

University of Pennsylvania
Department of Electrical and System Engineering
System-on-a-Chip Architecture

ESE532, Fall 2019

Final

Thursday, December 19

- Exam ends at 11:00AM; begin as instructed (target 9:00AM).
Do not open exam until instructed.
- Problems weighted as shown.
- Calculators allowed.
- Closed book = No text or notes allowed.
- Show work for partial credit consideration.
- Unless otherwise noted, answers to two significant figures are sufficient.
- Sign Code of Academic Integrity statement (see last page for code).

I certify that I have complied with the University of Pennsylvania's Code of Academic Integrity in completing this exam.

Name: Solution

1	2	3	4	5	6	7	8a	8b	8c	8d	Total
10	10	10	10	10	10	10	10	2	8	10	100

[Average 59, Std. Dev. 14](#)

Consider the following code to render augmented reality features on a real-time video stream

```

code_one.c          Fri Dec 20 10:14:13 2019          1

int WIDTH 4096
int HEIGHT 2048
int COLORS 3
int MASK 3

int VPARAMS 5
int VP_X 0
int VP_Y 1
int VP_XS 2
int VP_YS 3
int VP_ROT 4

int XOFF 1
int YOFF 1
int ROT 1
int XSCALE 2
int XFACT 2 // typo for XSFACT in original
int XSFACT 2
int YSCALE 2
int YSFAC 2 // typo for YSFAC in original
int YFACT 2

uint16_t reference[HEIGHT][WIDTH][COLORS];
uint16_t overlay[HEIGHT][WIDTH][COLORS+1]; // +1 for mask
int16_t sintable[360]; // -1 to 1 -- scaled by 2^14
int16_t costable[360];

void main() {
    while (true) { // loop Z
        augment_frame();
    }
}

void augment_frame() {
    uint16_t raw[HEIGHT][WIDTH][COLORS]; // uint16_t for 16b (2 byte) color per pixel
    uint16_t augment[HEIGHT][WIDTH][COLORS];
    uint16_t augmented[HEIGHT][WIDTH][COLORS];
    uint16_t old_viewpoint[VPARAMS];
    uint16_t viewpoint[VPARAMS];
    uint16_t *tmp_viewpoint;
    get_image(raw);
    tmp_viewpoint=old_viewpoint;
    old_viewpoint=viewpoint;
    viewpoint=tmp_viewpoint;
    compute_viewpoint(raw,reference,old_viewpoint,viewpoint);
    render_augmentation(viewpoint,overlay,augment);
    merge_frames(reference,viewpoint,raw,augment,augmented);
    send_image(augmented);
}

```

```

code_two.c      Mon Dec 23 18:18:34 2019      1

void compute_viewpoint(uint16_t ***image, uint16_t ***reference,
                       int16_t *old, int16_t *current)
{
    uint64_t best_score=MAXINT; // maximum representable integer

    for (int rot=old[VP_ROT]-ROT;rot<old[VP_ROT]+ROT;rot+=1) { // loop A
        int16_t sr=sintable[rot]; // result is a fraction
        int16_t cr=costable[rot];
        for (int x=old[VP_X]-XOFF;x<old[VP_X]+XOFF;x++) // loop B
            for (int y=old[VP_Y]-YOFF;y<old[VP_Y]+YOFF;y++) // loop C
                for (int xs=old[VP_XS]/XSCALE;xs<old[VP_XS]*XSCALE;xs*=XSFACT) // loop D
                    for (int ys=old[VP_YS]/YSCALE;ys<old[VP_YS]*YSCALE;ys*=YSFACT) // loop E
                        {
                            uint64_t score=0;
                            for (int iy=0;iy<HEIGHT;iy++) // loop F
                                for (int ix=0;ix<WIDTH;ix++) // loop G
                                    {
                                        uint16_t tx=((ix*cr+iy*sr)*xs)>>(14+8)+x; // 14 to scale sr, cr
                                        uint16_t ty=((ix*sr+iy*cr)*ys)>>(14+8)+y; // +8 for xscale, yscale
                                        if ((tx>=0) && (tx<WIDTH) && (ty>=0) && (ty<HEIGHT))
                                            for (int c=0;c<COLORS;c++) // loop H
                                                score+=abs(image[iy][ix][c]-reference[ty][tx][c]);
                                    }
                            if (score<best_score)
                                {
                                    best_score=score;
                                    current[VP_ROT]=rot;
                                    current[VP_X]=x;
                                    current[VP_Y]=y;
                                    current[VP_XS]=xs;
                                    current[VP_YS]=ys;
                                }
                        }
    }
}

void render_augmentation(int16_t *current, uint16_t ***overlay, uint16_t ***image)
{
    uint16_t rot=current[VP_ROT];
    uint16_t x=current[VP_X];
    uint16_t y=current[VP_Y];
    uint16_t xs=current[VP_XS];
    uint16_t ys=current[VP_YS];
    int16_t sr=sintable[rot]; // result is a fraction
    int16_t cr=costable[rot];
    for (int iy=0;iy<HEIGHT;iy++) // loop I
        for (int ix=0;ix<WIDTH;ix++) // loop J
            image[iy][ix]=UNMAPPED; // assume this runs like streaming data copy
    for (int iy=0;iy<HEIGHT;iy++) // loop K
        for (int ix=0;ix<WIDTH;ix++) // loop L
            {
                uint16_t tx=((ix*cr+iy*sr)*xs)>>(14+8)+x; // 14 to scale sr, cr
                uint16_t ty=((ix*sr+iy*cr)*ys)>>(14+8)+y; // +8 for xscale, yscale
                if ((tx>=0) && (tx<WIDTH) && (ty>=0) && (ty<HEIGHT)
                    && (overlay[ty][tx][MASK]>0))
                    for (int c=0;c<COLORS;c++) // loop M
                        image[iy][ix][c]=overlay[ty][tx][c];
            }
}

