



Duke Architecture

# Navigating Heterogeneous Processors with Market Mechanisms

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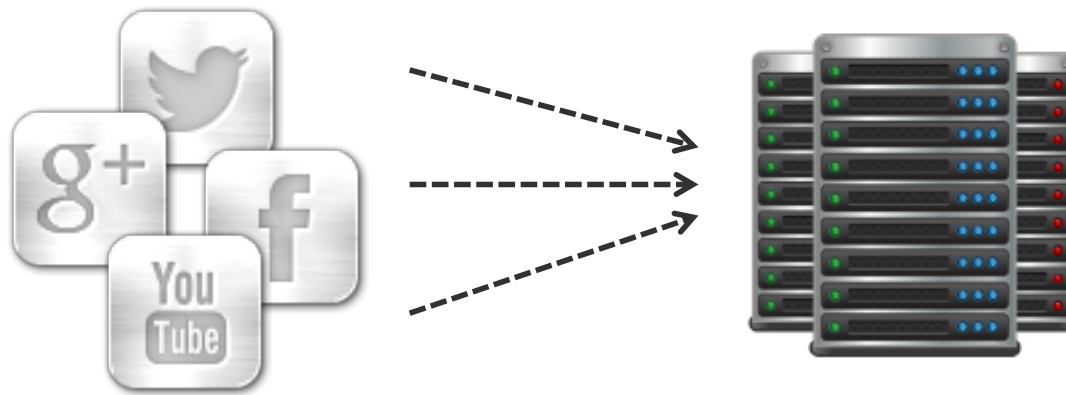
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# Growth of Cloud Applications



- Data volumes are growing geometrically
- Cloud applications are diversifying rapidly
- Computing capability must grow
- Datacenters consume tens of megawatts of power



# Datacenter Scaling Limitations

- Dennard scaling is over
  - On-chip power density is a constraint
- Adding servers is expensive
  - Power determines operating costs
- Heterogeneity improves energy efficiency
  - Small cores consume a fraction of big core power
  - Big cores ensure service quality



# Executive Summary

Goal:

- Improve service quality in heterogeneous datacenters

Methods:

- Leverage processor heterogeneity
- Mitigate performance risk using market

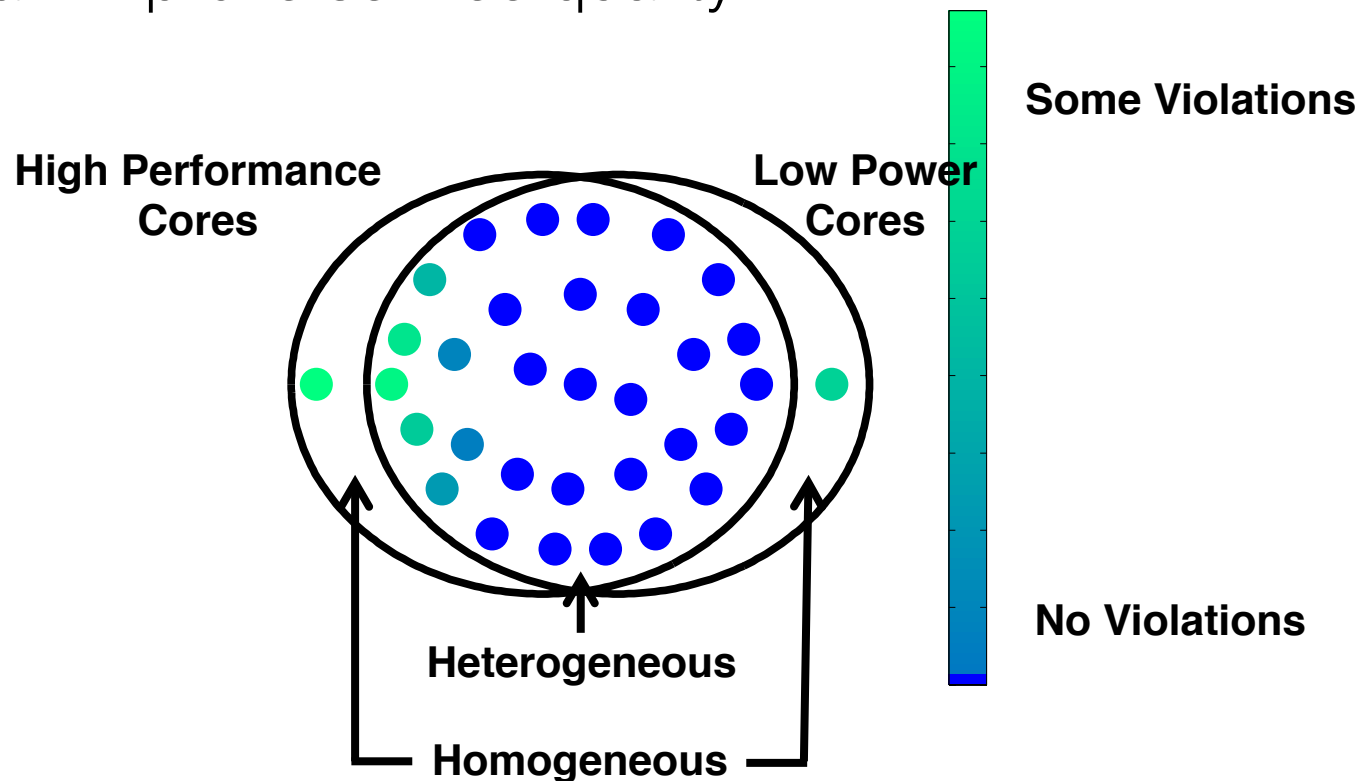
Evaluation:

