

ESE535: Electronic Design Automation

Day 8: February 11, 2009
Dataflow



Penn ESE535 Spring 2009 -- DeHon

Previously

- Behavioral model (C)
 - Based on sequential semantics
 - Control flow semantics
- Scheduling of concurrent operations

2

Control Flow Model

- Can extract some parallelism
 - “Instruction”-level parallelism within basic block (hyperblock)
- Must sequentialize blocks
 - Only one block run at a time
 - Limited Parallelism

3

Want to see

- There are other compute models
 - More natural for parallelism/hardware

4

Today

- Dataflow
- SDF
 - Single rate
 - Multirate
- Dynamic Dataflow
- Expression

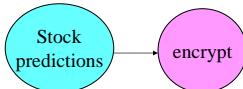
5

Parallelism Motivation

Penn ESE535 Spring 2009 -- DeHon

6

Producer-Consumer Parallelism

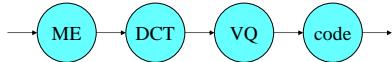


- Can run concurrently
- Just let consumer know when producer sending data
- Not described well in C
 - How put encryption tasks in same basic-block as prediction code?
 - Would prefer to develop separately
 - Tasks run own threads of control
 - Synchronize on data transfer

7

Penn ESE535 Spring 2009 -- DeHon

Pipeline Parallelism



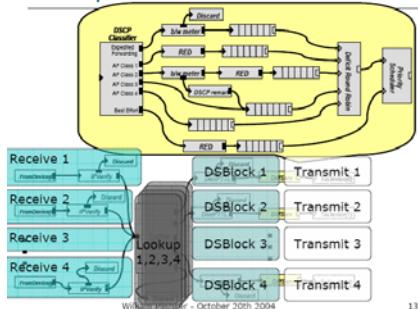
- Can potentially all run in parallel
- Like physical pipeline
- Useful to think about **stream** of data between operators

8

Penn ESE535 Spring 2009 -- DeHon

Plisker Task Example

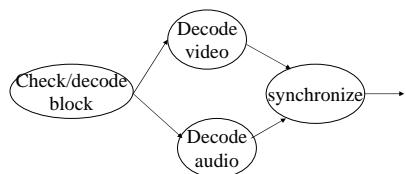
Example: 4 Port DiffServ



13 9

Penn ESE535 Spring 2009 -- DeHon

DAG Parallelism

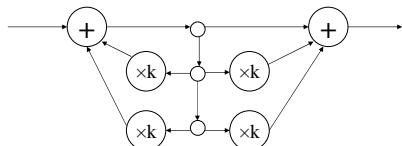


- Doesn't need to be linear pipeline
- Synchronize inputs

10

Penn ESE535 Spring 2009 -- DeHon

Graphs with Feedback



- In general may hold state
- Very natural for many tasks

11

Penn ESE535 Spring 2009 -- DeHon

Definitions

12

Penn ESE535 Spring 2009 -- DeHon

Dataflow / Control Flow

Dataflow

- Program is a graph of operators
- Operator consumes **tokens** and produces tokens
- All operators run concurrently

Control flow

- Program is a sequence of operations
- Operator reads inputs and writes outputs into common store
- One operator runs at a time
 - Defines successor

13

Penn ESE535 Spring 2009 -- DeHon

Token

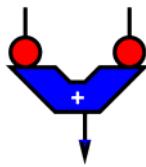
- Data value with presence indication
 - May be conceptual
 - Only exist in high-level model
 - Not kept around at runtime
 - Or may be physically represented
 - One bit represents presence/absence of data
 - Data comes in packets

14

Penn ESE535 Spring 2009 -- DeHon

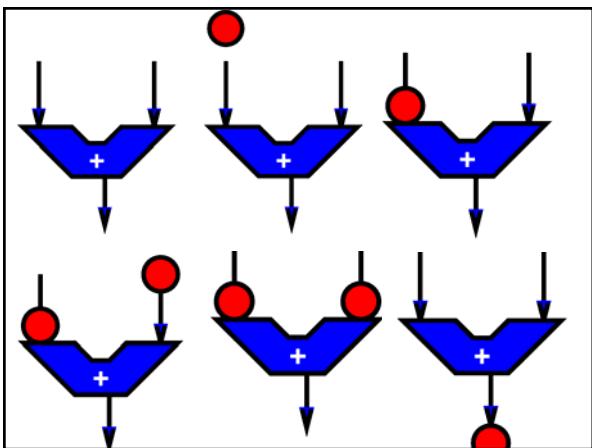
Operator

- Takes in one or more inputs
- Computes on the inputs
- Produces a result
- Logically self-timed
 - “Fires” only when input set present
 - Signals availability of output