```

```

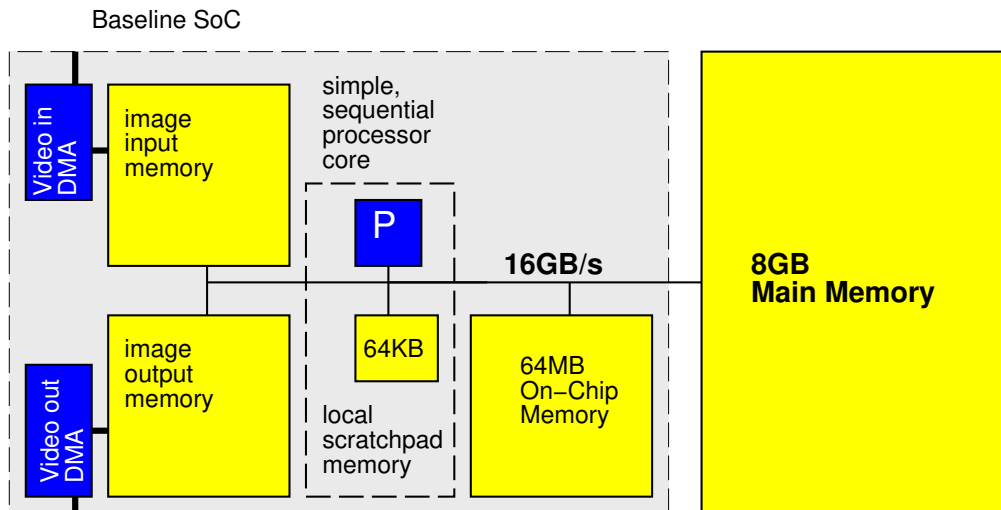
code_three.c      Thu Dec 19 05:26:56 2019      1
void merge_frames(uint16_t ***reference, int16_t *current,
                  uint16_t ***image, uint16_t ***augment, uint16_t ***augmented)
{
    uint16_t rot=current[VP_ROT];
    uint16_t x=current[VP_X];
    uint16_t y=current[VP_Y];
    uint16_t xs=current[VP_XS];
    uint16_t ys=current[VP_YS];
    int16_t sr=sintable[rot]; // result is a fraction
    int16_t cr=costable[rot];
    for (int iy=0;i<HEIGHT;iy++) // loop N
        for (int ix=0;i<WIDTH;ix++) // loop O
            {
                uint16_t tx=((ix*cr+iy*sr)*xs)>>(14+8)+x; // 14 to scale sr, cr
                uint16_t ty=((ix*sr+iy*cr)*ys)>>(14+8)+y; // +8 for xscale, yscale
                if ((tx>=0) && (tx<WIDTH) && (ty>=0) && (ty<HEIGHT)
                    && (augment[iy][ix]!=UNMAPPED))
                    {
                        uint32_t diff=0;
                        for (int c=0;c<COLORS;c++) // loop P
                            diff+=abs(image[iy][ix][c]-reference[ty][tx][c]);
                        if (diff<THRESH)
                            for (int c=0;c<COLORS;c++) augmented[iy][ix][c]=augment[iy][ix][c];
                        else
                            for (int c=0;c<COLORS;c++) augmented[iy][ix][c]=image[iy][ix][c];
                    }
            }
}

void get_image(uint16_t ***image)
{
    for (int iy=0;i<HEIGHT;iy++)
        for (int ix=0;i<WIDTH;ix++)
            for (int c=0;c<COLORS;c++)
                image[iy][ix][c]=image_in[iy][ix][c];
}

void send_image(uint16_t ***image)
{
    for (int iy=0;i<HEIGHT;iy++)
        for (int ix=0;i<WIDTH;ix++)
            for (int c=0;c<COLORS;c++)
                image_out[iy][ix][c]=image[iy][ix][c];
}

