

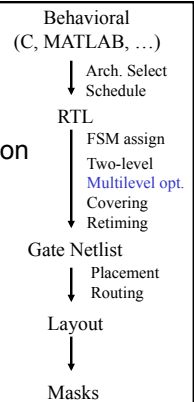
ESE535: Electronic Design Automation

Day 22: April 15, 2015
Multi-level Synthesis



Today

- Multilevel Synthesis/Optimization
 - Why
 - Transforms -- defined
 - Division/extraction
 - How we support transforms



Multi-level Logic

- General circuit netlist
- May have
 - sums within products
 - products within sum
 - arbitrarily deep
- $y = ((a(b+c)+e)fg+h)i$

Why Multi-level Logic?

- $ab(c+d+e)(f+g)$
- $abcf+abdf+abef+abcf+abdg+abeg$
- 6 product terms \rightarrow 23 2-input gates
- vs. 3 gates: and4,or3,or2 \rightarrow 6 2-input gates
- Aside from Pterm sharing between outputs,
 - two level cannot share sub-expressions

Why Multi-level Logic

- $a \text{ xor } b$
 - $a/b+/ab$
- $a \text{ xor } b \text{ xor } c$
 - $a/bc+/abc+/a/b/c+ab/c$
- $a \text{ xor } b \text{ xor } c \text{ xor } d$
 - $a/bcd+/abcd+/a/b/cd+ab/cd+/ab/c/d+a/b/c/d+abc/d+/a/bc/d$

Why Multilevel

- $a \text{ xor } b$
 - $a/b+/ab$
 - $a \text{ xor } b \text{ xor } c$
 - $a/bc+/abc+/a/b/c+ab/c$
 - $a \text{ xor } b \text{ xor } c \text{ xor } d$
 - $a/bcd+/abcd+/a/b/cd+ab/cd+/ab/c/d+a/b/c/d+abc/d+/a/bc/d$
- Compare
- $a \text{ xor } b$
 - $x1=a/b+/ab$
 - $a \text{ xor } b \text{ xor } c$
 - $x2=x1/c+/x1*c$
 - $a \text{ xor } b \text{ xor } c \text{ xor } d$
 - $x3=x2/d+/x2*d$

Why Multilevel

- $a \text{ xor } b$
 - $x1 = a/b + ab$
- $a \text{ xor } b \text{ xor } c$
 - $x2 = x1/c + x1 * c$
- $a \text{ xor } b \text{ xor } c \text{ xor } d$
 - $x3 = x2/d + x2 * d$
- Multi-level
 - exploit common sub-expressions
 - linear complexity
- Two-level
 - exponential complexity

Penn ESE353 Spring 2015 -- DeHon

7

Harder than Two-Level

- Two-level already NP-hard
 - has canonical representation
 - clean formulation
 - observed can limit to Primes
 - identified opportunities for pruning
- Multi-level has more flexibility
 - \rightarrow larger space to explore
 - Not formulated cleanly
- Solution more heuristic ... art
 - ... all problems start this way, some stay...

Penn ESE353 Spring 2015 -- DeHon

8

Goal

- Find the structure
- Exploit to minimize gates
 - Total (area)
 - In path (delay)

Penn ESE353 Spring 2015 -- DeHon

9

Multi-level Transformations

- Decomposition
- Extraction
- Factoring
- Substitution
- Collapsing

[copy these to board so stay up as we move forward]

Penn ESE353 Spring 2015 -- DeHon

10

Decomposition

- $F = abc + abd + a/c/d + b/c/d$
- $F = XY + /X/Y$
- $X = ab$
- $Y = c + d$

Penn ESE353 Spring 2015 -- DeHon

11

Decomposition

- $F = abc + abd + a/c/d + b/c/d$
 - 4 3-input + 1 4-input \rightarrow 11 2-input gates
- $F = XY + /X/Y$
- $X = ab$
- $Y = c + d$
 - 5 2-input gates
- Note: use X and /X, use at multiple places

Penn ESE353 Spring 2015 -- DeHon

12

Extraction

- $F=(a+b)cd+e$
- $G=(a+b)/e$
- $H=cde$
- $F=XY+e$
- $G=X/e$
- $H=Ye$
- $X=a+b$
- $Y=cd$

Extraction

- $F=(a+b)cd+e$
- $G=(a+b)/e$
- $H=cde$
- 2-input: 4
- 3-input: 2
- $F=XY+e$
- $G=X/e$
- $H=Ye$
- $X=a+b$
- $Y=cd$
- 2-input: 6