- Big/Small core heterogeneity improves service quality
- Three core types reduce service violations by 12x



# Risks of Heterogeneity

- Heterogeneity introduces performance risk
- Yet it can improve service quality



# Managing Risk

1. Types of processors?
  2. Number of each?
  3. How to allocate?
- Resource allocation that mitigates performance risk
    - Hide hardware complexity
    - Trade-off performance and power
    - Allocate small cores when possible
  - Coordinate design and management

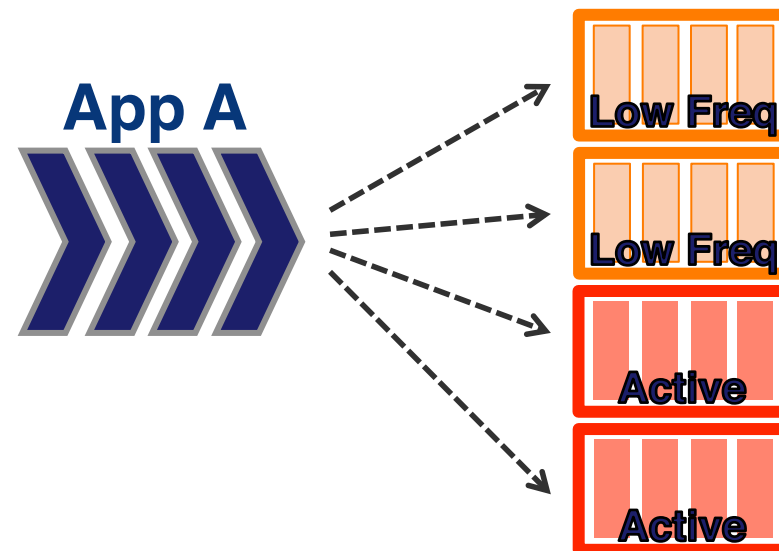


# Allocating Time

Market periodically allocates time on hardware resource

$$\frac{\text{Tasks}}{\text{Sec}} = \frac{\text{Tasks}}{\text{Cycles}} \times \frac{\text{Cycles}}{\text{Sec}}$$

Right-size datacenter via server activation and DVFS

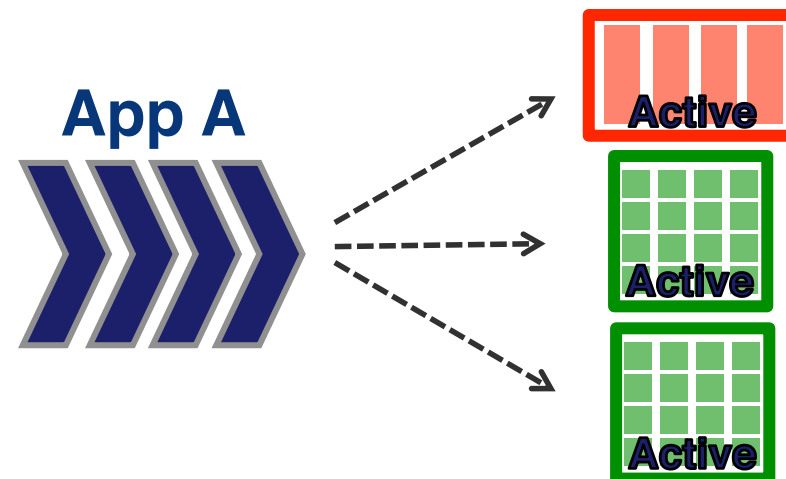


# Accommodating Heterogeneity

Profile task-specific performance on each processor:

$$\frac{\text{Tasks}}{\text{Sec}} = \frac{\text{Tasks}}{\text{Inst}} \times \boxed{\frac{\text{Insts}}{\text{Cycle}}} \times \frac{\text{Cycles}}{\text{Sec}}$$

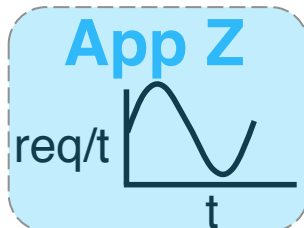
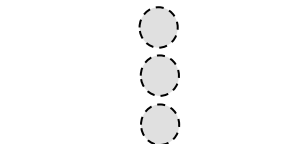
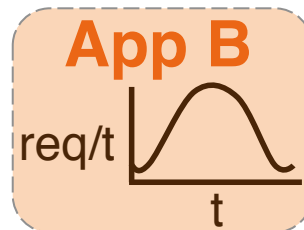
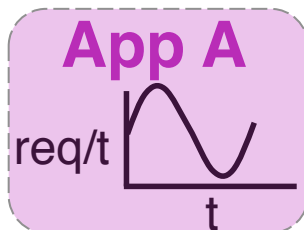
Store  $IPC_{a,m}$  as scaling factor relative to baseline  $IPC_0$





