# CIS 5480 Midterm Review

February 27, 2025

Jeopardy template made by Eric Curts

#### **FINAL JEOPARDY**

#### **JEOPARDY BOARD**

Process Management	Signals	File Descriptors & I/O	File Systems	Concurrency & Synchronization	CPU Scheduling
<b>\$100</b>	<b>\$100</b>	\$100	\$100	<b>\$100</b>	<b>\$100</b>
<b>\$200</b>	<b>\$200</b>	<b>\$200</b>	<b>\$200</b>	<b>\$200</b>	<b>\$200</b>
\$300	\$390	\$300	\$300	\$300	\$300
\$400	\$400	\$400	<b>\$400</b>	\$400	\$400
\$500	\$500	\$500	\$500	\$500	\$500

#### Process Management - \$100 Question

What does the wait() system call do in a parent process? Why is calling wait() important after forking a child process?



#### Process Management - \$100 Answer

wait() makes the parent process pause execution until one of its child processes terminates. It allows the parent to retrieve the child's exit status and prevents the creation of zombie processes.



#### **Process Management - \$200 Question**

# What are the usual return values of fork() in the parent and child processes upon execution?



#### Process Management - \$200 Answer

The child process returns 0. The parent process returns the pid of the child. On failure, it returns -1.



#### **Process Management - \$300 Question**

How does the exec() family of functions differ from fork(), and why are they often used together?



Process Management - \$300 Answer

fork() creates a new process by duplicating the current one.

exec() replaces the current process with a new program.

They are commonly used together so that the child process created by fork() can immediately load and run a different program using exec().



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#### **Process Management - \$400 Question**

A simple shell spawns a child process to run a foreground job using fork(). However, the user notices that pressing Ctrl+C does not terminate the child process but instead affects the shell itself. What mistake in terminal control is likely causing this issue, and how can it be fixed?



Process Management - \$400 Answer

The shell likely failed to call tcsetpgrp() to assign the child process's group as the terminal's foreground process group. To fix the issue, we would need to run: tcsetpgrp(STDIN\_FILENO, child\_pgid); After forking.



#### **Process Management - \$500 Question**

What are the resulting process groups and process assignments after executing the following code? Assume that the operating system assigns process IDs (PIDs) based on the first available integer greater than 100.



#### Process Management - \$500 Answer





# Signals - \$100 Question

What does the alarm() system call do, and how can it be used to implement a timeout mechanism?



# Signals - \$100 Answer

Alarm schedules a SIGALRM to be delivered to the process after a given number of seconds. To implement a timeout, we call alarm(s)



## Signals - \$200 Question

Describe the difference between signals and interrupts; say a user enters CTRL+C. How is SIGINT actually sent/handled?



# Signals - \$200 Answer Interrupts originate from hardware, whereas signals are higher-level (software). SIGINT is sent by the OS, which checks the foreground process' PCB for the signal disposition (term, ign, ...) or a handler if one is set.



#### Signals - \$300 Question

# Type question here



#### Signals - \$300 Answer

# Type answer here



# Signals - \$400 Question A process sets a signal handler for SIGUSR with sigaction(). Inside the handler, a global variable is modified. What are some potential issues if the handler is preempted by another signal

Click to see answer



## Signals - \$400 Answer

# Race condition: concurrent access to the global variable, inconsistent state of the global var



# Signals - \$500 Question

```
void handler(int sig) {
         printf("Caught SIGINT!\n");
     int main() {
         sigset t block_set, old_set;
11
12
         signal(SIGINT, handler);
13
14
         sigemptyset(&block set);
         sigaddset(&block set, SIGINT);
         sigprocmask(SIG BLOCK, &block set, &old set);
17
18
         sleep(3);
19
20
         sigprocmask(SIG SETMASK, &old set, NULL);
21
22
         while (1) sleep(1);
23
         return 0;
     }
25
```

Explain what this example does; if **CTRL+C** is pressed once every second for 6 seconds, how many times is "Caught SIGINT!" printed?



# Signals - \$500 Answer

Four times. While SIGINT is blocked, the OS does not queue standard signals. So only one of the three SIGINTs during the first three seconds will be handled. In the remaining three seconds, the SIGINTs will be handled correctly.



#### File Descriptors & I/O - \$100 Question

# How does the file descriptor table function in a Unix-like operating system? What is stored in it?



#### File Descriptors & I/O - \$100 Answer

Each process maintains its own file descriptor table, which maps these integers to the underlying open files or devices, including associated metadata like file offsets and access modes.



#### File Descriptors & I/O - \$200 Question

# What roles do the open() and close() system calls play in managing file descriptors?



#### File Descriptors & I/O - \$200 Answer

The open() system call opens a file or device, returning a file descriptor that references it, while close() terminates the association with that file descriptor, freeing the resource for future use.



#### File Descriptors & I/O - \$300 Question

How does file descriptor inheritance work when a process calls fork(), and what impact does it have on I/O redirection?