→ 8 2-input gates

Common sub-expressions over **multiple output**

Factoring

- $F=ac+ad+bc+bd+e$
- $F=(a+b)(c+d)+e$

Factoring

- $F=ac+ad+bc+bd+e$
 - 4 2-input, 1 5-input → 8 2-input gates
 - 9 literals
- $F=(a+b)(c+d)+e$
 - 4 2-input
 - 5 literals

Substitution

- $G=a+b$
- $F=a+bc$
- Substitute G into F
- $F=G(a+c)$
 - (verify) $F=(a+b)(a+c)=aa+ab+ac+bc=a+bc$
- useful if also have $H=a+c$, then $F=GH$

Collapsing

- $F=Ga+Gb$
- $G=c+d$
- $F=ac+ad+b/c/d$
- opposite of substitution
 - sometimes want to collapse and refactor
 - especially for delay optimization [next lecture]

Moves

- These transforms define the “moves” we can make to modify our network.
- Goal is to apply, usually repeatedly, to minimize gates
 - ...then apply as necessary to accelerate design
- MIS/SIS
 - Applies to canonical 2-input gates
 - Then covers with target gate library
 - Day 3

Division

Division

- **Given:** function (f) and divisor (p)
- **Find:** quotient (q) and remainder (r)
 $f=pq+r$

E.g.

$$f=abc+abd+ef, p=ab$$
$$q=c+d, r=ef$$

Algebraic Division

- Use basic rules of algebra, rather than full boolean properties
- Computationally simple
- Weaker than boolean division
- $f=a+bc$ $p=(a+b)$
- **Algebra:** not divisible
- **Boolean:** $q=(a+c)$, $r=0$

Algebraic Division

- Given:** function (f) and divisor (p)
- Find:** quotient and remainder
 $f=pq+r$
- f and p are expressions (lists of cubes)
 - $p=\{a_1, a_2, \dots\}$
- Define: $h_i = \{c_j \mid a_i * c_j \in f\}$
- $f/p = h_1 \cap h_2 \cap h_3 \dots$

Algebraic Division Example (adv to alg.; work ex on board)

- $f=abc+abd+de$
- $p=ab+e$

Algebraic Division

- f and p are expressions (lists of cubes)
- $p = \{a_1, a_2, \dots\}$
- $h_i = \{c_j \mid a_i * c_j \in f\}$
- $f/p = h_1 \cap h_2 \cap h_3 \dots$

Algebraic Division Example

- $f = abc + abd + de$, $p = ab + e$
- $p = \{ab, e\}$
- $h_1 = \{c, d\}$
- $h_2 = \{d\}$
- $h_1 \cap h_2 = \{d\}$
- $f/p = d$
- $r = f - p * (f/p)$
- $r = abc + abd + de - (ab + e)d$
- $r = abc$

Algebraic Division Time

- $O(|f||p|)$ as described
 - compare every cube pair
- Sort cubes first
 - $O((|f|+|p|)\log(|f|+|p|))$

Primary Divisor

- f/c such that c is a cube
- $f = abc + abde$
- $f/a = bc + bde$ is a primary divisor

Cube Free

- The only cube that divides p is 1
- $c + de$ is cube free
- $bc + bde$ is not cube free

Kernel

- Kernels of f are
 - cube free primary divisors of f
 - *Informally*: sums w/ cubes factored out
- $f = abc + abde$
- $f/ab = c + de$ is a kernel
- ab is **cokernel** of f to $(c + de)$
 - cokernels always cubes

Factoring

- Gfactor(f)
 - if (terms==1) return(f)
 - p=CHOOSE_DIVISOR(f)
 - (h,r)=DIVIDE(f,p)
 - f=Gfactor(h)*Gfactor(p)+Gfactor(r)
 - return(f) // factored

Factoring

- Trick is picking divisor
 - pick from kernels
 - goal minimize literals **after** resubstitution
 - Re-express design using new intermediate variables
 - Variable and complement

Kernel Extraction

- Kernel1(j,g)
 - R=g
 - N max index in g
 - for(i=j+1 to N)
 - if (l_i in 2 or more cubes)
 - c_i=largest cube divide g/l_i
 - if (forall k≤i, l_k∉c_i)
 - » R=R ∪ KERNEL1(i,g/(l_i∩c_i))
 - return(R)
- Must be to Generate Non-trivial kernel

- Find c_f = largest cube factor of f
- K=Kernel1(0,f/c_f)
- if (f is cube-free)
 - return(f∪K)
- else
 - return(K)

Consider each literal for cokernel once
(largest cokernels will already have been found)

Kernel Extract Example (ex. on board; adv to return to alg.)