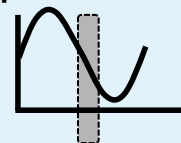
# Market Resource Allocation

$$\text{maximize } \sum_{a \in \text{App}} (\text{Value}_a - \text{Cost})$$



## Proxy

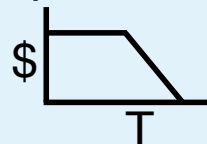
$\lambda \leftarrow$  predict demand



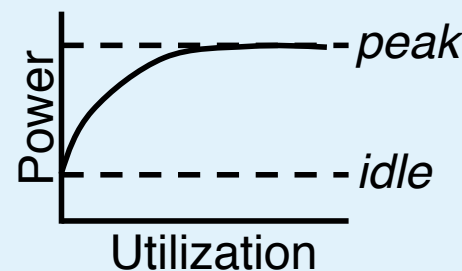
$T \leftarrow$  predict wait time

$\lambda \rightarrow$    $\rightarrow \mu$

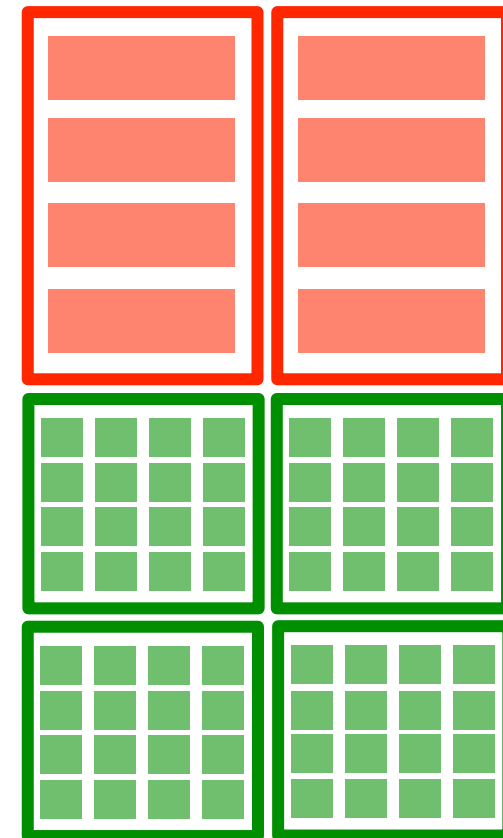
$V \leftarrow$  predict value



## Cost Model



$C \leftarrow$  energy  $\times$  price



# Experimental Methodology

Real  
tasks

Diurnal  
arrivals

Java implementation,  
CPLEX solver for optimization

System profiles

Simulation

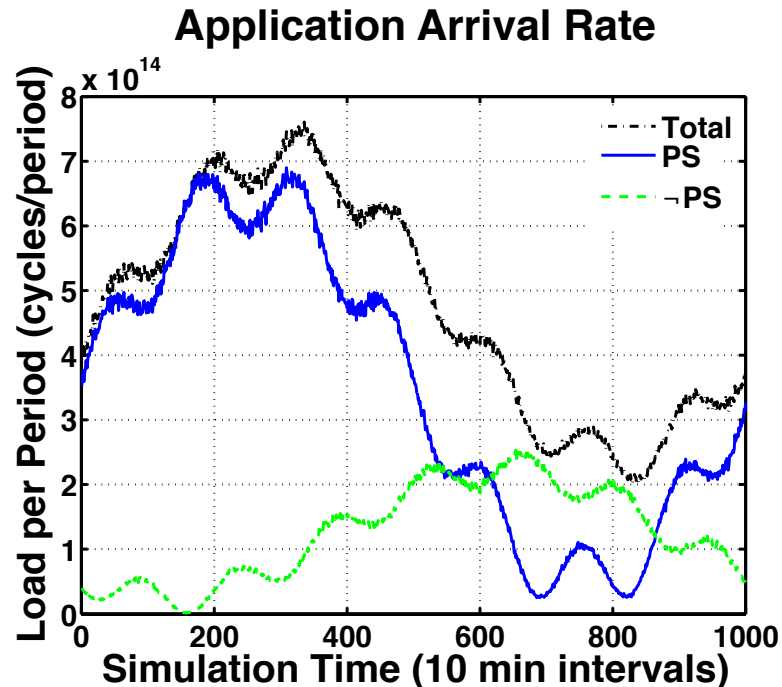


# Defining Big and Small Cores

- Within fixed power budget, vary number of:
  - 4-core Xeon servers
  - 16-core Atom servers
- Core measurements [ISCA'10]
  - 0.3 – 1.0 relative IPC
  - 1.5 W Atom vs. 15 W Xeon
- System model
  - Equal die area
  - Fixed system power overhead (65 W)



# Modeling Application Behavior



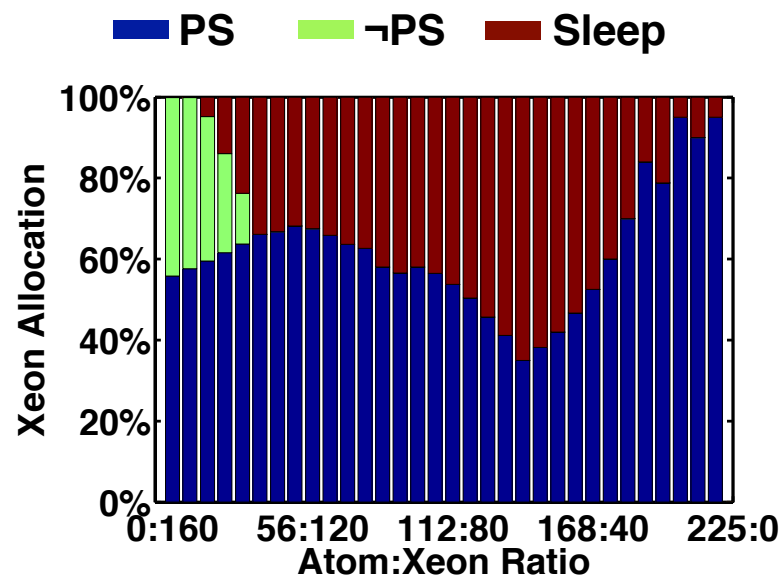
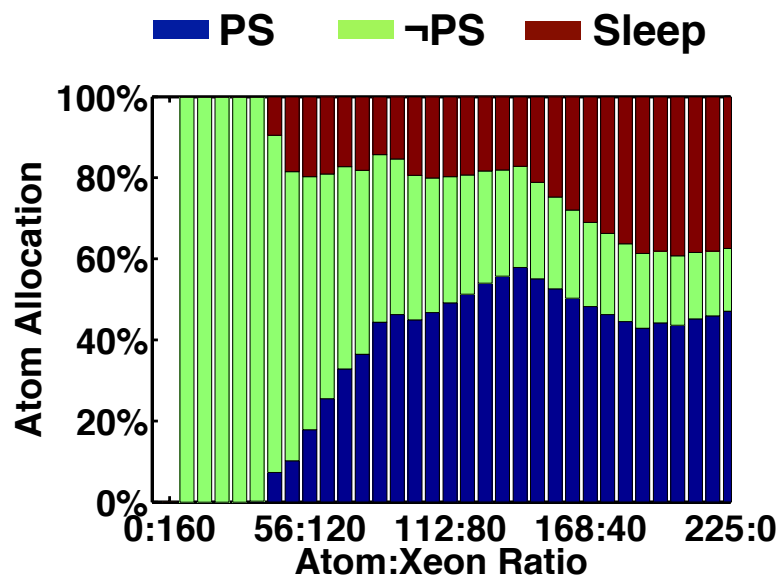
- One week of requests
- Diurnal pattern

