

- ❖ Processes are managed by a shared chunk of memory-resident OS code called the *kernel*
 - Important: the kernel is not a separate process, but rather runs as part of some existing process.
- ❖ Control flow passes from one process to another via a *context switch*



OS: The Scheduler

- ❖ When switching between processes, the OS will run some kernel code called the “Scheduler”
- ❖ The scheduler runs when a process:
 - starts (“arrives to be scheduled”),
 - Finishes
 - Blocks (e.g., waiting on something, usually some form of I/O)
 - Has run for a certain amount of time
- ❖ It is responsible for scheduling processes
 - Choosing which one to run
 - Deciding how long to run it

Scheduler Considerations

- ❖ The scheduler has a scheduling algorithm to decide what runs next.
- ❖ Algorithms are designed to consider many factors:
 - Fairness: Every program gets to run
 - Liveness: That “something” will eventually happen
 - Throughput: Number of “tasks” completed over an interval of time
 - Wait time: Average time a “task” is “alive” but not running
 - A lot more...
- ❖ More on this later. **For now: think of scheduling as non-deterministic**, details handled by the OS.

fork() example

```
→ printf("Hello!\n");  
pid_t fork_ret = fork();  
int x;  
  
if (fork_ret == 0) {  
    x = 1234;  
} else {  
    x = 5678;  
}  
printf("%d\n", x);
```

Always prints "Hello"

fork() example

```
printf("Hello!\n");  
→ pid_t fork_ret = fork();  
int x;  
  
if (fork_ret == 0) {  
    x = 1234;  
} else {  
    x = 5678;  
}  
printf("%d\n", x);
```

Always prints "Hello"

fork() example

Parent Process (PID = X)

```
printf("Hello!\n");
→ pid_t fork_ret = fork();
int x;

if (fork_ret == 0) {
    x = 1234;
} else {
    x = 5678;
}

printf("%d\n", x);
```

fork_ret = Y

Always prints "Hello"

Child Process (PID = Y)

```
printf("Hello!\n");
→ pid_t fork_ret = fork();
int x;

if (fork_ret == 0) {
    x = 1234;
} else {
    x = 5678;
}

printf("%d\n", x);
```

fork_ret = 0

Does NOT print "Hello"

fork()

fork() example

Parent Process (PID = X)

```
printf("Hello!\n");
pid_t fork_ret = fork();
int x;

if (fork_ret == 0) {
    x = 1234;
} else {
    x = 5678;
}

printf("%d\n", x);
```

fork_ret = Y

Child Process (PID = Y)

```
printf("Hello!\n");
pid_t fork_ret = fork();
int x;

if (fork_ret == 0) {
    x = 1234;
} else {
    x = 5678;
}

printf("%d\n", x);
```

fork_ret = 0

fork()

Always prints "Hello"

Always prints "5678"

Always prints "1234"

Exiting a Process

```
void exit(int status);
```

- Causes the current process to exit normally
- Automatically called by **main()** when main returns
- Exits with a return status (e.g. **EXIT_SUCCESS** or **EXIT_FAILURE**)
 - This is the same int returned by **main()**
- The exit status is accessible by the parent process with **wait()** or **waitpid()**.

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```
int global_num = 1;

void function() {
    global_num++;
    printf("%d\n", global_num);
}

int main() {
    pid_t id = fork();

    if (id == 0) {
        function();
        id = fork();
        if (id == 0) {
            function();
        }
        return EXIT_SUCCESS;
    }

    global_num += 2;
    printf("%d\n", global_num);
    return EXIT_SUCCESS;
}
```

- ❖ How many numbers are printed?
- ❖ What number(s) get printed from each process?

Poll Everywhere

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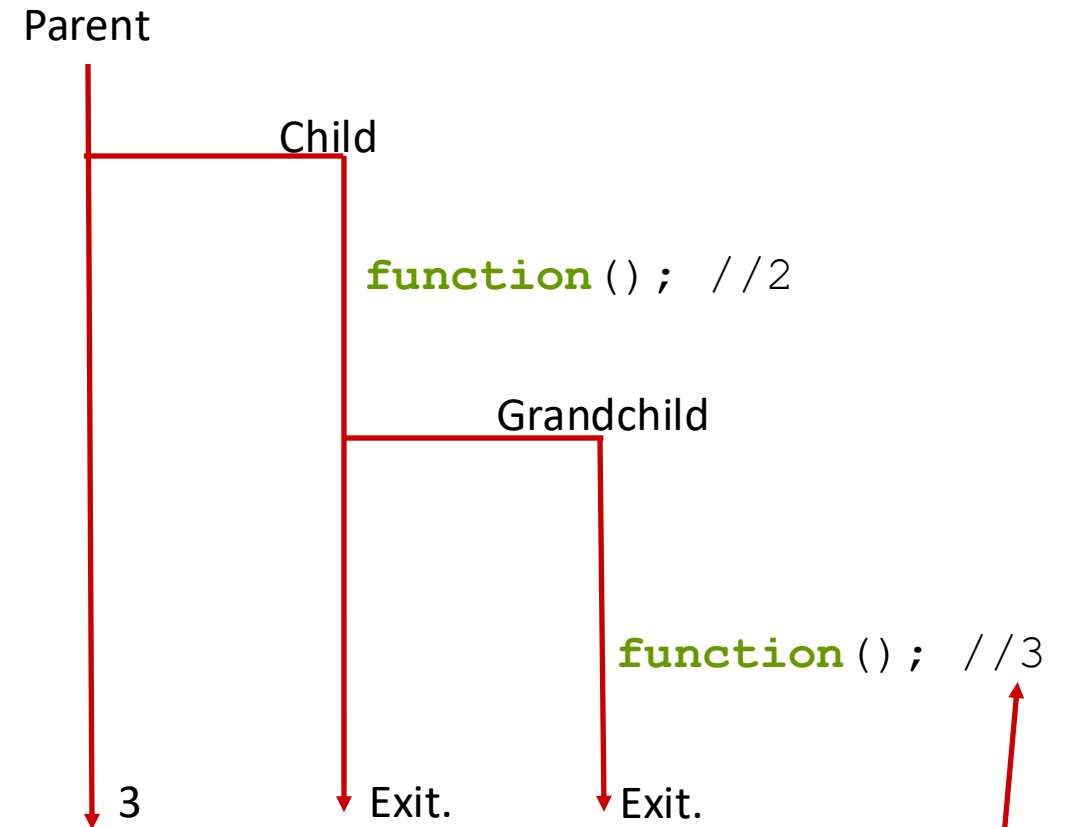
```
int global_num = 1;