- f=abcd+abce+abef

Kernel Extraction

- Kernel1(j,g)
 - R=g
 - N max index in g
 - for(i=j+1 to N)
 - if (l_i in 2 or more cubes)
 - c_i=largest cube divide g/l_i
 - if (forall k≤i, l_k∉c_i)
 - » R=R ∪ KERNEL1(i,g/(l_i∩c_i))
 - return(R)
- Must be to Generate Non-trivial kernel

- Find c_f = largest cube factor of f
- K=Kernel1(0,f/c_f)
- if (f is cube-free)
 - return(f∪K)
- else
 - return(K)

Consider each literal for cokernel once
(largest cokernels will already have been found)

Kernel Extract Example (stay on prev. slide, ex. on board)

- f=abcd+abce+abef
- c_f=ab
- f/c_f=cd+ce+ef
- R={cd+ce+ef}
- N=6
- a,b not present
- (cd+ce+ef)/c=e+d
- largest cube 1
- Recurse→e+d
- R={cd+ce+ef,e+d}
- only 1 d
- (d+ce+ef)/e=c+f
- Recurse→c+f
- R={cd+ce+ef,e+d,c+f}

Extraction

Identify cube-free expressions in many functions
(common sub expressions)

1. Generate kernels for each function
2. select pair such that $k1 \cap k2$ is not a cube
 - Note: $k1=k2$ is simplest case of this
 - ...but intersection case is more powerful
 - Example to come
3. new variable from intersection
 - $v = k1 \cap k2$
4. update functions (resubstitute)
 - $f_i = v * (f_i / v) + r_i$
 - (similar for common cubes)

Extraction Example

- $X = ab(c(d+e)+f+g)+g$
- $Y = ai(c(d+e)+f+j)+k$

Extraction Example

- $X = ab(c(d+e)+f+g)+g$
- $Y = ai(c(d+e)+f+j)+k$
- $d+e$ kernel of both
- $L = d+e$
- $X = ab(cL+f+g)+h$
- $Y = ai(cL+f+j)+k$

Extraction Example

- $L = d+e$
- $X = ab(cL+f+g)+h$
- $Y = ai(cL+f+j)+k$
- kernels: $(cL+f+g), (cL+f+j)$
- extract: $M = cL+f$
- $X = ab(M+g)+h$
- $Y = ai(M+f)+h$

Extraction Example

- | | |
|---------------------|-------------------|
| • $L = d+e$ | • $N = aM$ |
| • $M = cL+f$ | • $M = cL+f$ |
| • $X = ab(M+g)+h$ | • $L = d+e$ |
| • $Y = ai(M+j)+h$ | • $X = b(N+ag)+h$ |
| • no kernels | • $Y = i(N+aj)+k$ |
| • common cube: aM | |

Extraction Example

- | | |
|-------------------|---|
| • $N = aM$ | • Can collapse |
| • $M = cL+f$ | - L into M into N |
| • $L = d+e$ | - Only used once |
| • $X = b(N+ag)+h$ | • Get larger common kernel N |
| • $Y = i(N+aj)+k$ | - maybe useful if components becoming too small for efficient gate implementation |

Resubstitution

- Also useful to try complement on new factors
- $f=ab+ac+b/cd$
- $X=b+c$
- $f=aX+b/cd$
- $/X=b/c$
- $f=aX+/Xd$
- ...extracting complements not a direct target

Multilevel Optimization

- Unlike Two-level, very heuristic
- Not clear when done
- Goal find common terms to share
- Often start with two-level optimization
 - Identifies product term sharing
- Identify kernels and cubes
- Factor them out
- If can be used many places, get benefit
- SIS included common recipes
- More after timing analysis

Summary

- Want to exploit structure in problems to reduce (contain) size
 - common sub-expressions
- Identify component elements
 - decomposition, factoring, extraction
- Division key to these operations
- Kernels give us divisors

Big Ideas

- Exploit freedom
 - form
- Exploit structure/sharing
 - common sub expressions
- Techniques
 - Iterative Improvement
 - Refinement/relaxation

Admin

- Reading for Monday on canvas
- Updated energy model:
update_model.tar
- Milestone Thursday