- Processor Sensitive (PS) – Atom throughput  $\frac{1}{3}$  that of Xeon
- Processor Insensitive ( $\neg$ PS) – same throughput



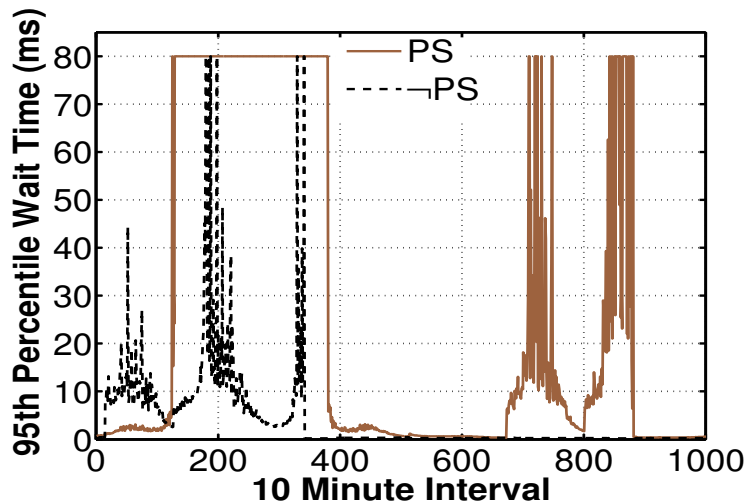
# Understanding Datacenter Dynamics

- Vary Atom to Xeon ratio
- Examine allocations to each task
- Identify a balanced mix (e.g. 147:55)