```

We start with a baseline, single processor system as shown.



- For simplicity throughout, we will treat non-memory indexing adds (subtracts count as adds), compares, abs, shifts, and multiplies as the only compute operations. We'll assume the other operations take negligible time or can be run in parallel (ILP) with the adds, abs, shift, multiplies, and memory operations. (Some consequences: You may ignore loop and conditional overheads in processor runtime estimates; you may ignore computations in array indices.)
- Baseline (simple, sequential) processor can execute one multiply, compare, shift, abs, or add per cycle and runs at 1 GHz.
- Data can be transferred between pairs of memory (including main memory) at 16 GB/s when streamed in chunks of at least 2048B. Assume for loops that only copy data can be auto converted into streaming operations.
- Non-streamed access to the main memory takes 100 cycles and can move 8B.
- Non-streamed access to image and 64 MB on-chip memories takes 10 cycles and can move 8B.
- Baseline processor has a local scratchpad memory that holds 64KB of data. Data can be streamed into the local scratchpad memory at 16 GB/s. Non-streamed accesses to the local scratchpad memory take 1 cycle.
- Baseline processor is 1 mm² of silicon including its 64KB local scratchpad.
- By default, all arrays live in the 8 GB main memory.
- `image_in` and `image_out` live in the respective image input and image output memories.
- Arrays for `sintable`, `costable` and viewpoints (`old_viewpoint`, `viewpoint`) live in local scratchpad memory.
- Assume scalar (non-array) variables can live in registers.
- Assume all additions are associative.
- Assume comparisons, adds, and multiplies take 1 ns when implemented in hardware accelerator, so fully pipelined accelerators also run at 1 GHz. A compare-mux operation can also be implemented in 1 ns. Consider abs and shift free in hardware.
- Data can be transferred to accelerator local memory at the same 16 GB/s when streamed in chunks of at least 2048B.

1. Simple, Single Processor Resource Bounds

- (a) Based only on the resource bound for compute operations, what throughput can a simple, single processor system achieve [answer in frames/second, or equivalently, `augment_frame` calls per second]?

<code>get_image</code>	all DMA	0
<code>compute_viewpoint</code>	$2^5 \times 4096 \times 2048 \times (12 + 3 \times 3)$ cycles	5.6 s
<code>render_augmentation</code>	$4096 \times 2048 \times 12$ cycles	0.10 s
<code>merge_frames</code>	$4096 \times 2048 \times (12 + 3 \times 3)$ cycles	0.18 s
<code>send_image</code>	all DMA	0
Total		5.9 s

0.17 frames/second

- (b) Based only on the resource bound for memory operations, what throughput can a simple, single processor system achieve [answer in frames/second]?

