

Useful Properties:

$$W_N^{k(N-n)} = W_N^{-kn} = (W_N^{kn})^* \quad W_N^{kn} = W_N^{k(n+N)} = W_N^{(k+N)n} \quad W_N^2 = W_{N/2}$$

Example 1:

A long *periodic* sequence x of period $N = 2^r$ (r is an integer) is to be convolved with a finite-length sequence h of length K .

- Show that the output y of this convolution (filtering) is *periodic*. What is its period?
- Let $K = mN$ where m is an integer; N is large. How would you implement this convolution *efficiently*? Explain your analysis clearly. Compare the *total number of multiplications* required in your scheme to that in the direct implementation of FIR filtering. (Consider the case $r = 10$, $m = 10$).

Example 2:

A sequence $x = \{x[n], n = 0, 1, \dots, N-1\}$ is given; let $X(e^{j\omega})$ be its DTFT.

- Suppose $N=10$. You want to evaluate both $X(e^{j2\pi 7/12})$ and $X(e^{j2\pi 3/8})$. The only computation you can perform is one DFT, on any one input sequence of your choice. Can you find the desired DTFT values? (*Show your analysis and explain clearly.*)

- Suppose N is large. You want to obtain $X(e^{j\omega})$ at the following $2M$ frequencies:

$$\omega = \frac{2\pi}{M}m, m = 0, 1, \dots, M-1 \quad \text{and} \quad \omega = \frac{2\pi}{M}m + \frac{2\pi}{N}m, m = 0, 1, \dots, M-1$$

Here $M = 2^\mu \ll N = 2^\nu$

A standard radix-2 FFT algorithm is available. You may execute the FFT algorithm *once or more than once*, and *multiplications* and *additions* outside of the FFT are *allowed*, if necessary.

- You want to get the $2M$ DTFT values with as few *total multiplications* as possible (*including those in the FFT*). Give explicitly the best method you can find for this, with an estimate of the *total number of multiplications* needed in terms of M and N .
- Does your result change if extra multiplications outside of FFTs are *not* allowed?