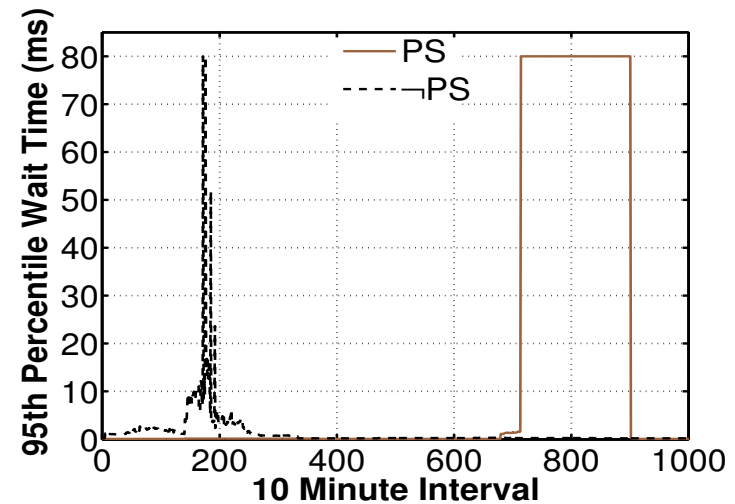


# Improving Service Quality

## Homogeneous Xeon



## Homogeneous Atom

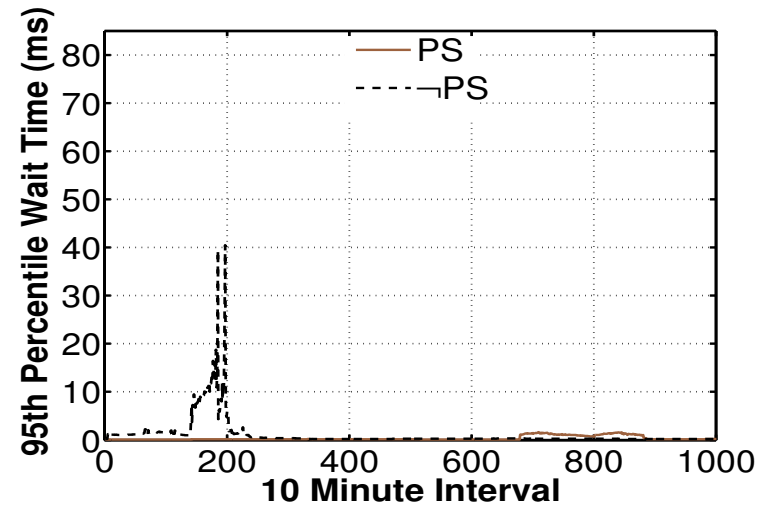


- Xeon-only has insufficient resources
- Atom-only incurs violations due to cost



# Improving Service Quality

## Heterogeneous Xeon/Atom



- Xeon/Atom mix reduces waiting time
  - Atoms mostly allocated to ¬PS
  - Xeons freed to service PS peaks



# Defining Greater Heterogeneity

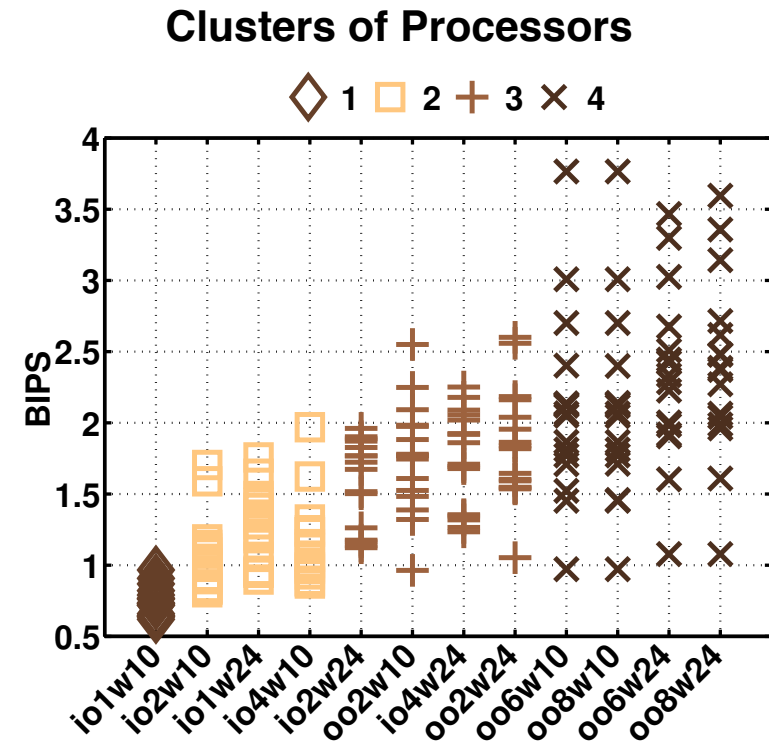
- Within fixed power budget, vary core designs
  - Dynamic scheduling (IO vs OO)
  - Issue Width (1,2,4,6,8)
  - Frequency (1.0,2.4 GHz)
- Processor simulation
  - 0.4 – 1.5 relative IPC (gem5)
  - 1.1 W – 28 W (McPAT)
- System model
  - Equal die area
  - Fixed system power overhead (65 W)