<code>get_image</code>	$\frac{4096 \times 2048 \times 3 \times 2}{16 \times 10^9}$	0.0031 s
<code>compute_viewpoint</code>	$2^5 \times 4096 \times 2048 \times (2 \times 100)$ cycles image[iy][ix] and reference[ty][tx] is single read ignore small terms sin/costable, current update	54 s
<code>render_augmentation</code>	$\frac{4096 \times 2048 \times 3 \times 2}{16 \times 10^9} +$ $4096 \times 2048 \times 2 \times 100$ cycles overlay[ty][tx] including mask is single read image[iy][ix] is single write	1.7s
<code>merge_frames</code>	$4096 \times 2048 \times 4 \times 100$ cycles	3.4 s
<code>send_image</code>	$\frac{4096 \times 2048 \times 3 \times 2}{16 \times 10^9}$	0.0031
Total		59 s

0.017 frames/second

2. Based on the simple, single processor mapping from Question 1:

(a) What function is the bottleneck? (circle one)

get_image
(compute_viewpoint)
render_augmentation
merge_frames
send_image

(b) What is the Amdahl's Law speedup if you only accelerate the identified function?

$$\frac{64.9}{5.4} = 11$$

3. Data Parallel and Reduce: Classify Loops

Loop	Data Parallel?	Associative Reduce?	Must be Sequential?
A		X	
F		X	
K	X		
N	X		
Z (main)			X

Z must compute new viewpoint from one iteration/image before starting computation on next image.

A is a min-reduce on best_score.

F is a sum-reduce for score.

Computation for image[iy][ix] (K) and augmented[iy][ix] (N) are each independent of other elements of the respective arrays.

4. Data Streaming:

- (a) Can the producer and consumer operate concurrently on the same input image? or must the consumer work on a different (earlier) input image? (“Same Image?” column)
- (b) How big (minimum size) does the buffer (or other data storage space) need to be between the identified loops in order to allow the loops to profitably execute concurrently?

(Hint: Based on data dependencies, under what scenarios and granularity can the identified loops act as a producer-consumer pair in a pipeline.)

Loop Pair	(a) Same Image?	(b) Size (bytes)
<code>get_image</code> → <code>compute_viewpoint</code>	N	48 MB
<code>compute_viewpoint</code> → <code>render_augmentation</code>	N	10 B
<code>render_augmentation</code> → <code>merge_frames</code>	Y	6 B
<code>merge_frames</code> → <code>send_image</code>	Y	6 B

Explain size choices for partial credit consideration.

Must hold onto an entire image from `get_image` to perform the search in `compute_viewpoint`.

Need to process entire search in `compute_viewpoint` before have a new viewpoint ($5 \times 2B = 10B$) to pass to `render_augmentation`. `render_augmentation` needs the viewpoint to process any image pixels.

As `render_augmentation` completes a pixel ($3 \times 2B = 6B$), it is ready to use, in the same order, in `merge_frames`.

As `merge_frames` completes a pixel ($3 \times 2B = 6B$), it is ready to be sent by `send_image` in the same order produced.

5. Latency Bound

- (a) What is the critical path (latency bound) for the entire computation as captured in the `augment_frame` function?

compute_ viewpoint	read sintable, costable	1
	multiply by sine, cos	1
	add sin/cos terms	1
	scale	1
	(shifts for free)	0
	add offset	1
	read image and reference	100
	subtract	1
	(abs for free)	0
	sum reduce	$\log_2(4096 \times 2048 \times 3) = 25$
	min reduce	$\log_2(2^5) = 5$
render_ augmentation	read sintable, costable	1
	multiply by sine, cos	1
	add sin/cos terms	1
	scale	1
	(shifts for free)	0
	add offset	1
	read overlay	100
	(don't write image, just use it below)	0
merge_ frames	compute tx, ty with above	0
	(reads, if needed, happen with overlay above)	0
	subtract	1
	(abs for free)	0
	sum reduce	$\log_2(3) = 2$
	(don't write image, just use for output)	0
Total		244

- (b) What is the latency bound Iteration Interval (II) for the `main` computation?
(Hint: builds on part (a).)

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Only need to compute new viewpoint.

6. Consider rewriting the body of `compute_viewpoint` to minimize the memory resource bound by exploiting the scratchpad memory and the 64MB on-chip memory and streaming data transfers.

- (a) Identify new temporary arrays allocated to scratchpad memory or 64MB on-chip memory (and specify which memory each new array is in).

```
uint16_t image_line[WIDTH][COLORS]; // scratchpad
uint16_t ref_copy[HEIGHT][WIDTH][COLORS]; // in 64MB on-chip memory
```

- (b) Describe how you use these arrays.

Copy reference image into 64MB on-chip memory at beginning of function and operate on it from there.

Copy each line ($4096 \times 3 \times 2B$) into `image_line` in the body of F before starting G. All references to `image[iy][ix]` now go to `image_line`.