#### File Descriptors & I/O - \$300 Answer

When a process forks, the child process inherits a copy of the parent's file descriptor table. This inheritance means both processes can access the same open files, which facilitates I/O redirection and inter-process communication.



#### File Descriptors & I/O - \$400 Question

What is the difference between using dup() and dup2() for file descriptor duplication, and how are they typically used in shell redirection?



#### File Descriptors & I/O - \$400 Answer

dup(fd) returns the lowest available new file descriptor referencing the same resource, while dup2(oldfd, newfd) forces duplication into a specific file descriptor number (closing it first if open). Shells use dup2() to redirect standard I/O streams to files or pipes.



#### File Descriptors & I/O - \$500 Question

How do Unix-like shells implement a multi-stage pipeline (cmd1 | cmd2 | cmd3) using file descriptors, and what role do the pipe ends play?



#### File Descriptors & I/O - \$500 Answer

For each pipeline stage, the shell uses pipe() to create a pipe with read and write file descriptors. It then forks each command, connecting the write end of one pipe to the read end of the next using dup2().



## File Systems - \$100 Question

# What is the smallest unit of work for a file system (R/W)? a. A bit b. A byte c. A block d. A page



#### File Systems - \$100 Answer

# C. A block

# Even if you want to change just one byte in a file, you must write an entire block of that file.



#### File Systems - \$200 Question

# In an implicit linked list architecture of the file system...

#### What if I want to grow File D by 2 blocks?

Disk:

Bit- map	Root Dir	File D	free	File B	Also File B	File D Blk 2	File A	File C Blk 2	free	File C	File E
BO	B1	B2	B3	B4	<b>У</b> <sub>В5</sub>	) <sup>B6</sup>	Β7	B8	89	B10	B11



# File Systems - \$200 Answer



- Scan the bitmap to find which blocks are free
- Allocate the blocks and set up pointers to them Disk:

Bit- map	Root Dir	File D	File D Blk 3	File B	Also File B	File D Blk 2	File A	File C Blk 2	File D Blk 4	File C	File E
во	B1	B2	B3	B4	/ <sub>B5</sub>	B6	B7	B8	182	B10	B11



# File Systems - \$300 Question

Block #	File X Block Y		
0	FAT		
1	Root directory	Block #	Next
2	File A Block 1	0	BITMAP/SPECIAL
3	?	1	END
		2	6
4	File B Block 1	3	9
5	Empty	4	END
6	?	5	EMPTY / UNUSED
		6	3
7	File C Block 1	7	END
8	?	8	END
9	?	9	END
10	File D Block 1	10	8

# File Systems - \$300 Answer

Block #	File X Block Y		
0	FAT		
1	Root directory	Block #	Next
2	File A Block 1	0	BITMAP/SPECIAL
3	File A Block 3	1	END
1	Filo B Block 1	2	6
		3	9
5	Empty	4	END
6	File A Block 2	5	EMPTY / UNUSED
		6	3
		7	END
8	<u>File D Block 2</u>	8	END
9	File A Block 4	9	END
10	File D Block 1	10	8

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## File Systems - \$400 Question

With an Inode struct like below, what's the largest file size possible if...
Each block is 512 bytes
Each block\_no\_t is 4 bytes? (Express your answer as an equation)

struct inode st { attributes t metadata; block no t blocks[12]; block no t \*single ind; block no t \*\*double ind; block no t \*\*\*triple ind;



## File Systems - \$400 Answer

- Each block is 512 bytes
- Each block\_no\_t is 4 bytes

```
struct inode_st {
   attributes_t metadata;
   block_no_t blocks[12];
   block_no_t *single_ind;
   block_no_t **double_ind;
   block_no_t ***triple_ind;
};
```

Direct blocks: 12 \* 512 = 6,144 bytes Single indirect: 512/4 = 128 pointers, 128 \* 512 = 65,536 bytes Double indirect: 128 \* 128 \* 512 = 8,388,608 bytes Triple indirect: 128^3 \* 512 = 1,073,741,824 bytes

Total: 6144 + 65536 + 8388608 + 1073741824 ≈ 1.01GB



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## File Systems - \$500 Question

#### Assume... • A block is 1,024 bytes • Each inode is 128 bytes • There are 12 directory entries (dirents) in each inode

<pre>struct dirent {</pre>
<pre>ino_t d_ino;</pre>
<pre>char d_name[]</pre>
uint8 t d type

#### maximum

What is the minimum number of directory entries (struct dirents) that must be traversed to resolve the path /dir\_one/dir\_two/dir\_three/file.txt and access the data blocks of file.txt in the best-case scenario? Assume you have not accessed the root inode.