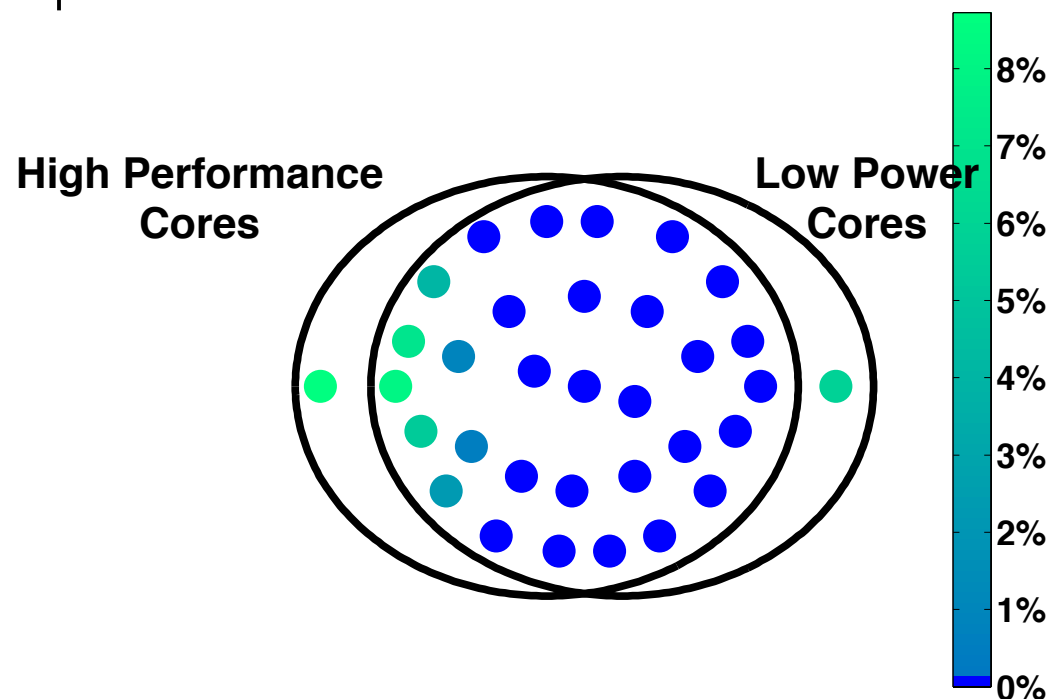
# Clustering Heterogeneous Processors

- Cluster cores with similar SPEC performance
- Select core with lowest performance variation from each cluster
- Evaluate with diverse SPEC task streams



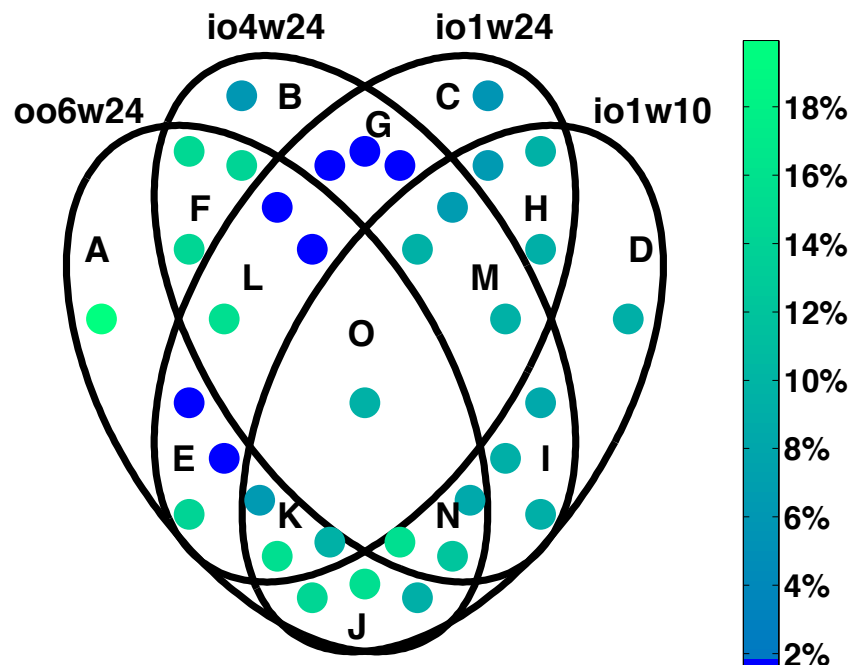
# Visualizing the Design Space

- Ellipses represent core types
- Points are combinations of cores
- Colors represent service violations



# Design Space with Four Processors

Service Quality Violations



- Identify right core types
- Prune design space
- Best configuration is heterogeneous
  - RT violations reduced from 15.5% to 1.6%

# Conclusions and Future Directions

- Leverage market to mitigate heterogeneity's risk
  - Embed microarchitectural insight into the market
  - Allocate multiple resources
- Deploy heterogeneous hardware in a datacenter
  - Optimal balance improves service
  - Sophisticated trade-offs require further study
- Propose a datacenter research methodology
  - Simulating detailed server architecture
  - Modeling user and datacenter dynamics
  - Tractability for web search, memcached, map/reduce





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