Common Problem: `reference` is accessed randomly. A line buffer will not work for it.

- (c) Account for total memory usage in the local scratchpad.

24KB in `image_line`; 1440B in `sintable` and `costable`; 20B in `old` and `current`. Less than 26KB

- (d) Estimate the new memory resource bound for your optimized `compute_viewpoint`.

$$\frac{4096 \times 2048 \times 3 \times 2}{16 \times 10^9} + 2^5 \times 2048 \times \frac{4096 \times 3 \times 2}{16 \times 10^9} + \frac{2^5 \times 4096 \times 2048 \times (10+1)}{10^9}$$

3.1 seconds

(This page intentionally left mostly blank for answers.)

7. Consider a multiprocessor design that included N copies of a vector processor with 16 vector lanes, each operating on 16b data. This is a single-issue vector processor that can either issue one vector or one scalar operation on each cycle. Assume the loops you identified as data parallel or reduce operation in Question 3 are perfectly vectorizable. Each vector processor requires 2 mm^2 including a local 64KB data scratchpad.

- (a) Based on computational requirements alone, how many vector processors do you need to achieve a 30 frame per second frame rate? [for this problem, ignore memory and communication]

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- (b) Identify how the processors are used.

Everything that takes significant time is data parallel or an associative reduce.

<code>compute_viewpoint</code>	81
<code>render_augmentation</code>	1
<code>merge_frames</code>	(shared)

(This page intentionally left mostly blank for pagination and answers.)

8. Considering a custom hardware accelerator implementation for `compute_viewpoint` where you are designing both the compute operators and the associated memory architecture. How would you use loop unrolling and array partitioning to achieve guaranteed throughput of 30 frames per second of throughput while minimizing area? Use the following area model in units of mm^2 :

- n -bit counters: $n \times 10^{-5}$
- n -bit adder: $n \times 10^{-5}$
- 16×16 multiplier: 2.5×10^{-3}
- p -port, w -bit wide memory holding d words: $w(1+p)(d+6) \times 10^{-7}$

Make the (probably unreasonable) assumption that reads from these memories can be completed in one cycle.

Start by assuming we unroll H; we need to understand how much unrolling of the rest of the loops is required. Since the loops are associative reduce, the inner loop can be pipelined to $\Pi=1$. $\frac{2^5 \times 4096 \times 2048}{A \times 10^9} \leq \frac{1}{30}$, giving us A a little over 8. This suggests unrolling about a factor of 16 beyond H will be sufficient.

Common Problem: Not accounting for the operations that can be pipelined.

- (a) Unrolling for each loop?

Loop	Unroll Factor
A	1
B	1
C	1
D	1
E	1
F	1
G	16
H	3

- (b) For the unrolling, how many multipliers and adders?

Multipliers	$6 \times 16 = 96$
Adders	$16 \times (4 + 3 \times 2) = 160$

(note: for upcoming area calculation, you will need to break down adders by size.)

64b: $3 \times 16 = 48$, 16b: $(3 + 4) \times 7 = 49$

(c) Array partitioning for each array?

Note: blank rows left for local arrays you may have added when optimizing memory in Question 6.

Array	Array Partition	Ports	Width	Depth/partition
old[]	none	1	16	10
current[]	none	1	16	10
sintable[]	none	1	16	360
costable[]	none	1	16	360
image[]	n/a			
reference[]	n/a			
image_line[]	cyclic 16 dim 1, x complete dim 2 (and pack), c	1	48	256
ref_tmp[]	none	16	48	8,388,608

Common Problem: **reference** needs ports rather than partitioning since it is accessed randomly.

(d) Estimate the area for the accelerator.

Resource	Count	Area/resource	Area
16×16 multipliers	96	2.5×10^{-3}	0.24
64b-adders	48	64×10^{-5}	0.03072
16b-adders	112	16×10^{-5}	0.01792
3b-counters	5	3×10^{-5}	15×10^{-5}
8b-counters	1	8×10^{-5}	8×10^{-5}
12b-counters	1	12×10^{-5}	12×10^{-5}
memory (1,16,10)	2	5.1×10^{-5}	1.0×10^{-4}
memory (1,16,360)	2	1.2×10^{-3}	2.4×10^{-3}
memory (1,48,256)	16	2.5×10^{-3}	0.0402
memory(16,48,8388608)	1	680	680
Total			680

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