void function() {
    global_num++;
    printf("%d\n", global_num);
}

int main() {
    pid_t id = fork();

    if (id == 0) {
        function();
        id = fork();
        if (id == 0) {
            function();
        }
        return EXIT_SUCCESS;
    }

    global_num += 2;
    printf("%d\n", global_num);
    return EXIT_SUCCESS;
}
```



*Remember, the child starts with an identical copy of the it's parents memory.

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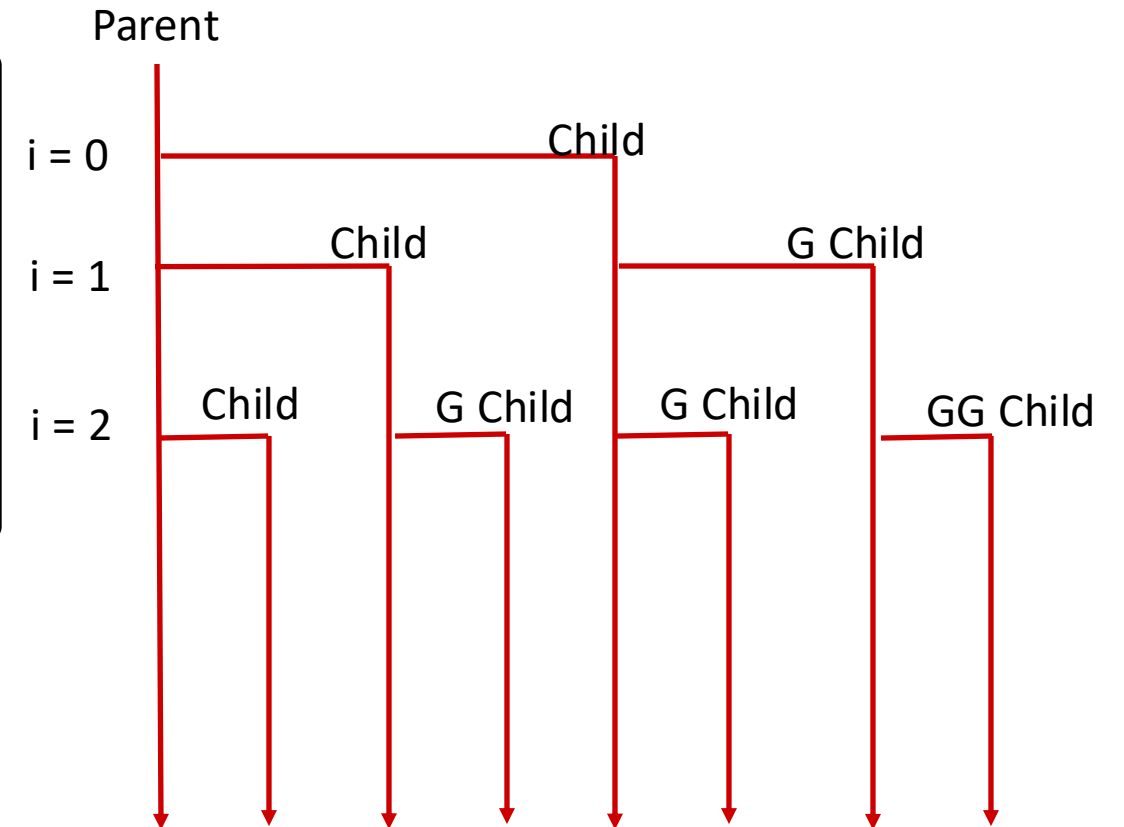
❖ How many times is ":)" printed?

```
int main(int argc, char* argv[]) {  
    for (int i = 0; i < 4; i++) {  
        fork();  
    }  
  
    printf(":) \n"); // "\n" is similar to endl  
    return EXIT_SUCCESS;  
}
```

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❖ How many times is ":)" printed?

```
int main(int argc, char* argv[]) {  
    for (int i = 0; i < 4; i++) {  
        fork();  
    }  
  
    printf(":) \n"); // "\n" is similar to endl  
    return EXIT_SUCCESS;  
}
```



Processes & Fork Summary

- ❖ Processes are instances of programs that:
 - Each have their own independent ***address space (more on that later!)***
 - Each process is scheduled by the OS
 - Without using some functions we have not talked about (yet), there is no way to guarantee the order processes are executed
 - Processes are created by `fork()` system call
 - Only difference between processes is their process id and the return value from `fork()` each process gets

Lecture Outline

❖ C “Refresher”

- Dynamic Memory vs the Stack
- Structs

❖ **Processes**

- Overview
- `fork()`
- **`exec()`**

That's all!

- ❖ See y'all on Tuesday! Enjoy your extended weekend!
- ❖ Check-In 00 will be released tonight!