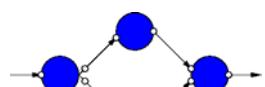
15

Penn ESE535 Spring 2009 -- DeHon



Dataflow Graph

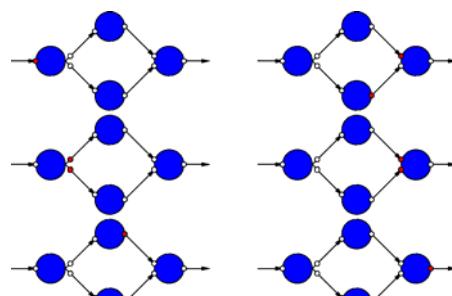
- Represents
 - computation sub-blocks
 - linkage
- Abstractly
 - controlled by data presence



17

Penn ESE535 Spring 2009 -- DeHon

Dataflow Graph Example



18

Penn ESE535

Stream

- Logical abstraction of a persistent point-to-point communication link
 - Has a (single) source and sink
 - Carries data presence / flow control
 - Provides in-order (FIFO) delivery of data from source to sink

19

Penn ESE535 Spring 2009 -- DeHon

Streams

- Captures communications structure
 - Explicit producer → consumer link up
- Abstract communications
 - Physical resources or implementation
 - Delay from source to sink

20

Penn ESE535 Spring 2009 -- DeHon

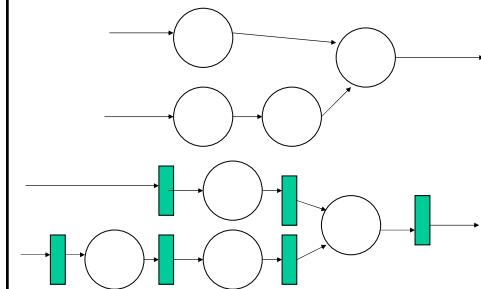
Dataflow Abstracts Timing

- Doesn't say
 - on which cycle calculation occurs
- Does say
 - What order operations occur in
 - How data interacts
 - i.e. which inputs get mixed together
- Permits
 - Scheduling on different # of resources
 - Operators with variable delay
 - Variable delay in interconnect

21

Penn ESE535 Spring 2009 -- DeHon

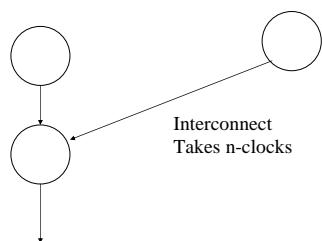
Difference: Dataflow Graph/Pipeline



22

Penn ESE535 Spring 2009 -- DeHon

Clock Independent Semantics



23

Penn ESE535 Spring 2009 -- DeHon

Semantics

- Need to implement semantics
 - i.e. get same result as if computed as indicated
- But can implement any way we want
 - That preserves the semantics

24

Penn ESE535 Spring 2009 -- DeHon

Dataflow Variants

Penn ESE535 Spring 2009 – DeHon

25

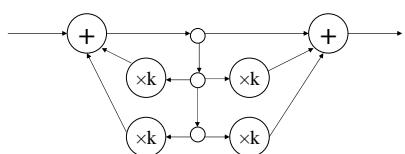
Synchronous Dataflow

- Particular, restricted form of dataflow
- Each operator
 - Consumes a fixed number of input tokens
 - Produces a fixed number of output tokens
 - When full set of inputs are available
 - Can produce output
 - Can fire any (all) operators with inputs available at any point in time

Penn ESE535 Spring 2009 – DeHon

26

Synchronous Dataflow



27

SDF: Execution Semantics

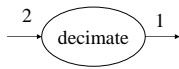
```
while (true)
    Pick up any operator
    If operator has full set of inputs
        Compute operator
        Produce outputs
        Send outputs to consumers
```

Penn ESE535 Spring 2009 – DeHon

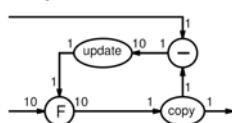
28

Multirate Synchronous Dataflow

- Rates can be different
 - Allow lower frequency operations
 - Communicates rates to CAD
 - Something not clear in RTL
 - Use in scheduling, provisioning
 - Rates must be constant



Penn ESE535 Spring 2009 – DeHon



29

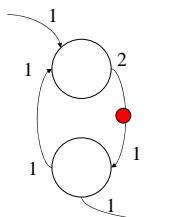
SDF

- Can validate flows to check legal
 - Like KCL → token flow must be conserved
 - No node should
 - be starved of tokens
 - Collect tokens
- Schedule onto processing elements
 - Provisioning of operators
- Provide real-time guarantees
- Simulink is SDF model

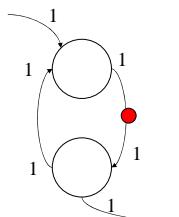
Penn ESE535 Spring 2009 – DeHon

30

SDF: good/bad graphs

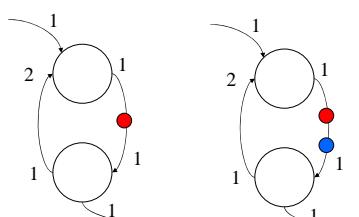


Penn ESE535 Spring 2009 -- DeHon



31

SDF: good/bad graphs



Penn ESE535 Spring 2009 -- DeHon

32

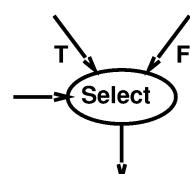
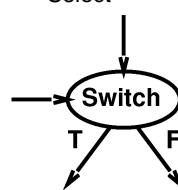
Dynamic Rates?