(assume file.txt does exist in this path)

Click to see answer

/\* File inode number \*/

/\* Null terminated name of file \*/

/\* Indicator of file type ONLY IN ext2, 3, 4 \*/



# File Systems - \$500 Answer 5 inodes 48 dirents



#### **Concurrency & Synchronization - \$100 Question**

What is the unit of scheduling in most modern OS? a. A process b. A thread c. A Procedure



#### Concurrency & Synchronization - \$100 Answer

# b. A thread



#### Concurrency & Synchronization - \$200 Question

With a 4-core machine, what's the most amount of performance boost I can get when executing a task with multiple threads? Using parallelism or concurrency?



Concurrency & Synchronization - \$200 Answer

#### 4x with parallelism

Physical parallelism is limited by the number of available cores. With 4 cores, only 4 threads can truly execute instructions simultaneously.



Concurrency & Synchronization - \$300 Question

Which of the following are shared (1) between processes (2) between threads but not processes (3) by neither of the two? a. Stack b. Heap c. Registers d. Stack pointer e. Program counter f. Pipes via pipe()

#### Concurrency & Synchronization - \$300 Answer

- a. Stack
- b. Heap
- c. Registers
- d. Stack pointer
- e. Program counter
  - f. Pipes via pipe()

# (1). Between processes: f(2) Between threads but not processes: b(3) Neither: acde



**Concurrency & Synchronization - \$400 Question** 

Describe a scenario where using processes would be preferable to threads despite their higher overhead.



#### Concurrency & Synchronization - \$400 Answer

# One advantage for using processes is the isolation: if one process crashes, others can keep working.

Scenario: web server handling multiple client requests with separate processes.



#### Concurrency & Synchronization - \$500 Question

In a multithreaded program where one thread attempts to increment a global counter variable 1000 times and another thread attempts to decrement it 1000 times.

Why might the final value not be 0?



#### Concurrency & Synchronization - \$500 Answer

#### Race condition!

An example scenario:

• Thread A reads the counter value as 10

- Thread B reads the counter value as 10
   Thread A adds 1 and writes back 11
- Thread B subtracts 1 and writes back 9



#### Scheduling - \$100 Question

# Which scheduling algorithm(s) runs the risk of "starving" its jobs?



#### Scheduling - \$100 Answer

# Shortest Job First (SJF) Priority Round Robin



#### Scheduling - \$200 Question

If you're guaranteed that all incoming jobs take approximately the same amount of time to complete, which scheduling algorithm makes the most sense to implement and why?



## Scheduling - \$200 Answer

First Come First Served (FCFS):

- Easiest to implement
- Consistent duration of jobs negates any positive benefits of other scheduling algorithms.

 Any preemptive scheduling algorithms will cost more time with the extra context-switching



## Scheduling - \$300 Question

Given the following jobs and their estimated lengths, what time quantum would you choose for a Round Robin Scheduler?

Job A: 2-10ns Job B: 4-6ns Job C: 1-17ns Job D: 6-7ns Job E: 8-13ns Job F: 1-2ns Job G: 3-14ns Job H: 15-18ns Job I: 5-11ns Job J: 6-9ns



#### Scheduling - \$300 Answer

If you want 80-90% of jobs to finish within 1 attempt: need to use upper estimate

80% of jobs finish within 14ns 90% of jobs finish within 17ns So 14ns <= TQ <= 17n



## Scheduling - \$400 Question

Rank the following in order of preference in priority round-robin and explain your reasoning:

- 1. Involving little to no I/O
- 2. Involving a lot of I/O, but split into smaller chunks
- 3. One long sequence of I/O accesses at the first half of the job, but no I/O in the last half

Ties are allowed!



#### Scheduling - \$400 Answer

# Small I/O bursts >= One long sequence of I/O > No I/O

Want to make most utilization of I/O devices when needed scheduler can context-switch to another thread that needs CPU while I/O is accessed, and switch back quickly if I/O access is relatively short. If no I/O bursts are needed, then nothing useful is happening while the thread is in queue!



# Scheduling - \$500 Question

Job Name	Arrival Time	Running Time
Swellow	2	11
Pidove	0	8
Fletchling	12	20
Starly	7	15
Doduo	10	4

Using a Round Robin scheduling algorithm and a time quantum of 8...

- 1. what is the finishing time for each job?
- 2. What is the average waiting time?
- 3. Is 8 the best time quantum to use given the current info?



# Scheduling - \$500 Answer

Job Name	Finishing Time	Wait Time
Swellow	39	39 - 2 - 11 = 26
Pidove	8	8 - 0 - 8 = 0
Fletchling	58	58 - 12 - 20 = 26
Starly	46	46 - 7 - 15 = 24
Doduo	28	28 - 10 - 4 = 14

#### Average wait time: 18

No, this time quantum is quite short - only 40% of known jobs can finish within 1 attempt. A better time quantum would be between 15 and 19, inclusive.



# Other questions you can add about scheduling:

Q: what scheduling algorithm can be implemented using an array of linked lists? A: priority round robin (~200 pts)

# FINAL E PARJU

# Topic: Type topic here



#### **Final Jeopardy Question**

# Type question here



#### Final Jeopardy Answer

# Type answer here