- When might static rates be limiting?
 - Compress/decompress
 - Lossless
 - Even Run-Length-Encoding
 - Filtering
 - Discard all packets from gerald0
 - Anything data dependent

33

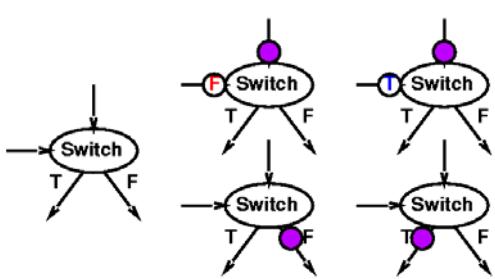
Data Dependence

- Add Two Operators
 - Switch
 - Select



Penn ESE535 Spring 2009 -- DeHon

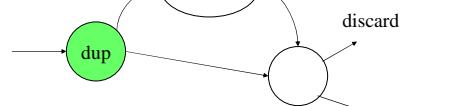
Switch



35

Penn ESE535 Spring 2009 -- DeHon

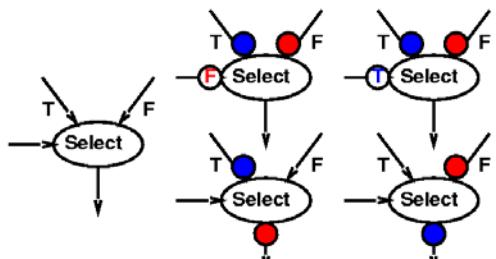
Filtering Example



36

Penn ESE535 Spring 2009 -- DeHon

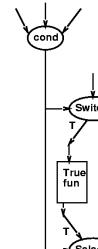
Select



Penn ESE535 Spring 2009 -- DeHon

37

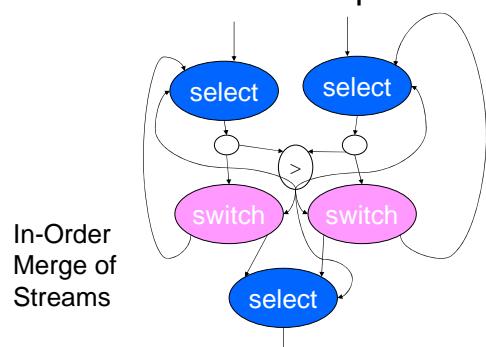
Constructing If-Then-Else



Penn ESE535 Spring 2009 -- DeHon

38

Select Example

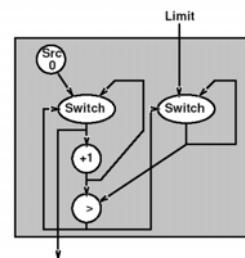


Penn ESE535 Spring 2009 -- DeHon

39

Looping

- For ($i=0; i < \text{Limit}; i++$)



Penn ESE535 Spring 2009 -- DeHon

i

Dynamic Challenges

- In general, cannot say
 - If a graph is well formed
 - Will not deadlock
 - How many tokens may have to buffer in stream
 - Right proportion of operators for computation

Penn ESE535 Spring 2009 -- DeHon

41

Expression

Penn ESE535 Spring 2009 -- DeHon

42

Expression

- Could express operators in C/Java
 - Each is own thread
- Link together with Streams
- E.g. SystemC

Penn ESE535 Spring 2009 -- DeHon

43

C Example

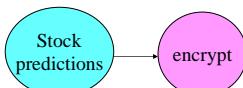
```
while (!(eos(stream_a) && !(eos(stream_b)))
    A=stream_a.read();
    B=stream_b.read();
    Out=(a+b)*(a-b);
    stream_out.write(Out);
```

Penn ESE535 Spring 2009 -- DeHon

44

Connecting up Dataflow

```
stream stream1=new stream();
operator prod=new stock(stream1);
operator cons=new encrypt(stream1);
```



Penn ESE535 Spring 2009 -- DeHon

45

Summary

- Dataflow Models
 - Simple pipelines
 - DAGs
 - SDF (single, multi)-rate
 - Dynamic Dataflow
- Allow
 - express parallelism
 - freedom of implementation

Penn ESE535 Spring 2009 -- DeHon

46

Admin

- Reading
 - Monday online
- Assignment 2A due Monday

Penn ESE535 Spring 2009 -- DeHon

47

Big Ideas:

- Dataflow
 - Natural model for capturing computations
 - Communicates useful information for optimization
 - Linkage, operator usage rates
- Abstract representations
 - Leave freedom to implementation

Penn ESE535 Spring 2009 -- DeHon

48